

Heuristic Optimization – Homework 2

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2 Exercise 2

2.1 Given

$$\forall x \in \mathbb{R}: e^x = \sum_{i=0}^{\infty} \frac{x^i}{i!} = 1 + x + \frac{x^2}{2} + \frac{x^3}{6} + \dots$$

2.2 Task a

To show:

$$\forall x \geq -1, k \in \mathbb{N}: (1+x)^k \geq 1+kx$$

Base case: $k=0$

$$(1+x)^0 \geq 1 \geq 1+0 \cdot x$$

True.

Induction hypothesis:

$(1+x)^k \geq 1+k \cdot x$ is correct for some natural number k .

Induction-step:

We now have to prove, that $(1+x)^{k+1} \geq 1+(k+1) \cdot x$.

$$(1+x)^k \geq 1+k \cdot x \quad (\text{induction hypothesis})$$

$$\Rightarrow (1+x)^k \cdot (1+x) \geq (1+x) + kx \cdot (1+x)$$

$$\Rightarrow (1+x)^{k+1} \geq 1+x+kx+kx^2$$

Since kx^2 is greater 0, we can make the right side even smaller by subtracting it.

$$\Rightarrow (1+x)^{k+1} \geq 1+kx+x$$

$$\Rightarrow (1+x)^{k+1} \geq 1+(k+1)x$$

Hence by mathematical induction the inequality is correct for all natural numbers.

q.e.d.

2.3 Task b

To show:

$$\forall x > -1, x \neq 0: \ln(1+x) < x$$

Proof:

$$\boxed{\ln(1+x) < x}$$

$$\Leftrightarrow 1+x < e^x$$

$$\Leftrightarrow 1+x < \sum_{i=0}^{\infty} \frac{x^i}{i!} \quad \Leftrightarrow 1+x < 1+x + \sum_{i=2}^{\infty} \frac{x^i}{i!}$$

$$\Leftrightarrow 0 < \sum_{i=2}^{\infty} \frac{x^i}{i!} = \sum_{i=1}^{\infty} \frac{x^{2i}}{(2i)!} + \frac{x^{2i+1}}{(2i+1)!}$$

If both summands together are greater zero, the sum also will be greater than zero.

$$\Leftrightarrow 0 < \frac{x^{2i}}{(2i)!} + \frac{x^{2i+1}}{(2i+1)!} = \frac{x^{2i} \cdot (2i+1)}{(2i)! \cdot (2i+1)} + \frac{x^{2i+1}}{(2i+1)!} = \frac{x^{2i} \cdot (2i+1)}{(2i+1)!} + x^{2i} \cdot \frac{x}{(2i+1)!}$$

$$\Leftrightarrow 0 < \frac{x^{2i} \cdot (2i+1+x)}{(2i+1)!}$$

$$\Leftrightarrow 0 < x^{2i} \cdot (2i+1+x)$$

$$\Leftrightarrow 0 < 2i+1+x, \quad x^{2i} > 0$$

$$\Leftrightarrow 0 < 1+x, \quad 2i > 0$$

$$\Leftrightarrow -1 < x$$

This is true, since it is given in the task.

Now, that the last statement is true, you can read the proof from bottom up, and you will get $\ln(1+x) < x$ in the end.

2.4 Task c

To show:

$$\forall x > -1, x \neq 0 \forall r > 0: (1+x)^r < e^{rx}$$

Proof:

$$\boxed{(1+x)^r < e^{rx}}$$

$$\Leftrightarrow 1+x < e^x$$

$$\Leftrightarrow \ln(1+x) < x$$

Already proven in 2.3(task b)

q.e.d.

3 Exercise 3

3.1 Given

The probability that a certain bit flips is $P(\text{a certain bit flips}) = \frac{1}{n}$ and the number of current zeros is k . $k > 0$, since all ones would already be the optimum.

3.2 Task a

To show:

$$P(\text{create all 1s individually}) = \Theta(n^{-k})$$

Proof:

$$P(\text{create all 1s individually}) = P(K)$$

$$P(K) = P(\text{all zeros flip to ones}) \cdot P(\text{all ones stay})$$

$$P(K) = P(\text{a certain zero flips to one})^k \cdot P(\text{a certain one stays})^{n-k}$$

$$P(K) = P(\text{a certain zero flips to one})^k \cdot (1 - P(\text{a certain one flips to zero}))^{n-k}$$

$$P(K) = \left(\frac{1}{n}\right)^k \cdot \left(1 - \frac{1}{n}\right)^{n-k} = n^{-k} \cdot \left(1 - \frac{1}{n}\right)^{n-k}$$

For $k > 0$, what is given, $\left(1 - \frac{1}{n}\right)^{n-k}$ goes towards e^{-1} from above. *

$$\Rightarrow \left(1 - \frac{1}{n}\right)^{n-k} > e^{-1}$$

Since $0 \leq 1 - \frac{1}{n} \leq 1$, and $n - k \geq 1$, $\Rightarrow \left(1 - \frac{1}{n}\right)^{n-k} \leq 1$

$$\Rightarrow \exists c_1, c_2, n_0 > 0 \forall n > n_0: c_1 \leq \left(1 - \frac{1}{n}\right)^{n-k} \leq c_2, \quad (c_1 = e^{-1} \wedge c_2 = 1)$$

Knowing, that n^{-k} is always positive, we can multiply it

$$\Leftrightarrow \exists c_1, c_2, n_0 > 0 \forall n > n_0: n^{-k} \cdot c_1 \leq n^{-k} \cdot \left(1 - \frac{1}{n}\right)^{n-k} \leq n^{-k} \cdot c_2$$

$$\Leftrightarrow \exists c_1, c_2, n_0 > 0 \forall n > n_0: n^{-k} \cdot c_1 \leq P(K) \leq n^{-k} \cdot c_2$$

By definition of Θ :

$$\boxed{\Leftrightarrow P(K) = \Theta(n^{-k})}$$

q.e.d.

3.3 Task b

To show:

$$\exists c > 0 : \forall n : P(\text{k is bigger after the next step}) \geq c \cdot \frac{k}{n}$$

Proof:

$$P(\text{k is bigger after the next step}) = P(B)$$

$$P(B) = P(\text{exactly one zero flips to one}) + P(\text{exactly two zeros flip to ones}) + \dots$$

Since there is no negative probability:

$$\Rightarrow P(B) \geq P(\text{exactly one zero flips to one})$$

$$\Rightarrow P(B) \geq P(\text{exact the first zero flips to one or the second zero flips to one or ...})$$

Since all these events are disjoint, the probability of their sum is just the sum of their probabilities.

$$\Rightarrow P(B) \geq P(\text{exact the first zero flips}) + P(\text{exact the second zero flips}) + \dots$$

All these events have the same probability

$$\Rightarrow P(B) \geq k \cdot P(\text{exact a certain zero flips})$$

$$\Rightarrow P(B) \geq k \cdot P(\text{a certain zero flips}) \cdot P(\text{all other bits stay})$$

$$\Rightarrow P(B) \geq k \cdot \frac{1}{n} \cdot \left(1 - \frac{1}{n}\right)^{n-1}$$

$\left(1 - \frac{1}{n}\right)^{n-1}$ goes towards e^{-1} from above *. Therefore: $\left(1 - \frac{1}{n}\right)^{n-1} > e^{-1}$

$$\Rightarrow P(B) \geq \frac{k}{n} \cdot \left(1 - \frac{1}{n}\right)^{n-1} > \frac{k}{n} \cdot e^{-1} = e^{-1} \cdot \frac{k}{n}$$

$$\Rightarrow P(B) \geq e^{-1} \cdot \frac{k}{n}$$

$$\boxed{\Rightarrow \exists c > 0 : \forall n : P(B) \geq c \cdot \frac{k}{n}}$$

q.e.d.

* Proof that for $n > k, k > 0$ $f(n) = \left(1 - \frac{1}{n}\right)^{n-k}$ goes towards e^{-1} from above.

First we want to show, that $f(n)$ goes towards e^{-1} :

$$\lim_{n \rightarrow \infty} f(n) = \lim_{n \rightarrow \infty} \left(1 - \frac{1}{n}\right)^{n-k} = \lim_{n \rightarrow \infty} \frac{\left(1 - \frac{1}{n}\right)^n}{\left(1 - \frac{1}{n}\right)^k} = \frac{e^{-1}}{1^k} = e^{-1}$$

To prove, that all values $n > k$ are greater than e^{-1} , it is enough to show, that the function is monotonically decreasing. Since f goes towards e^{-1} , they have to be greater than e^{-1} .

$f(n)$ Monotonically decreasing

$$\Leftrightarrow \forall n > k: f(n+1) \leq f(n) \text{ , (only for } f(n) \geq 0 \text{)}$$

$$\Leftrightarrow \forall n > k: \frac{f(n+1)}{f(n)} \leq 1$$

$$\Leftrightarrow \forall n > k: \frac{\left(1 - \frac{1}{n+1}\right)^{n-k+1}}{\left(1 - \frac{1}{n}\right)^{n-k}} = \frac{\left(\frac{n}{n+1}\right)^{n-k+1}}{\left(\frac{n-1}{n}\right)^{n-k}} = \frac{\left(\frac{n}{n+1}\right)^{n-k}}{\left(\frac{n-1}{n}\right)^{n-k}} \cdot \left(\frac{n}{n+1}\right) \leq 1$$

$$\Leftrightarrow \forall n > k: \left(\frac{n}{n+1}\right)^{n-k} \cdot \left(\frac{n}{n-1}\right)^{n-k} \cdot \frac{n}{n+1} = \left(\frac{n^2}{n^2-1}\right)^{n-k} \cdot \frac{n}{n+1} \leq 1$$

$$\Leftrightarrow \forall n > k: \frac{n^{2n-2k}}{n^{2n-2k}-1} \cdot \frac{n}{n+1} = \frac{n^{2n-2k+1}}{n^{2n-2k+1} + n^{2n-2k} - n - 1} \leq 1$$

$$\Leftrightarrow \forall n > k: n^{2n-2k+1} \leq n^{2n-2k+1} + n^{2n-2k} - n - 1$$

$$\Leftrightarrow \forall n > k: 0 \leq n^{2n-2k} - n - 1$$

$$\Leftrightarrow \forall n > k: n+1 \leq n^{2n-2k}$$

$$\Leftrightarrow \forall n > k: 1 + \frac{1}{n} \leq n^{2(n-k)-1}$$

Since $1 + \frac{1}{n}$ is always smaller or equal 2:

$$\Leftrightarrow \forall n > k: 2 \leq n^{2(n-k)-1}$$

$$\Leftrightarrow \forall n > k: 2 \leq n^{2-1} \wedge n^{2-1} \leq n^{2(n-k)-1}$$

$$\Leftrightarrow \forall n > k: 2 \leq n \wedge \text{true} \text{ , because } n > k$$

This is obviously true for $n > 1$. n can not be < 1 , since $n > k > 0$.

Therefore you can read the proof from the bottom up.

q.e.d.