## Hasso

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## Similarity measures

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## Duplicate Detection - Research



## Overview Similarity Measures



## Similarity measures

- $\operatorname{sim}(x, y)$
$\square x$ and $y$ can be strings, numbers, tuples, objects, images, ...
- Normalized: $\operatorname{sim}(x, y) \in[0,1]$
$\square \operatorname{sim}(x, y)=1$ for exact match
$\square \operatorname{sim}(x, y)=0$ for „completely different" $x$ and $y$.
$\square 0<\operatorname{sim}(x, y)<1$ for some approximate similarity
- Distance function / distance metric
$\square$ Reflexive: $\operatorname{dist}(x, x)=0$
- Positive: $\quad \operatorname{dist}(x, y) \geq 0$
- Symmetric: $\operatorname{dist}(x, y)=\operatorname{dist}(y, x)$
$\square$ Triangular inequation: $\operatorname{dist}(x, z) \leq \operatorname{dist}(x, y)+\operatorname{dist}(y, z)$
- $\operatorname{sim}(x, y)=1-\operatorname{dist}(x, y)$
- $\operatorname{sim}(x, y)=1 / \operatorname{dist}(x, y)$


## Overview Similarity Measures



## Exact and truncated match

- $\operatorname{sim}_{\text {exact }}(x, y)=\left\{\begin{array}{l}1 \text { if } x=y \\ 0 \text { if } x \neq y\end{array}\right.$
- $\operatorname{sim}_{\text {trunc_beg }}(x, y)=\left\{\begin{array}{l}1 \text { if } x[1: k]=y[1: k] \\ 0 \text { if } x[1: k] \neq y[1: k]\end{array}\right.$
- sim $_{\text {trunc_end }}(x, y)=\left\{\begin{array}{l}1 \text { if } x[k: n]=y[k: n] \\ 0 \text { if } x[k: n] \neq y[k: n]\end{array}\right.$
- $\operatorname{sim}_{\text {encode }}(x, y)=\left\{\begin{array}{l}1 \text { if encode }(x)=\operatorname{encode}(y) \\ 0 \text { if encode }(y) \neq \operatorname{encode}(y)\end{array}\right.$
- E.g., with a phonetic encoding


## Hamming distance

■ Number of positions in which two strings (of equal length) differ
$\square$ Minimum number of substitutions required to change one string into the other
$\square$ Minimum number of errors that could have transformed one string into the other.

- Used mostly for binary numbers and to measure communication errors.
$\square$ Hamming distance $=$ number of 1 s in $\times$ XOR $y$.

■ dist $_{\text {hamming }}($ peter, pedro $)=3$

## Edit distances

- Compare two strings based on individual characters
- Minimal number of edits required to transform one string into the other.
$\square$ Edits: Insert, Delete, Replace (and Match)
$\square$ Alternative: Smallest edit cost
- Give different cost to different types of edits
$\square$ Give different cost to different letters
■ Naive approach: editdistance(J ones,J ohnson)
- DDDDDIIIIII = 12
$\square$ But: Not minimal!
- Levenshtein distance: Basic form
- Each edit has cost 1


## Levenshtein Distance

- Minimum number of character insertions, deletions, and replacements necessary to transform $s_{1}$ into $s_{2}$
- Compute transcript based on dynamic programming algorithm
- Optimality principle: Best transcript of two substrings must be part of best overall solution

1. Initialize matrix $M$ of size $\left(\left|s_{1}\right|+1\right) \times\left(\left|s_{2}\right|+1\right)$
2. Fill matrix: $M_{i, 0}=i$ and $M_{0, j}=j$
3. Recursion: $M_{i, j}=\left\{\begin{array}{cc}M_{i-1, j-1} & \text { if } x[i]=y[j] \\ 1+\min \left(M_{i-1, j}, M_{i, j-1}, M_{i-1, j-1}\right) & \text { else }\end{array}\right.$
4. Distance: LevenshteinDist $(x, y)=M_{|x|,|y|}$

Levenshtein Similarity: $\operatorname{sim}_{\text {Levenshtein }}(x, y)=1-\frac{\text { LevenshteinDist }(x, y)}{\max (|x|,|y|)}$

## Levenshtein Distance

|  |  | $\mathbf{J}$ | $\mathbf{O}$ | $\mathbf{N}$ | $\mathbf{E}$ | $\mathbf{S}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | 1 | 2 | 3 | 4 | 5 |
| $\mathbf{J}$ | 1 |  |  |  |  |  |
| $\mathbf{O}$ | 2 |  |  |  |  |  |
| $\mathbf{H}$ | 3 |  |  |  |  |  |
| $\mathbf{N}$ | 4 |  |  |  |  |  |
| $\mathbf{S}$ | 5 |  |  |  |  |  |
| $\mathbf{O}$ | 6 |  |  |  |  |  |
| $\mathbf{N}$ | 7 |  |  |  |  |  |

$$
M_{i, j}=\left\{\begin{array}{cc}
M_{i-1, j-1} & \text { if } x[i]=y[j] \\
1+\min \left(M_{i-1, j}, M_{i, j-1}, M_{i-1, j-1}\right) & \text { else }
\end{array}\right.
$$

## Levenshtein Distance - Example

| $\mathbf{s}_{\mathbf{1}}$ | $\mathbf{s}_{\mathbf{2}}$ | Levenshtein <br> Distance | sim $_{\text {Levenshtein }}$ |
| :--- | :--- | :---: | :---: |
| Jones | J ohnson | 4 | 0.43 |
| Paul | Pual | 2 | 0.5 |
| Paul Jones | Jones, Paul | 11 | 0 |

## Levenshtein discussion

- Complexity
$\square$ Time: $O(|x| \cdot|y|)$ (fill in matrix)
$\square$ Space: $O(\min (|x|,|y|))$
$\diamond$ Trick: Store only two rows of the matrix
- Some properties
- $0 \leq$ LevenshteinDist $(x, y) \leq \max (|x|,|y|)$
$\square||x|-|y|| \leq$ LevenshteinDist $(x, y)$
$\diamond$ Often: Compare only strings with similar lengths
- Other cost models
$\square$ Insert, delete cost 1.0 but replace 0.5
$\diamond$ change in string length is punished, e.g. for zip codes
- Character based: $O C R(m \simeq n, 1 \simeq I)$ or keyboard ( $a \simeq s$ ) or brain $(6 \simeq 9)$ or biology $(a \simeq t)$


## Damerau-Levenshtein distance

- Similar to Levenshtein distance, but additionally considers transposed characters
- $M_{i, 0}=i$ and $M_{0, j}=j$
- $M_{i, j}=$

$$
\left\{\begin{array}{cc}
M_{i-1, j-1} \\
M_{i-1, j}, M_{i, j-1}, & i f x[i]=y[j] \\
M_{i-1, j-1}, \\
1+\min \binom{\text { in }}{M_{i-2, j-2} \text { if } x[i]=y[j-1] \text { and } x[i-1]=y[j]} & \text { else }
\end{array}\right.
$$

| $\mathbf{s}_{\mathbf{1}}$ | $\mathbf{s}_{\mathbf{2}}$ | Levenshtein <br> Distance | Damerau-Levenshtein <br> Distance | sim <br> Damerau- <br> Levenshtein |
| :--- | :--- | :---: | :---: | :---: |
| Jones | Johnson | 4 | 4 | 0.43 |
| Paul | Pual | 2 | 1 | 0.75 |
| Paul Jones | Jones, Paul | 11 | 11 | 0 |

## Jaro similarity

- Specifically designed for names at US Census Bureau
- Search for common characters
- m : number of matching characters
- Search range matching characters: $\frac{\max (|x|,|y|)}{2}-1$
- t : number of transpositions
- $\operatorname{sim}_{\text {jaro }}=\frac{1}{3}\left(\frac{m}{|x|}+\frac{m}{|y|}+\frac{m-t}{m}\right)$


## Jaro similarity - Example

■ $\operatorname{sim}_{\text {jaro }}=\frac{1}{3}\left(\frac{m}{|x|}+\frac{m}{|y|}+\frac{m-t}{m}\right)$

$$
\begin{array}{lll}
\mathrm{S}_{1} & \mathbf{P} & \mathbf{A} \\
& \downarrow & \mathbf{U} \\
\mathrm{~S}_{2} & \mathbf{P} & \mathbf{U} \mathbf{A} \\
m=4 & t=\frac{2}{2}=1 \\
\operatorname{sim}_{\text {jaro }}= & \frac{1}{3} \cdot\left(\frac{4}{4}+\frac{4}{4}+\frac{4-1}{4}\right) \approx 0.92
\end{array}
$$



$$
m=4 \quad t=\frac{0}{2}=0
$$

$m=4 \quad t=\frac{0}{2}=0$
$\operatorname{sim}_{\text {jaro }}=\frac{1}{3} \cdot\left(\frac{4}{5}+\frac{4}{7}+\frac{4-0}{4}\right) \approx 0.79$

## Winkler similarity

- Intuition 1: Similarity of first few letters is most important.
- Let p be the length of the common prefix of x and y .
- $\operatorname{sim}_{\text {winkler }}(x, y)=\operatorname{sim}_{\text {jaro }}(x, y)+\left(1-\operatorname{sim}_{\text {jaro }}(x, y)\right) \frac{p}{10}$
$\square=1$ if common prefix is $\geq 10$
- Intuition 2: Longer strings with even more common letters
- $\operatorname{sim}_{\text {winkler_long }}(x, y)=\operatorname{sim}_{\text {winkler }}(x, y)+\left(1-\operatorname{sim}_{\text {winkler }}(x, y)\right) \frac{c-(p+1)}{|x|+|y|-2(p-1)}$
$\square$ Where c is overall number of common letters
$\square$ Apply only if
$\diamond$ Long strings: $\min (|x|,|y|) \geq 5$
$\diamond$ Two additional common letters: $c-p \geq 2$
$\diamond$ At least half remaining letters of shorter string are in common: $c-p \geq \frac{\min (|x|,|y|)-p}{2}$

| String 1 | String 2 | $c$ | $f$ | $p$ | $c_{\text {sim }}$ | $\operatorname{sim}_{\text {jam }}$ | $\operatorname{sim}_{\text {binkler }}$ | $\operatorname{sim}_{\text {winkler_loug }}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| shackleford | shackelford | 11 | 1 | 4 | 0 | 0.9697 | 0.9818 | 0.9886 |
| nichleson | nichulson | 8 | 0 | 4 | 0.3 | 0.9259 | 0.9556 | 0.9667 |
| jones | johnson | 4 | 0 | 2 | 0.3 | 0.7905 | 0.8324 | 0.8491 |
| massey | massie | 5 | 0 | 4 | 0.3 | 0.8889 | 0.9333 | - |
| jeraldine | geraldine | 8 | 0 | 0 | 0.3 | 0.9259 | 0.9259 | 0.9519 |
| michelle | michael | 6 | 0 | 4 | 0.3 | 0.8690 | 0.9214 | 0.9302 |



- Forming words from sequence of characters
- General idea: Separate string into tokens using some separator
$\square$ Space, hyphen, punctuation, special characters
- Usually also convert to lower-case
- Problems
$\square$ Both hyphenated and non-hyphenated forms of many words are common
$\diamond$ Sometimes hyphen is not needed
- e-bay, wal-mart, active-x, cd-rom, t-shirts
$\diamond$ Sometimes hyphens should be considered either as part of the word or a word separator
- winston-salem, mazda rx-7, e-cards, pre-diabetes, t-mobile, spanishspeaking
- Apostrophes can be a part of a word, a part of a possessive, or just a mistake
$\diamond$ rosie o'donnell, can't, don't, 80's, 1890's, men's straw hats, master's degree, england's ten largest cities, shriner's
- Numbers can be important, including decimals
$\diamond$ nokia 3250, top 10 courses, united 93 , quicktime 6.5 pro, 92.3 the beat, 288358
- Periods can occur in numbers, abbreviations, URLs, ends of sentences, and other situations
$\diamond ~ I . B . M ., ~ P h . D ., ~ c s . u m a s s . e d u, ~ F . E . A . R$.

Die Kapostropheum-Gruselgalerie Kategorie „Völlig willenlos"
http://www.apostroph.de/


## n-grams (aka q-grams)

- Split string into short substrings of length n.
$\square$ Sliding window over string
- $\mathrm{n}=2$ : Bigrams
- $\mathrm{n}=3$ : Trigrams
- Variation: Pad with $\mathrm{n}-1$ special characters
$\diamond$ Emphasizes beginning and end of string
$\square$ Variation: Include positional information to weight similarities
- Number of $n$-grams $=|x|-n+1$
- Count how many n -grams are common in both strings

| String | Bigrams | Padded bigrams | Positional bigrams | Trigrams |
| :---: | :---: | :---: | :---: | :---: |
| gail | ga, ai, il | $\odot \mathrm{g}, \mathrm{ga}, \mathrm{ai}, \mathrm{il}, \mathrm{l} \otimes$ | $(\mathrm{ga}, 1),(\mathrm{ai}, 2),(\mathrm{il}, 3)$ | gai, ail |
| gayle | ga, ay, yl, le | $\odot \mathrm{g}, \mathrm{ga}, \mathrm{ay}, \mathrm{yl}, \mathrm{le}, \mathrm{e} \otimes$ | $(\mathrm{ga}, 1),(\mathrm{ay}, 2),(\mathrm{yl}, 3),(\mathrm{le}, 4)$ | gay, ayl, yle |
| peter | pe, et, te, er | $\odot \mathrm{p}$, pe, et, te, er, r $\otimes$ | $(\mathrm{pe}, 1),(\mathrm{et}, 2),(\mathrm{te}, 3),(\mathrm{er}, 4)$ | pet, ete, ter |
| pedro | pe, ed, dr, ro | $\odot \mathrm{p}, \mathrm{pe}, \mathrm{ed}, \mathrm{dr}, \mathrm{ro}, \mathrm{o} \otimes$ | $(\mathrm{pe}, 1),(\mathrm{ed}, 2),(\mathrm{dr}, 3),(\mathrm{ro}, 4)$ | ped, edr, dro |

## Token-based Similarity Measures

- Token similarity
- Overlap coefficient: $\operatorname{sim}_{\text {overlap }}(x, y)=\frac{|\operatorname{tok}(x) \cap \operatorname{tok}(y)|}{\min (|\operatorname{tok}(x)|,|\operatorname{tok}(y)|)}$
- Jaccard coefficient:

$$
\operatorname{sim}_{j \operatorname{jaccard}}(x, y)=\frac{|\operatorname{tok}(x) \cap \operatorname{tok}(y)|}{|\operatorname{tok}(x)|+|\operatorname{tok}(y)|-|\operatorname{tok}(x) \cap \operatorname{tok}(y)|}=\frac{|\operatorname{tok}(x) \cap \operatorname{tok}(y)|}{|\operatorname{tok}(x) \cup \operatorname{tok}(y)|}
$$

- Dice's coefficient: $\operatorname{sim}_{\text {dice }}(x, y)=\frac{2 \cdot|\operatorname{tok}(x) \cap \operatorname{tok}(y)|}{|\operatorname{tok}(x)|+|\operatorname{tok}(y)|}$
- Tokens („Paul Jones")
- Words / Terms („Paul" „J ones")
- Padded n-grams (_P, Pa, au, ul, I_, _J, Jo, on, ne, es, s_)

| $\mathbf{s}_{\mathbf{1}}$ | $\mathbf{s}_{\mathbf{2}}$ | Jaccard | Dice |
| :--- | :--- | :---: | :---: |
| Jones | Johnson | 0.17 | 0.29 |
| Paul | Pual | 0.33 | 0.40 |
| Paul Jones |  | Jones, Paul | 0.77 |



| Digit | Letters |
| :---: | :--- |
| 1 | B, F, P, V |
| 2 | C, G, J, K, Q, S, X, Z |
| 3 | D, T |
| 4 | L |
| 5 | M, N |
| 6 | R |

3. Two adjacent letters with the same number are coded as a single number
4. Continue until you have one letter and three numbers. If you run out of letters, pad with 0s.

- If a surname has a prefix, such as Van, Con, De, Di, La, or Le, code both with and without the prefix
- Example
- PAUL: P400

■ PUAL: P400

- JONES: J520
- JOHNSON: J525
- Soundex codes a last name based on the way a last name sounds

1. Retain first letter of the name and drop all other occurrences of A, E, H, I, O, U, W, Y
2. Replace consonants with digits

## Soundex on WolframAlpha

## 参WolframAlphai menatas <br> WolframAlphà enamas

Soundex Levenshtein

Input interpretation:
Soundex Levenshtein

Soundex code:
L152

Soundex-close English words:
More
Livingstone | lebensraum | Livingston | lovemaking
Computed by Wolfram 2Nobematica
Download page

| Letter | Context | Code |
| :---: | :---: | :---: |
| A, E, I, J, O, U, Y |  | 0 |
| H |  | - |
| B |  | 1 |
| $P$ | not before H |  |
| D, T | not before C, S, Z | 2 |
| F, V, W |  | 3 |
| P | before H |  |
| G, K, Q |  |  |
| C | in the initial sound before A, H, K, L, O, Q, R, U, X | 4 |
|  | before A, H, K, O, Q, U, X but not after $\mathrm{S}, \mathrm{Z}$ |  |
| X | not after C, K, Q | 48 |
| L |  | 5 |
| M, N |  | 6 |
| R |  | 7 |
| S, Z |  |  |
|  | after S, Z |  |
| C | in the initial sound, but not before A, H, K, L, O, Q, R, U, X | 8 |
|  | not before $A, H, K, O, Q, U, X$ |  |
| D, T | before $\mathrm{C}, \mathrm{S}, \mathrm{Z}$ |  |
| X | after C, K, Q |  |

## Felix Naumann | Data Profiling and Data Cleansing | Summer 2013

## Metaphone

- Improves on the Soundex algorithm
$\square$ Knows variations and inconsistencies in English spelling and pronunciation
- Further improvements
- Double Metaphone
$\diamond$ Includes other languages: Slavic, Germanic, Celtic, Greek, French, Italian, Spanish, Chinese
$\diamond$ Accuracy 89\%
- Metaphone 3
$\diamond$ Accuracy over 99\% (says author)


## Original Metaphone Algorithm

- 16 consonant symbols OBFHJ KLMNPRSTWXY
- '0' represents "th", 'X' represents "sh" or "ch"

1. Drop duplicate adjacent letters, except for C .
2. If the word begins with 'KN', 'GN', 'PN', 'AE', 'WR', drop the first letter.
3. Drop ' B ' if after ' M ' at the end of the word.
4. ' C ' transforms to ' X ' if followed by 'IA' or ' $\mathrm{H}^{\prime}$ (unless in latter case, it is part of '- SCH - ', in which case it transforms to ' $K$ '). 'C' transforms to 'S' if followed by 'I', ' $E$ ', or ' $Y$ '. Otherwise, ' $C^{\prime}$ transforms to ' $K$ '.
5. 'D' transforms to 'J' if followed by 'GE', 'GY', or 'GI'. Otherwise, 'D' transforms to 'T'.
6. Drop ' $G$ ' if followed by ' H ' and ' H ' is not at the end or before a vowel. Drop ' G ' if followed by ' N ' or 'NED' and is at the end.
7. 'G' transforms to 'J' if before 'I', ' $E$ ', or ' $Y$ ', and it is not in ' $G G^{\prime}$. Otherwise, ' $G$ ' transforms to ' $K$ '.
8. Drop ' H ' if after vowel and not before a vowel.
9. 'CK' transforms to ' K '.
10. 'PH' transforms to ' F '.
11. 'Q' transforms to ' K '.
12. 'S' transforms to ' X ' if followed by ' H ', 'IO', or 'IA'.
13. ' $T$ ' transforms to ' X ' if followed by ' IA ' or ' IO '. 'TH' transforms to ' 0 '. Drop ' $T$ ' if followed by ' CH '.
14. 'V' transforms to ' F '.
15. 'WH' transforms to 'W' if at the beginning. Drop 'W' if not followed by a vowel.
16. ' $X$ ' transforms to ' S ' if at the beginning. Otherwise, ' X ' transforms to 'KS'.
17. Drop ' Y ' if not followed by a vowel.
18. 'Z' transforms to 'S'.
19. Drop all vowels unless it is the beginning


## Monge-Elkan

■ Hybrid: Token-based and internal similarity function for tokens
$\square$ Find best match for each token

- $\operatorname{sim}_{\text {MongeElkan }}(x, y)=\frac{1}{|x|} \sum_{i=1}^{|x|} \max _{j=1,|y|} \operatorname{sim}^{\prime}(x[i], y[j])$
$\square|x|$ is number of tokens in $x$
$\square$ sim' is internal similarity function (e.g., Levenshtein)
- If strings contain just one token each
$\square \operatorname{sim}_{\text {MongeElkan }}(x, y)=\operatorname{sim}^{\prime}(x, y)$

■ Complexity: Quadratic in number of tokens

## Monge-Elkan - Example

- $\operatorname{sim}_{\text {MongeElkan }}(x, y)=\frac{1}{|x|} \sum_{i=1}^{|x|} \max _{j=1,|y|} \operatorname{sim}^{\prime}(x[i], y[j])$
- Peter Christen vs. Christian Pedro
$\square \operatorname{sim}_{\text {jaro }}$ (peter, christian) $=0.3741$
$\square \operatorname{sim}_{\text {jaro }}$ (peter, pedro) $=0.7333$
$\square \operatorname{sim}_{\text {jaro }}($ christen, christian) $=0.8843$
$\square \operatorname{sim}_{\text {jaro }}($ christen, pedro $)=0.4417$
- $\operatorname{sim}_{\text {MongeElkan }}($ 'peter christen',' christian pedro' $)=\frac{1}{2}(0.7333+0.8843)=$ 0.8088

■ $\operatorname{sim}_{\text {MongeElkan }}(x, y) \neq \operatorname{sim}_{\text {MongeElkan }}(y, x)$


[^0]
## Extended J accard Similarity

- $\operatorname{sim}_{j \operatorname{accard}}(x, y)=\frac{|\operatorname{tok}(x) \cap \operatorname{tok}(y)|}{|\operatorname{tok}(x)|+|\operatorname{tok}(y)|-|\operatorname{tok}(x) \cap \operatorname{tok}(y)|}=\frac{|\operatorname{tok}(x) \cap \operatorname{tok}(y)|}{|\operatorname{tok}(x) \cup \operatorname{tok}(y)|}$
- If strings contain multiple words, choose words as tokens.
- Use internal similarity function to calculate similarity between all pairs of tokens.
■ Shared tokens:

$$
S=\left\{\left(x_{i}, y_{j}\right) \mid x_{i} \in \operatorname{tok}(x) \wedge y_{j} \in \operatorname{tok}(y): \operatorname{sim}^{\prime}\left(x_{i}, y_{j}\right) \geq \theta\right\}
$$

$\square$ Unique tokens: $U_{\operatorname{tok}(x)}=\left\{x_{i} \mid x_{i} \in \operatorname{tok}(x) \wedge y_{j} \in \operatorname{tok}(y) \wedge\left(x_{i}, y_{j}\right) \notin S\right\}$

- $\operatorname{sim}_{j_{\text {accad_ext }}}(x, y)=\frac{|S|}{|S|+\left|U_{\text {tok }(x)}\right|+\left|U_{\text {tok }(x)}\right|}$


## Vector Space Model

- Each document ranked by distance between points representing query and document
- Popular measure: Cosine similarity
$\square$ Cosine of angle between document and query vectors
- Normalized dot-product

$$
\operatorname{Cosine}\left(D_{i}, Q\right)=\frac{\sum_{j=1}^{t} d_{i j} \cdot q_{j}}{\sqrt{\sum_{j=1}^{t} d_{i j}{ }^{2} \cdot \sum_{j=1}^{t} q_{j}{ }^{2}}}
$$


http://www.euclideanspace.com/math s/geometry/trig/derived/index.htm

## Similarity Calculation - Example

- Consider three documents $D_{1}, D_{2}, D_{3}$ and query $Q$
$\square D_{1}=(0.5,0.8,0.3), D_{2}=(0.9,0.4,0.2), D_{3}=(0,0.9,0.1)$
$\square \mathrm{Q}=(1.5,1.0,0)$
- Vector space model reflects some term weights and number of matching terms (in contrast to Boolean retrieval)

$$
\begin{aligned}
\operatorname{Cosine}\left(D_{1}, Q\right) & =\frac{(0.5 \times 1.5)+(0.8 \times 1.0)}{\sqrt{\left(0.5^{2}+0.8^{2}+0.3^{2}\right)\left(1.5^{2}+1.0^{2}\right)}} \\
& =\frac{1.55}{\sqrt{(0.98 \times 3.25)}}=0.87 \\
\operatorname{Cosine}\left(D_{2}, Q\right) & =\frac{(0.9 \times 1.5)+(0.4 \times 1.0)}{\sqrt{\left(0.9^{2}+0.4^{2}+0.2^{2}\right)\left(1.5^{2}+1.0^{2}\right)}} \\
& =\frac{1.75}{\sqrt{(1.01 \times 3.25)}}=0.97 \quad \operatorname{Cosine}\left(D_{3}, Q\right)=0.55
\end{aligned}
$$

■ But: How to assign term weights?

## Term Weights - tf.idf

■ Term frequency weight tf measures importance of term $k$ in document i: $t f_{i k}=\frac{f_{i k}}{\sum_{j=1}^{t} f_{i j}}$
$\square \log \left(f_{i k}\right)$ to reduce this impact of frequent words

■ Inverse document frequency idf measures importance in collection: $i d f_{k}=\log \frac{N}{n_{k}}$

- Reflects "amount of information" carried by term

■ tfidf by multiplying tf and idf with some heuristic modifications

## SoftTFIDF

- Apply idea to records or values: Much shorter than documents
- $\operatorname{CLOSE}(\theta, \operatorname{tok}(x), \operatorname{tok}(y))$ is set of tokens from x that have at least one sufficiently similar token in y .
- $\operatorname{sim}_{\text {softtfidf }}(x, y)=$
$\sum_{t \in \operatorname{CLOSE}(\theta, \operatorname{tok}(x), \operatorname{tok}(y))} V(t, \operatorname{tok}(x)) \cdot V(t, \operatorname{tok}(y)) \cdot N(t, \operatorname{tok}(y))$
$\square V(t, \operatorname{tok}(x))$ is TFIDF weight of token t in all tokens of x
$\square N(t, \operatorname{tok}(y))=\max \left(\left\{\operatorname{sim}^{\prime}\left(t, y_{j}\right) \mid y_{j} \in \operatorname{tok}(y)\right\}\right)$
$\diamond$ Similarity of best matching token
- Soft: Tokens are considered a partial match if they get a good score using an internal similarity measure (CLOSE).
- Problem: Weights are calculated over entire database
- Scan all data
- Store weight for each unique token



## Numerical comparison

- $\operatorname{sim}_{\text {num_abs }}(n, m)=\left\{\begin{array}{cc}1-\left(\frac{|n-m|}{d_{\max }}\right) & \text { if }|n-m|<d_{\text {max }} \\ 0 & \text { else }\end{array}\right.$
- Linear extrapolation between 0 and $d_{\max }$
- Example:
- $\mathrm{d}_{\max }=\$ 1,000$
$\square \quad \operatorname{sim}_{\text {num_abs }}(2,000,2,500)=1-\frac{500}{1,000}=0.5$
$\square \operatorname{sim}_{\text {num_abs }}(200,000,200,500)=1-\frac{500}{1,000}=0.5$
- $\operatorname{sim}_{\text {num_perc }}(n, m)=\left\{\begin{array}{cc}1-\left(\frac{p c}{p c_{\max }}\right) & \text { if } p c<p c_{\max } \\ 0 & \text { else }\end{array}\right.$
$\square p c=\frac{|n-m|}{\max (|n|,|m|)} \cdot 100$ is percentage difference
$\square \mathrm{pc}_{\max }=33 \%$
$\square \operatorname{sim}_{\text {num_perc }}(2,000,2,500)=1-\frac{20}{33}=0.394$ because $p c=\frac{|2,000-2,500|}{2,500} \cdot 100=20$
$\square \operatorname{sim}_{\text {num_perc }}(200,000,200,500)=1-\frac{0,25}{33}=0.993$ because $p c=\frac{500}{200,500} \cdot 100=0,25 \%$


## Time and space comparisons

- Calculate difference in days and use sim $_{\text {num_abs }}$
- Special cases
$\square$ Swapped day and month and both $\leq 12$ : Return some fixed similarity, e.g. 0.5
$\square$ Single error in month could exceed $\mathrm{d}_{\text {max }}$ : Return some fixed similarity, e.g., 0.75
- Dates of birth can be converted to age (num days or num years)
$\square$ Then apply numerical measures
$\square \operatorname{sim}_{\text {age_perc }}(n, m)=\left\{\begin{array}{cc}1-\left(\frac{a p c}{a p c_{\max }}\right) & \text { if apc }<a p c_{\text {max }} \\ 0 & \text { else }\end{array}\right.$
$\square a p c=\frac{|n-m|}{\max (|n|,|m|)} \cdot 100$ is percentage difference
■ Geographical data: Compute distance based on some projection



## Similarity function packages

- SecondString
- Java classes
$\square$ All basic string comparisons
- MongeElkan, SoftTFIDF
- Similarity learner
- http://sourceforge. net/projects/secondstring/
- SimMetrics
$\square$ Java package
- All basic string comparisons
$\square$ Long sequences: Needleman-Wunsch, Smith-Waterman, Smith-Waterman-Gotoh
- http://sourceforge.net/projects/simmetrics/
- Geographiclib for geographic similarity
- http://geographiclib.sourceforge.net


[^0]:    Felix Naumann | Data Profiling and Data Cleansing | Summer 2013

