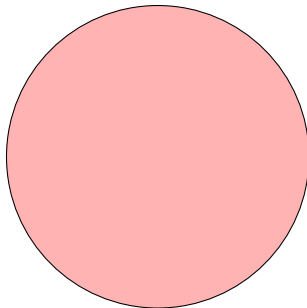


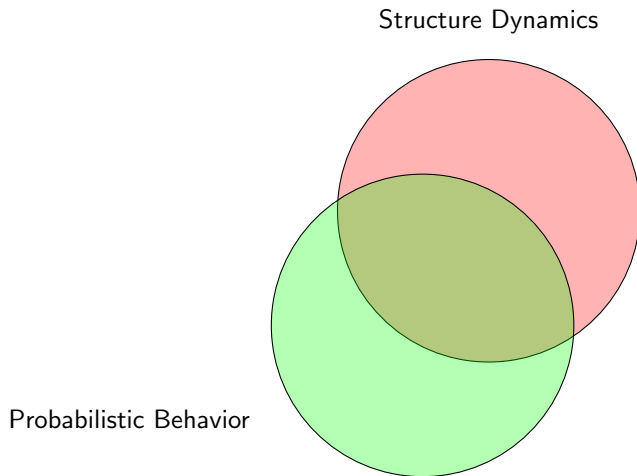
# Probabilistic Timed Graph Transformation Systems

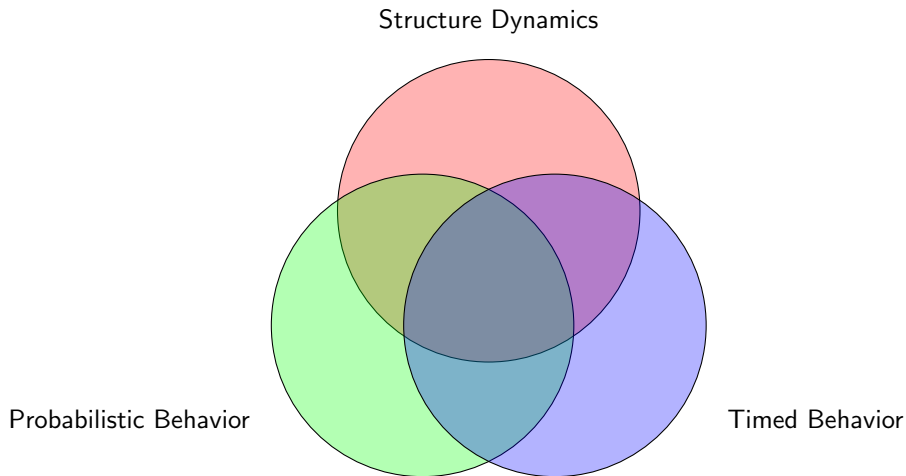
**Maria Maximova**, Holger Giese, Christian Krause

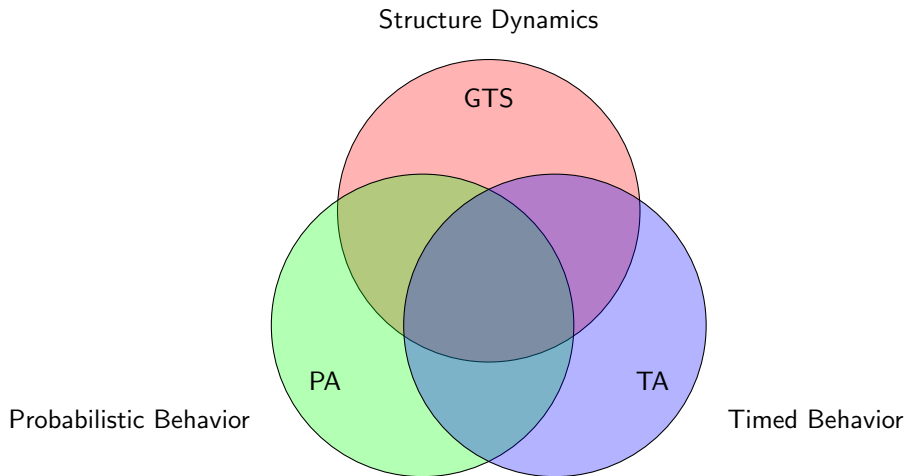
19th July 2017

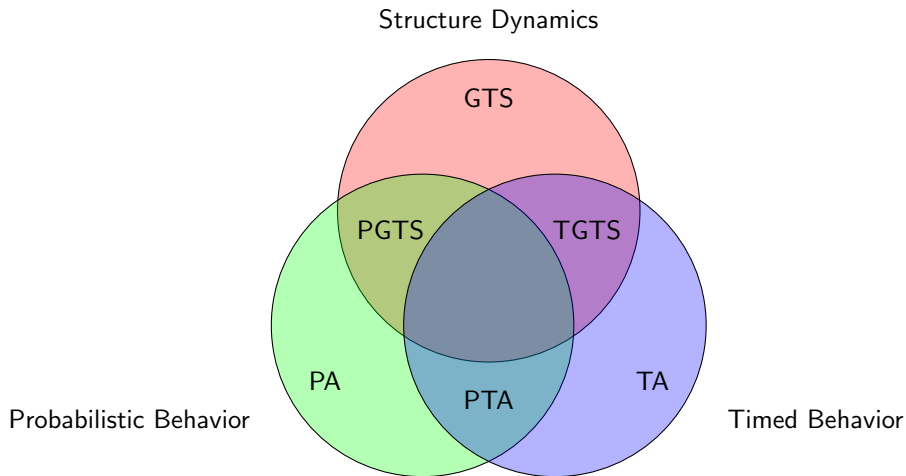
## Structure Dynamics

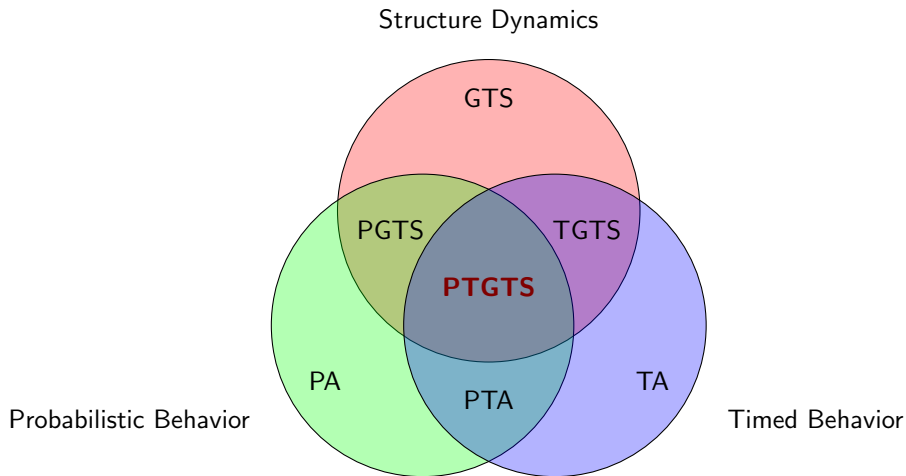


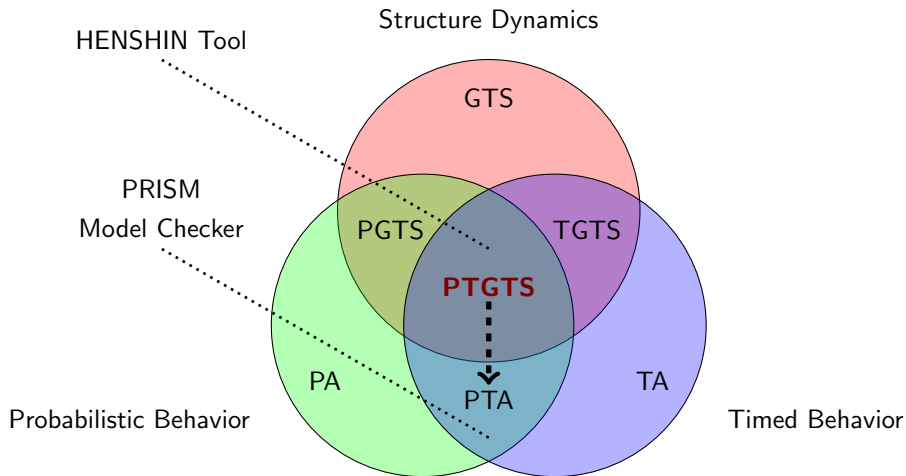












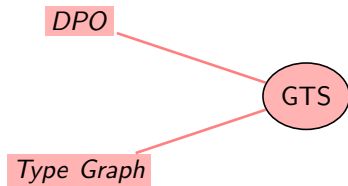


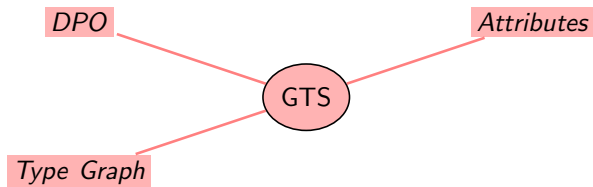
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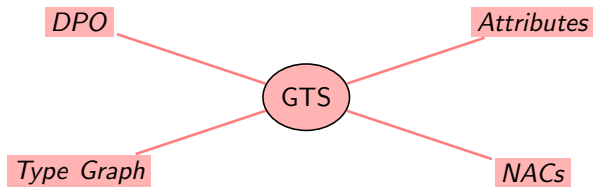
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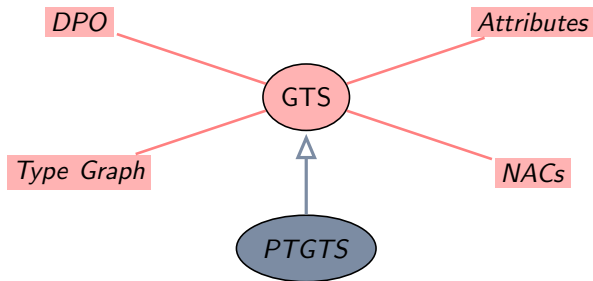




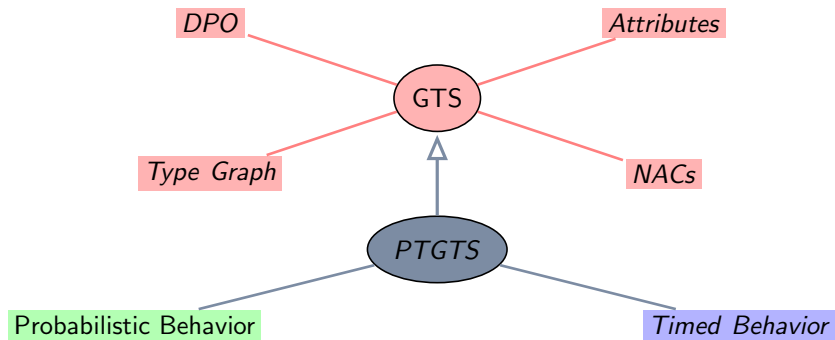


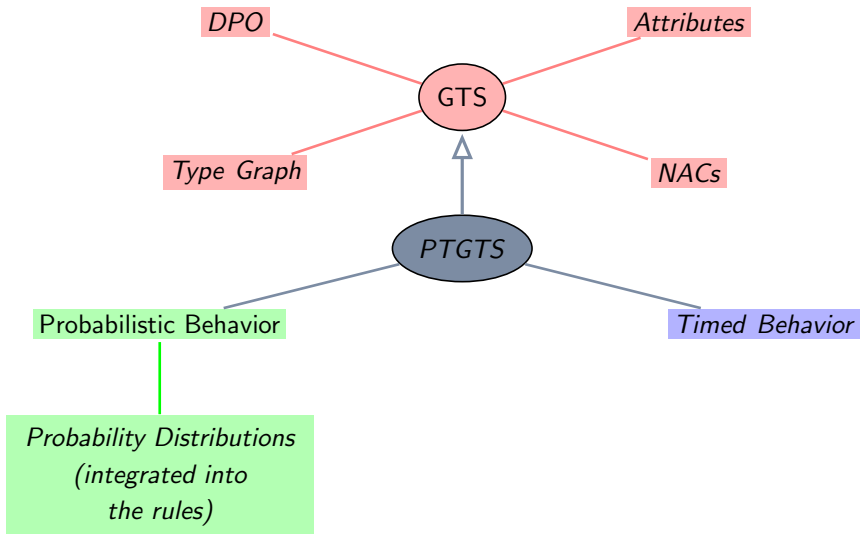


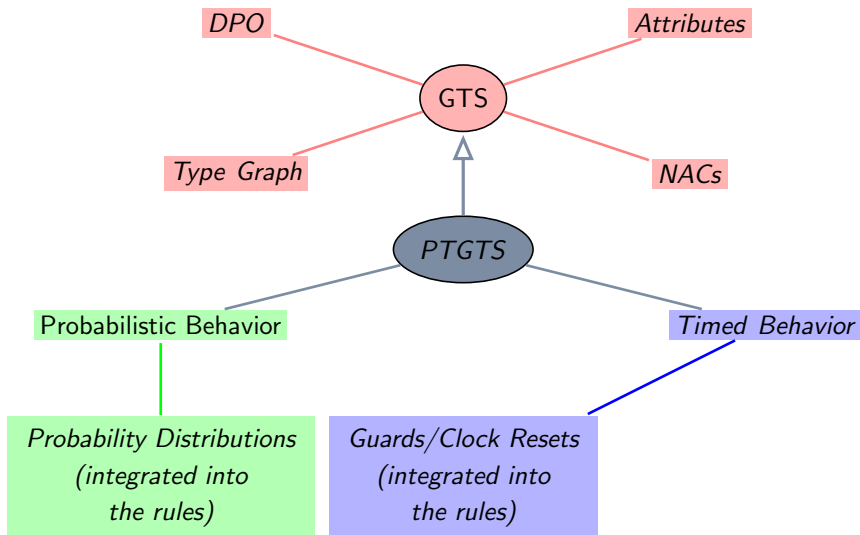


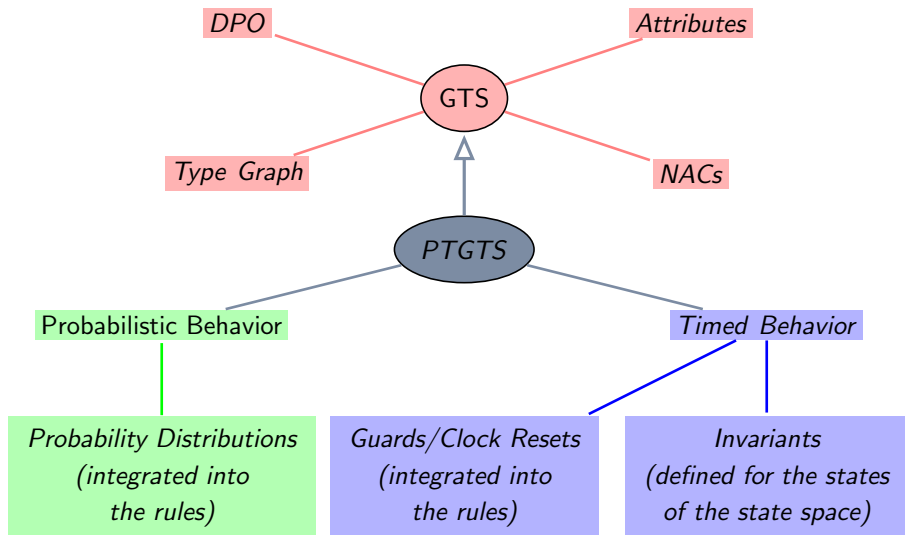












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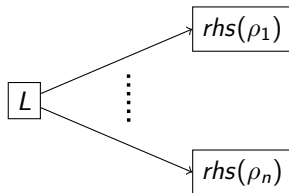
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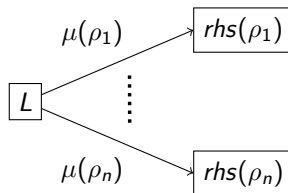
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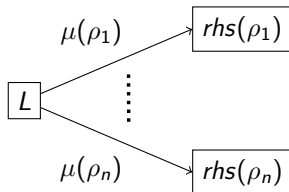


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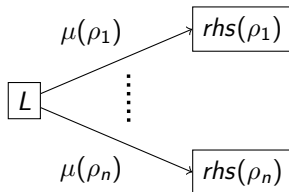


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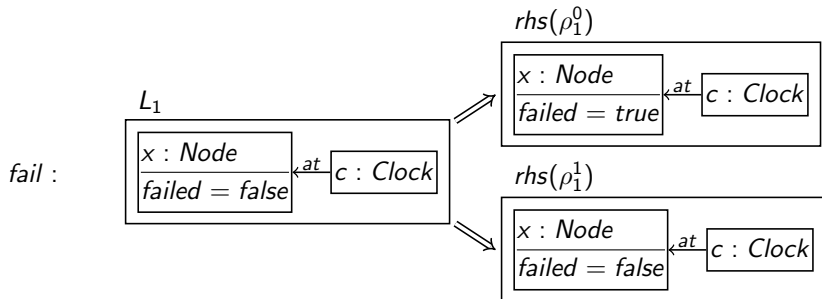
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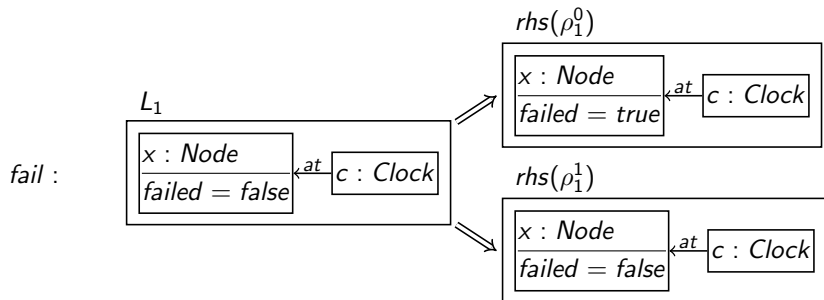
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- $r_C \subseteq \text{ClockNodes}(L)$  is a set of nodes of the type *Clock* to be reset.



$$P = \{\rho_1, \dots, \rho_n\}$$





$fail = (L_1, P_1, \mu_1, \phi_1, r_{C_1})$  with

$P_1 = \{\rho_1^0, \rho_1^1\}$

$\mu_1 = \{(\rho_1^0, 0.1), (\rho_1^1, 0.9)\}$

$\phi_1 = (c \geq 2)$

$r_{C_1} = \{c\}$

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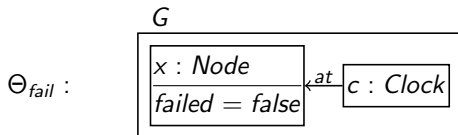
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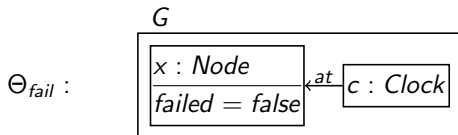
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$$\Theta_{fail} = (G, \phi) \text{ with } \phi = (c \leq 5)$$

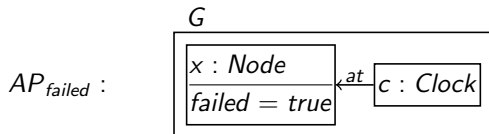
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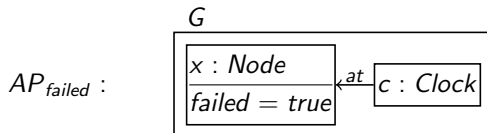
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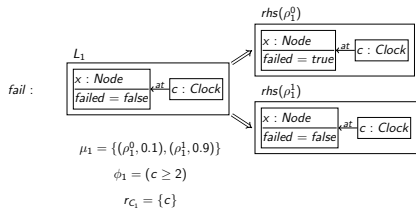


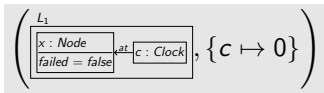
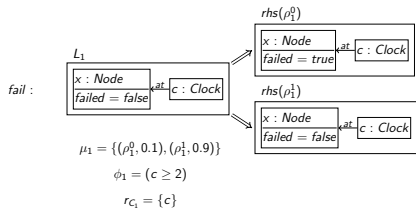
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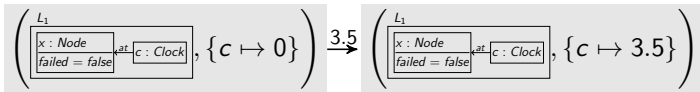
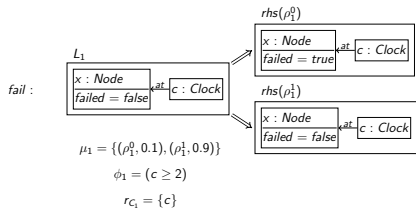


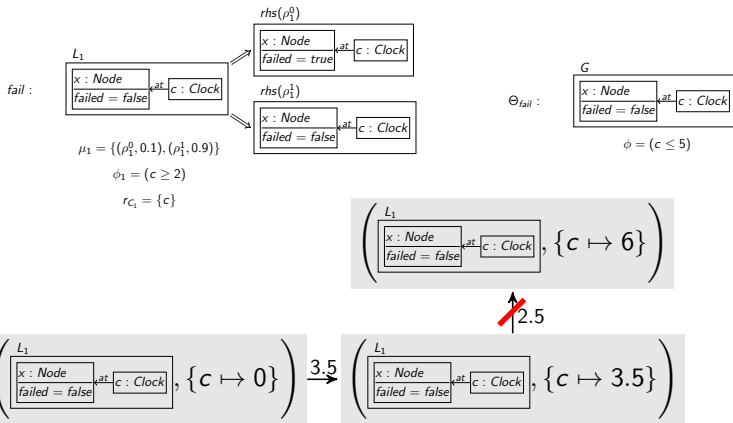
$$AP_{failed} = (G, \phi) \text{ with } \phi = true$$

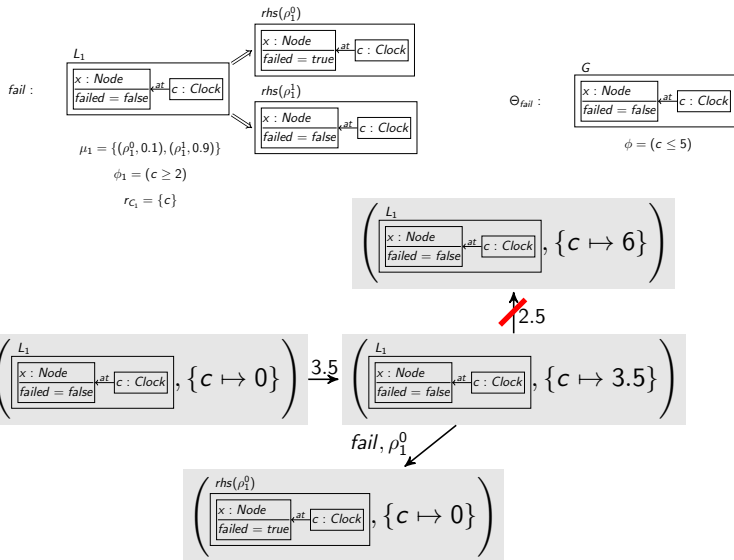












$\left( \begin{array}{c} L_1 \\ \boxed{x : \text{Node} \\ \text{failed} = \text{false}} \leftarrow_{\text{at}} c : \text{Clock} \end{array} , \{c \mapsto 6\} \right)$

$\uparrow$  ~~2.5~~

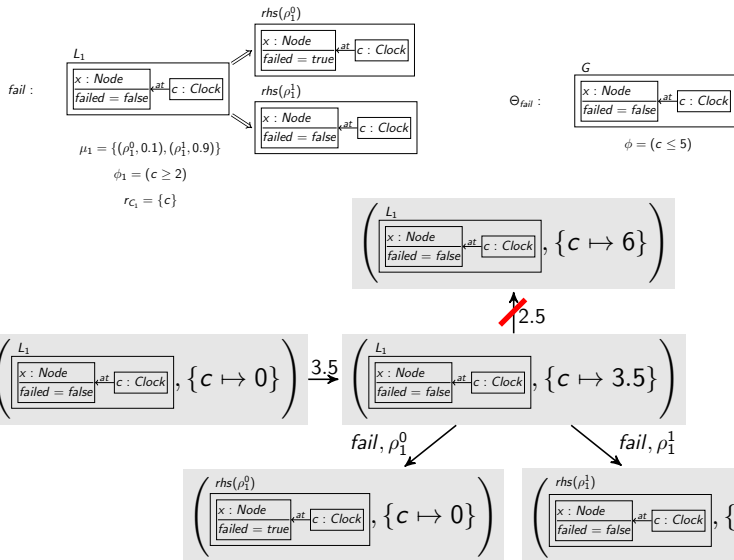
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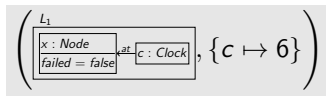
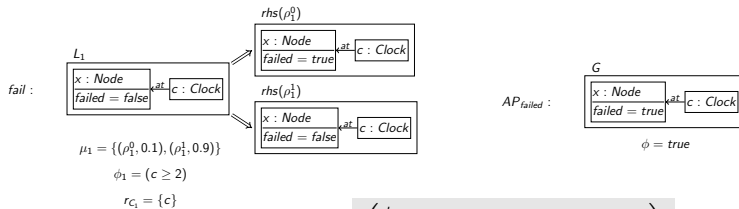
$\rightarrow$

$\left( \begin{array}{c} L_1 \\ \boxed{x : \text{Node} \\ \text{failed} = \text{false}} \leftarrow_{\text{at}} c : \text{Clock} \end{array} , \{c \mapsto 3.5\} \right)$

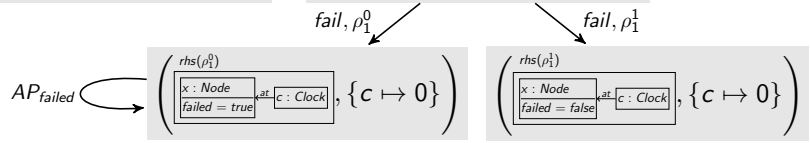
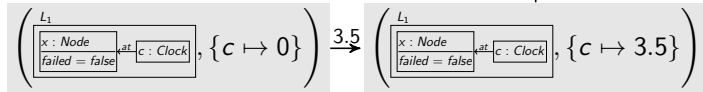
*fail*,  $\rho_1^0$

$\left( \begin{array}{c} rhs(\rho_1^0) \\ \boxed{x : \text{Node} \\ \text{failed} = \text{true}} \leftarrow_{\text{at}} c : \text{Clock} \end{array} , \{c \mapsto 0\} \right)$



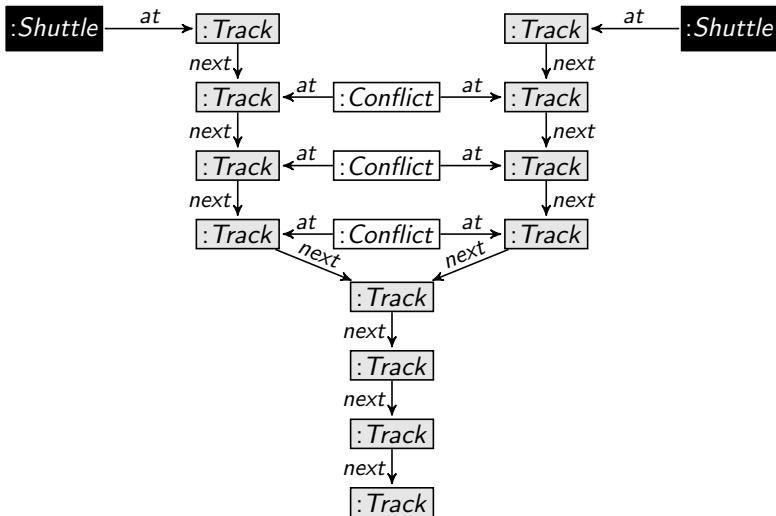


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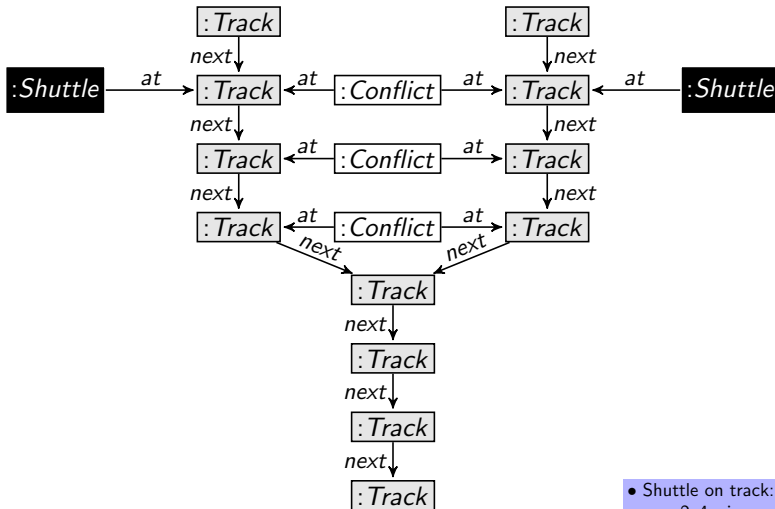
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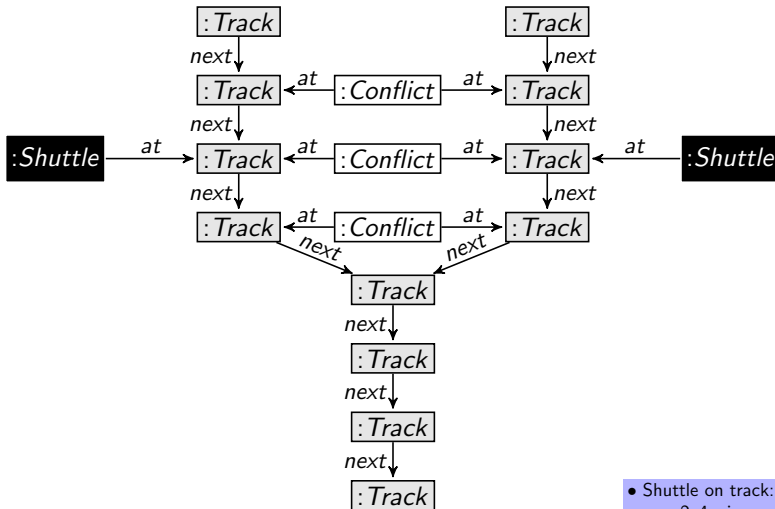






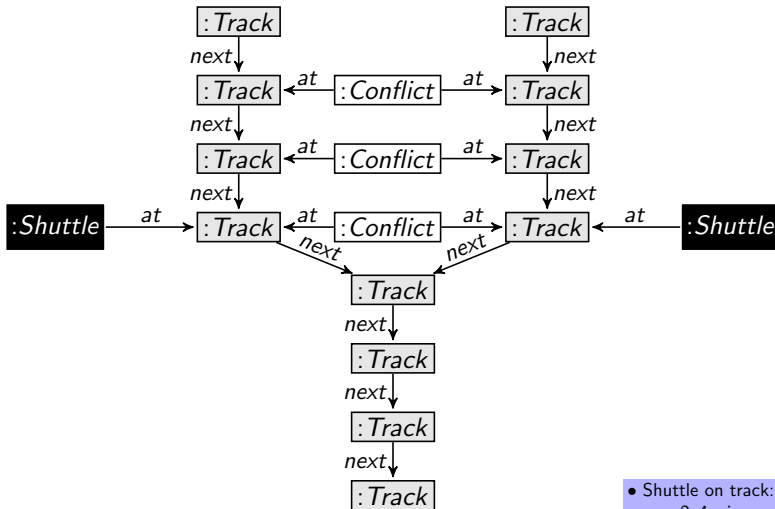
- Shuttle on track:  
2–4 min





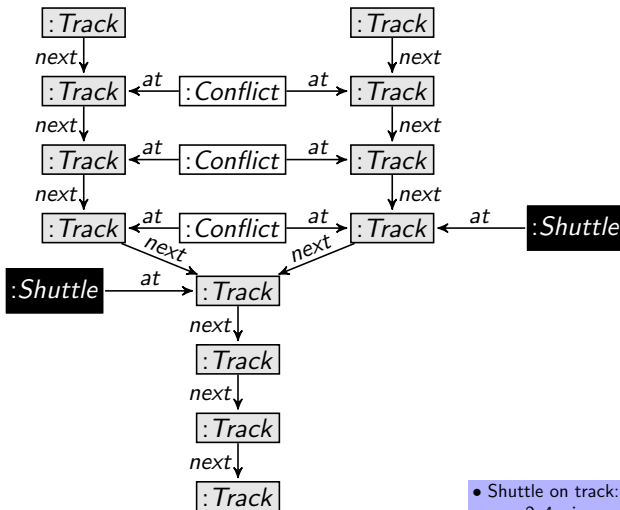
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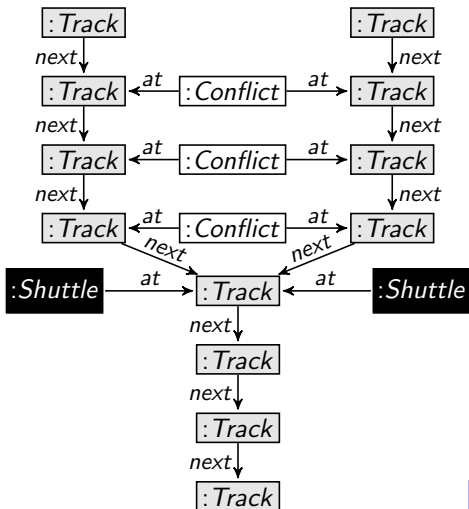




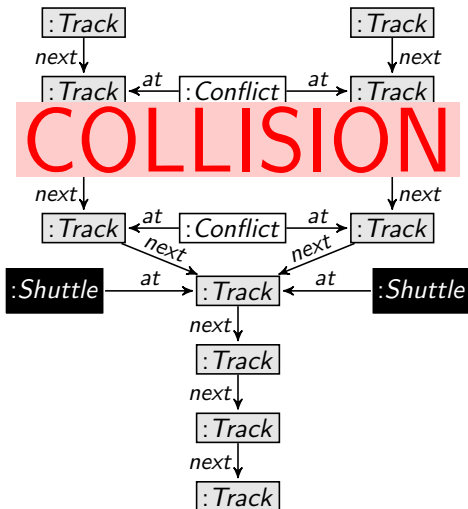
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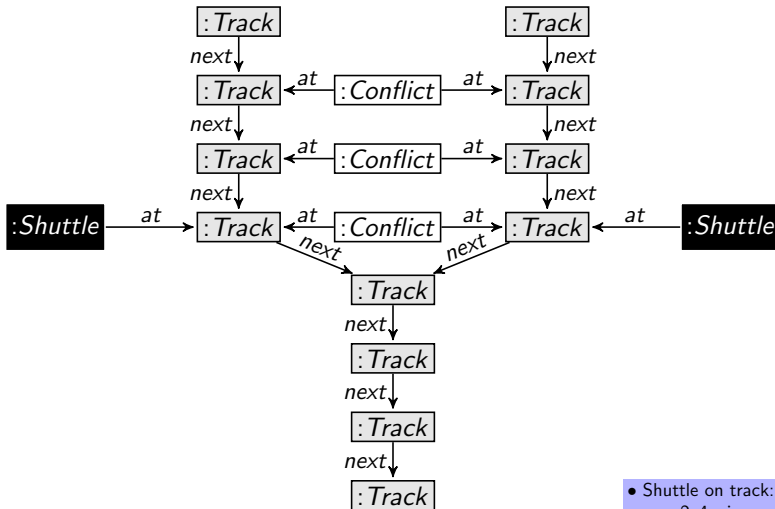


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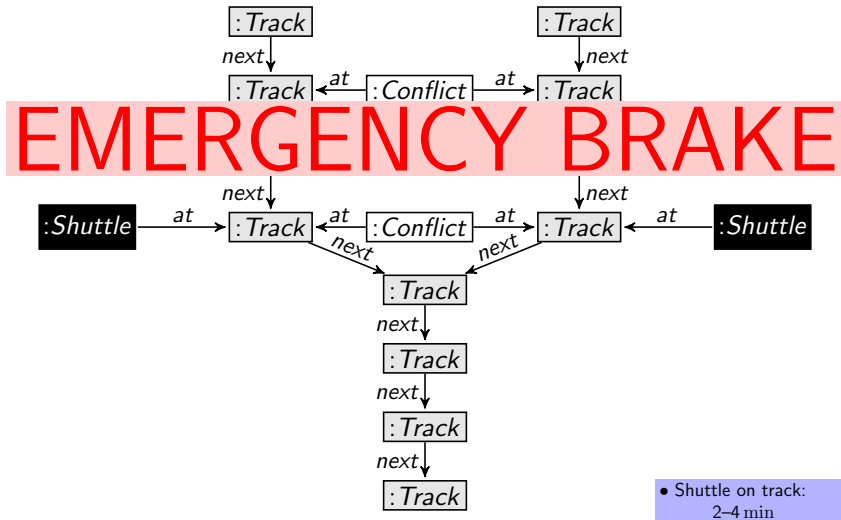


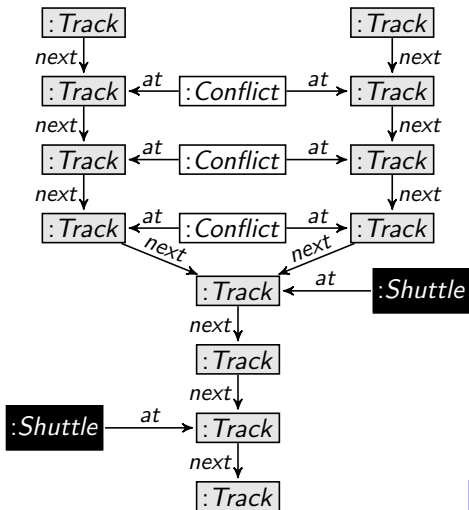
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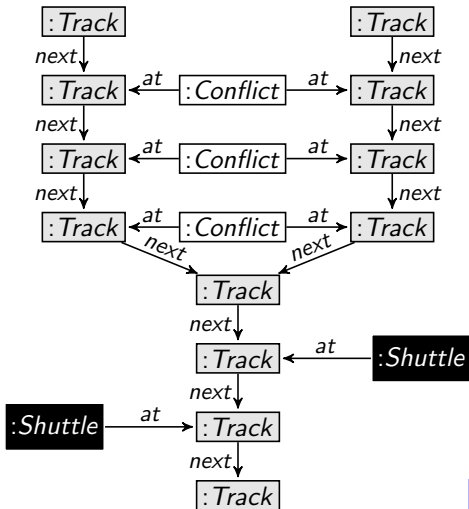


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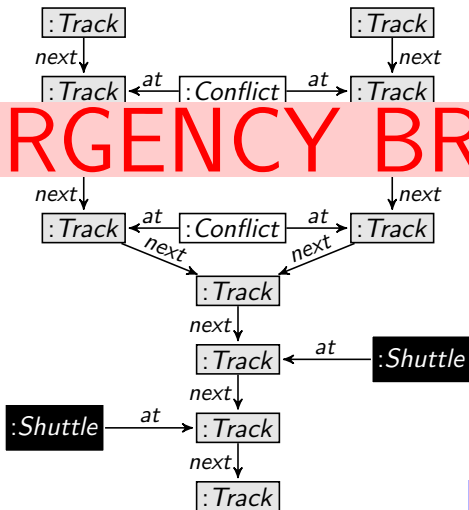
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## EMERGENCY BRAKE



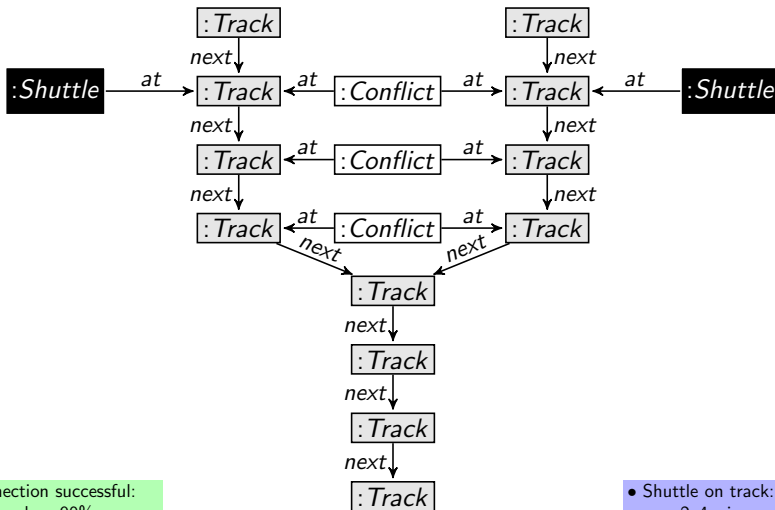
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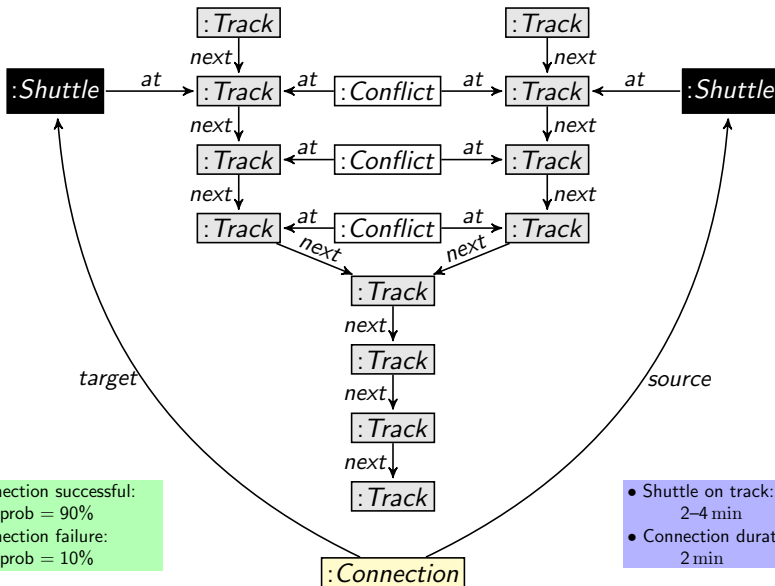




- Connection successful:  
prob = 90%

- Shuttle on track:  
2-4 min
- Connection duration:  
2 min

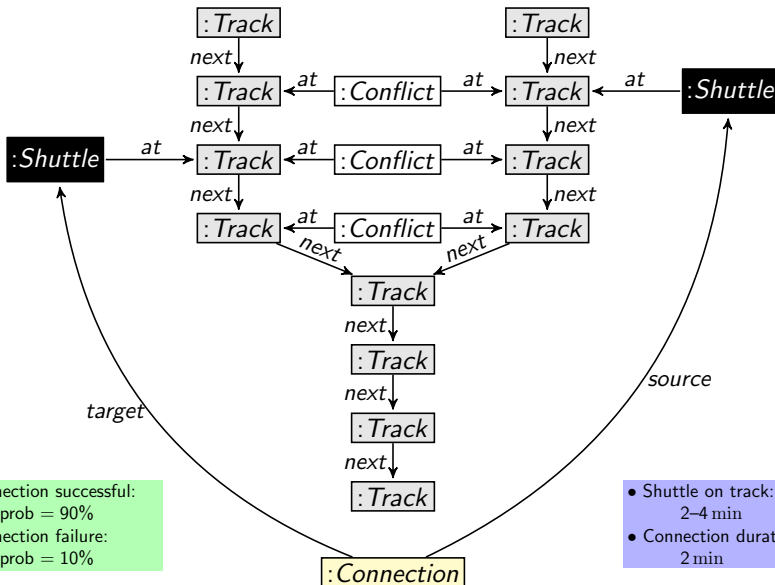




- Connection successful:  
prob = 90%
- Connection failure:  
prob = 10%

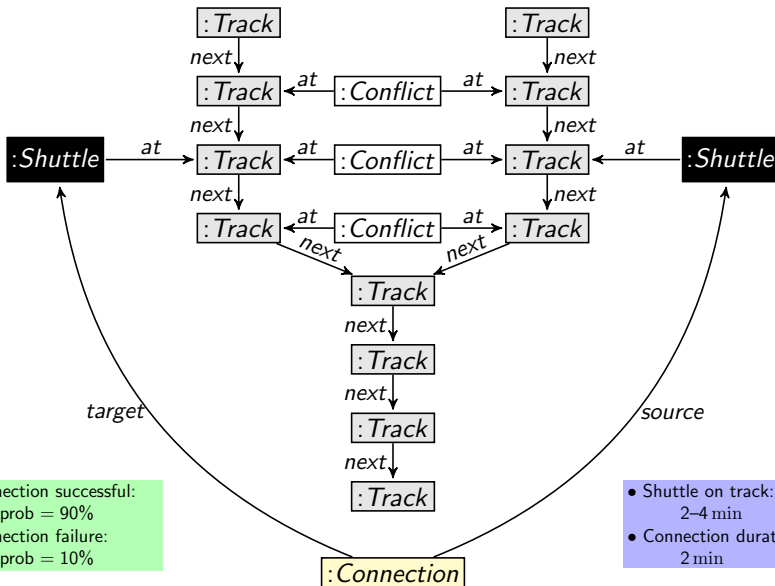
- Shuttle on track:  
2-4 min
- Connection duration:  
2 min





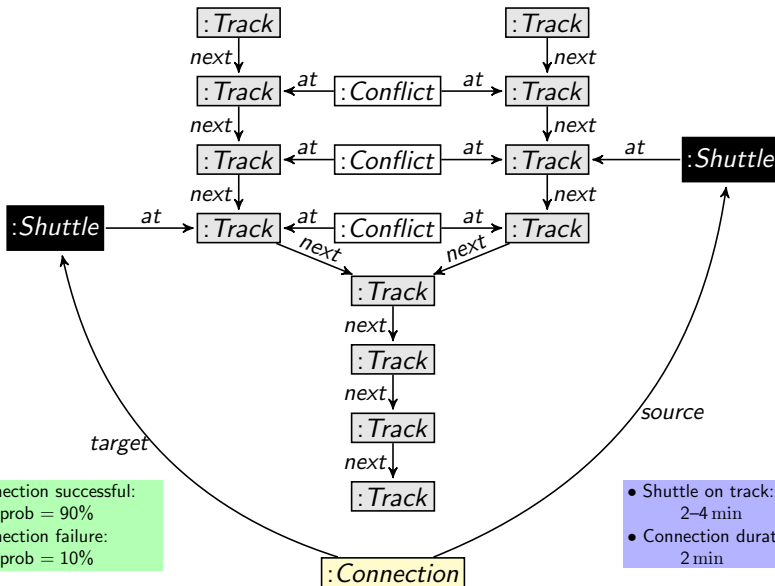
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2-4 min
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2 min



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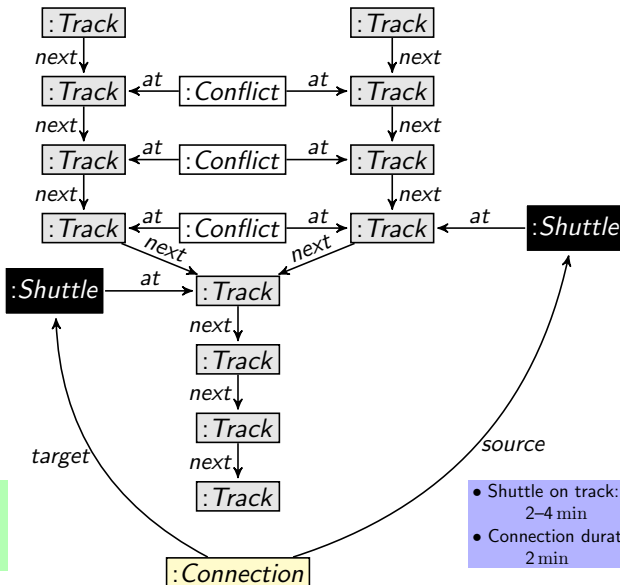
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- Connection failure:  
prob = 10%

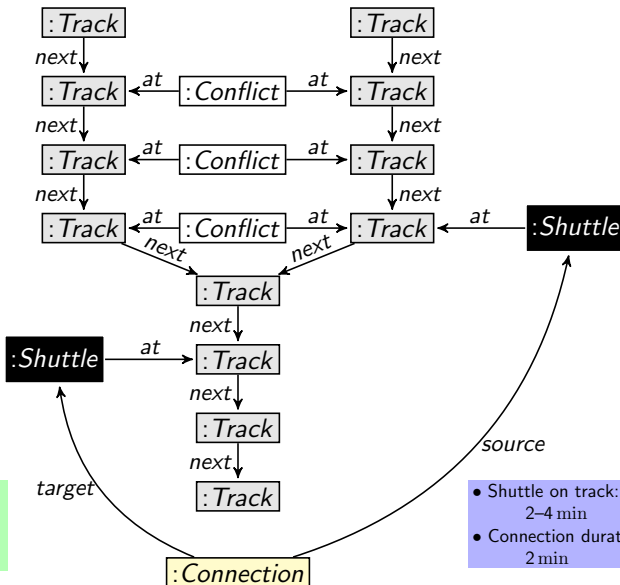
- Shuttle on track:  
2-4 min
- Connection duration:  
2 min





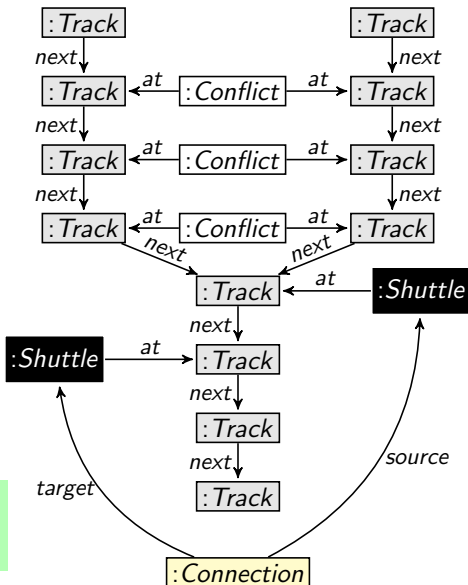
- Connection successful:  
prob = 90%
- Connection failure:  
prob = 10%

- Shuttle on track:  
2–4 min
- Connection duration:  
2 min



- Connection successful: prob = 90%
- Connection failure: prob = 10%

- Shuttle on track: 2-4 min
- Connection duration: 2 min



- Connection successful:  
prob = 90%
- Connection failure:  
prob = 10%

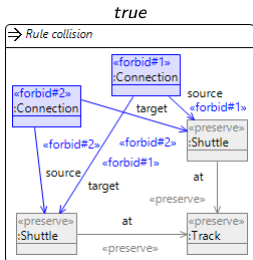
- Shuttle on track:  
2-4 min
- Connection duration:  
2 min





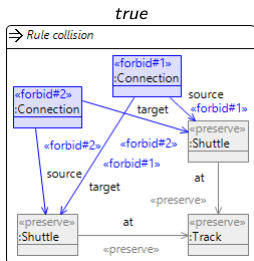
Can the considered shuttle scenario under any circumstances exhibit a collision?

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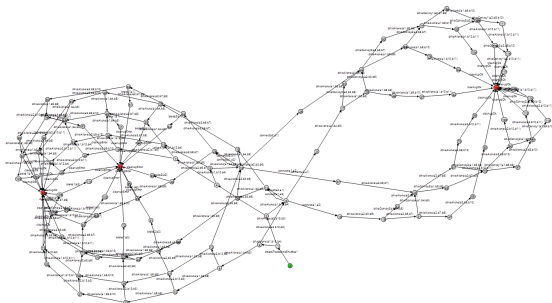


Atomic proposition *collision*

Can the considered shuttle scenario under any circumstances exhibit a collision?

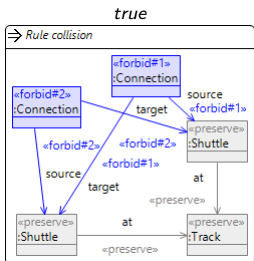


Atomic proposition *collision*

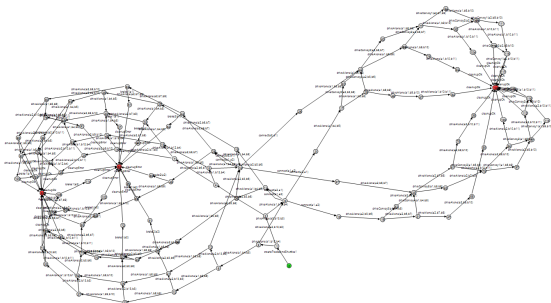


State space for the topology with 3 conflict nodes

Can the considered shuttle scenario under any circumstances exhibit a collision?



Atomic proposition *collision*



State space for the topology with 3 conflict nodes

**Result:** No collisions are possible.

What is the maximal probability that a shuttle executes an emergency brake?

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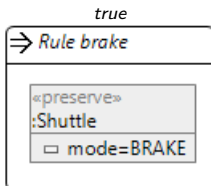
- State spaces for topologies with 2–6 conflict nodes generated by HENSHIN

What is the maximal probability that a shuttle executes an emergency brake?

- State spaces for topologies with 2–6 conflict nodes generated by HENSHIN
- State spaces exported to PRISM

What is the maximal probability that a shuttle executes an emergency brake?

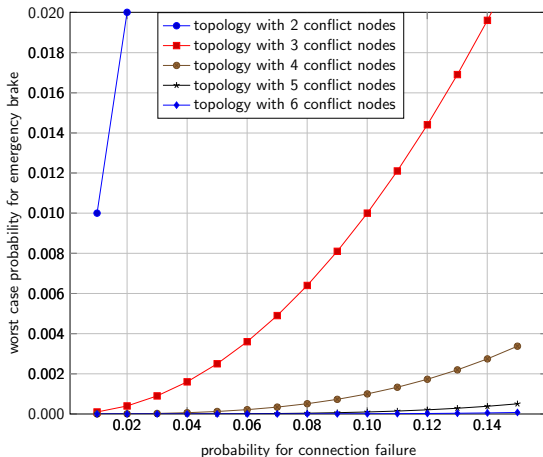
- State spaces for topologies with 2–6 conflict nodes generated by HENSHIN
- State spaces exported to PRISM
- Property for braking behavior in PRISM notation:  $Pmax = ? [F \text{ „brake“}]$

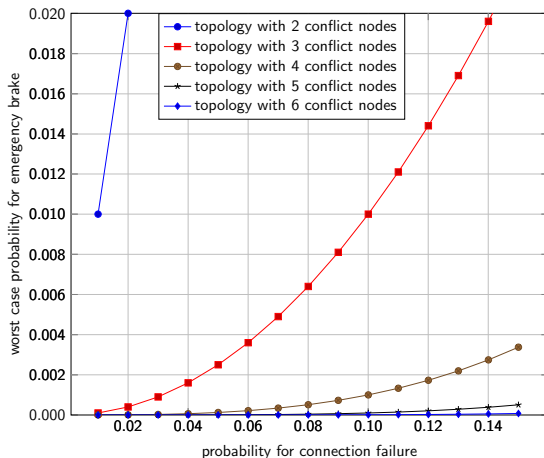


Atomic proposition *brake*



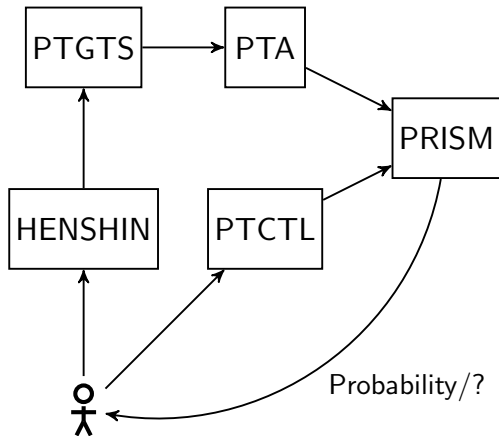
# Experiment 2 (Probability for Emergency Brake)





**Result:** For a desired worst case probability (y-axis) and a communication service quality (x-axis) we can determine the minimal required number of conflict nodes.

- ① Probabilistic Timed Graph Transformation Systems
- ② Modeling and Analysis of a Shuttle Example
- ③ Conclusion and Future Work



- Extension of PTGTSs to interval probabilities (and therefore to IPTA)
- Extension of PTCTL to path properties to specify structure dynamics
- Extension of PRISM to verify path properties of PTGTSs

Thank You For Your Attention