

Anonymous Credentials for the EUDI Wallet: Closing the Gaps

EUDI ON – 25.6.2026

Andrea Flamini

Anja Lehmann

Alexandros Zacharakis

§ 16. The technical framework of the European Digital Identity Wallet shall:

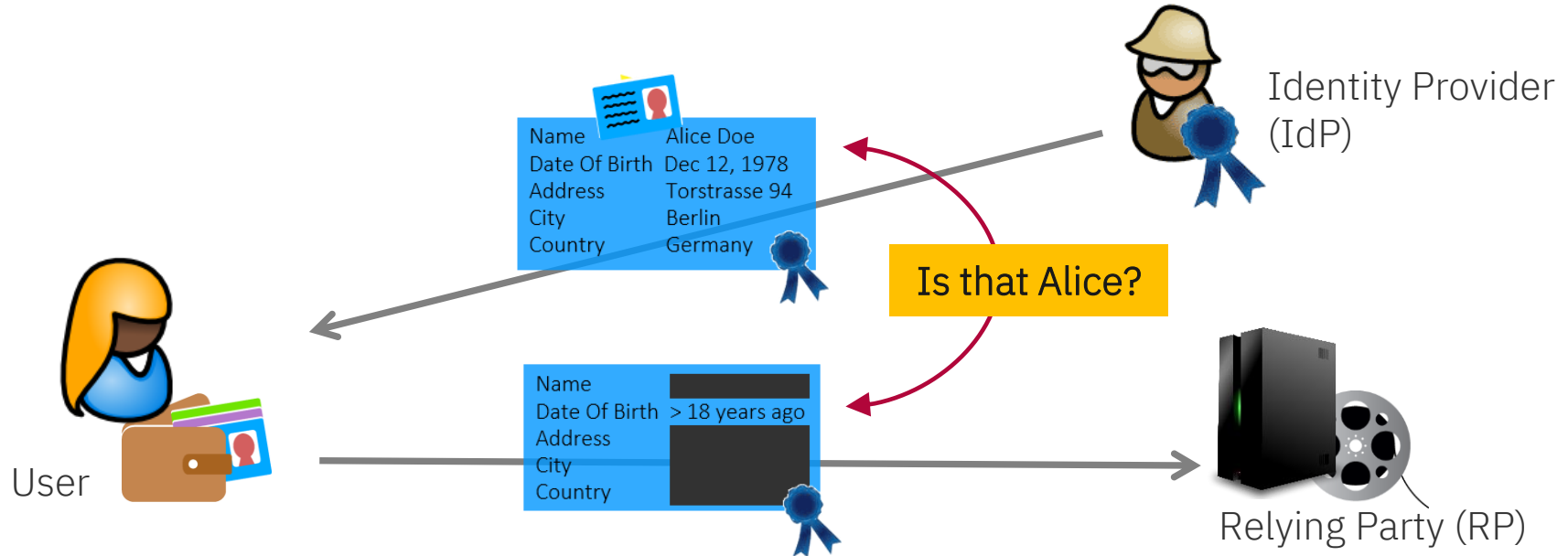
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- (b) enable privacy preserving techniques which ensure ~~unlikeability~~

Corrigendum 9.4.25: unlinkability! 😊

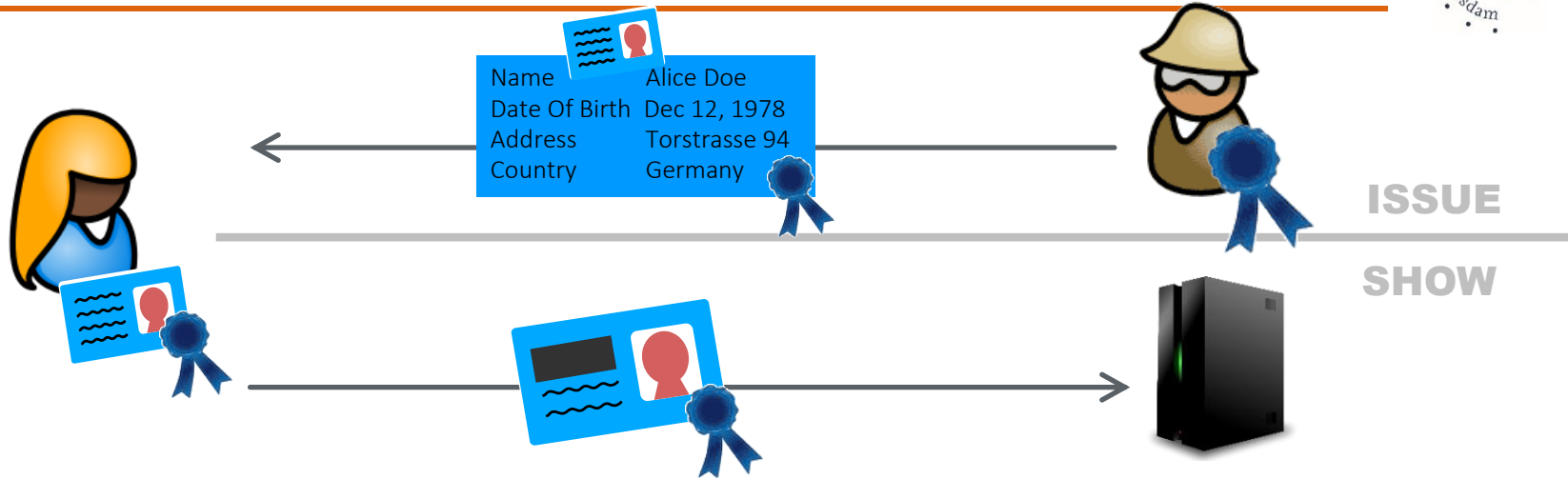
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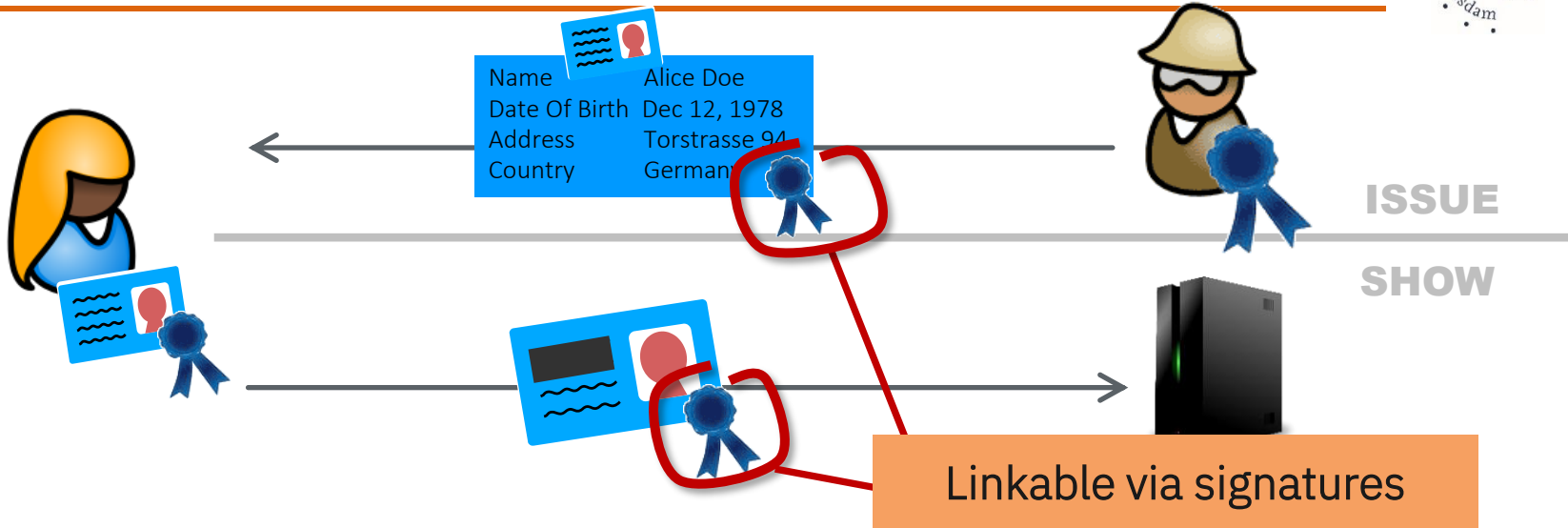
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EUDI | Solution with Classical Signatures

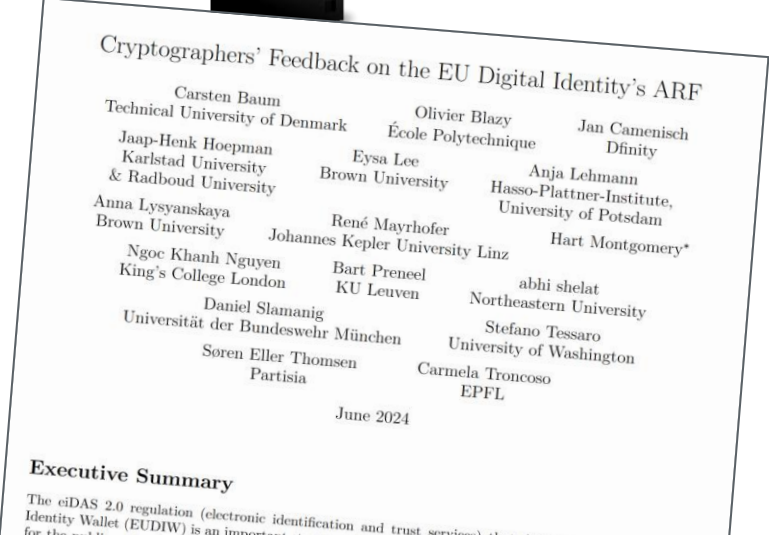
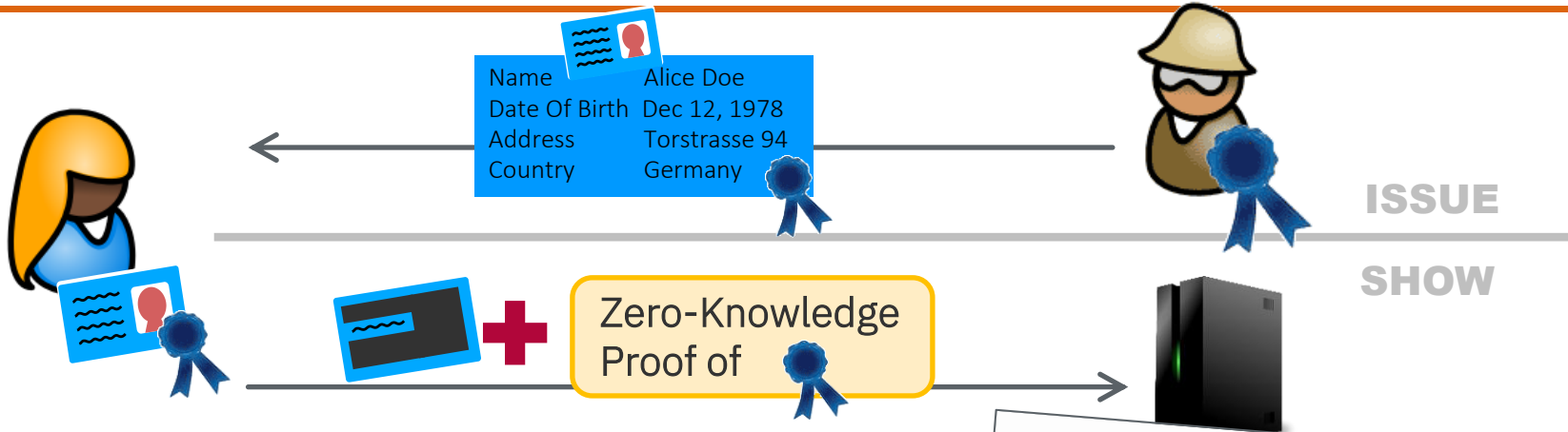


- Digital identity credentials based on cryptographic signatures

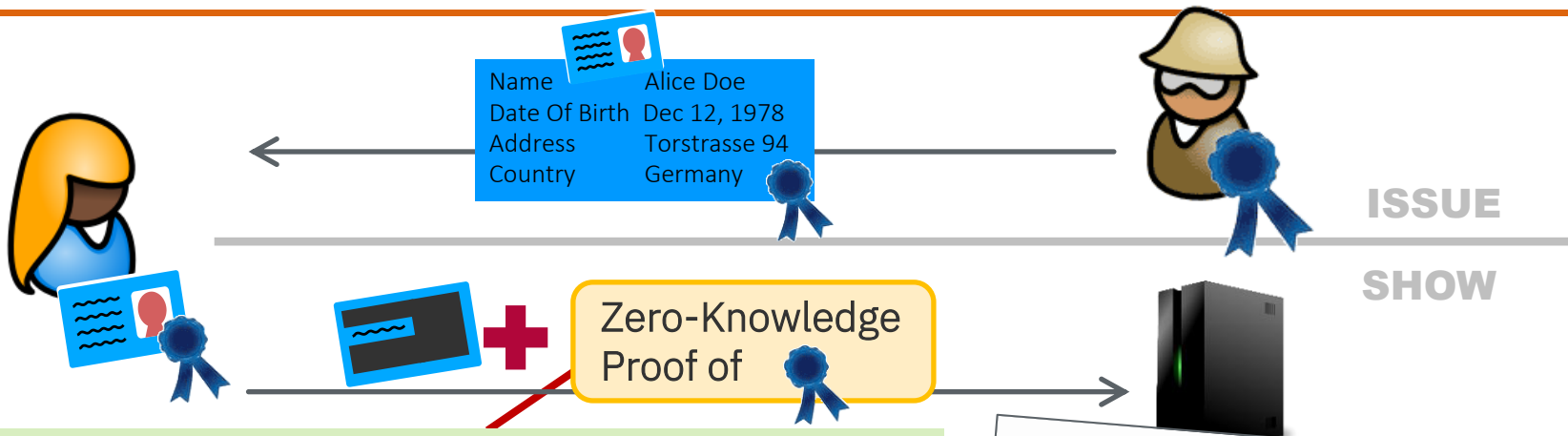


- Digital identity credentials based on cryptographic signatures
- Digital Signatures are great for security, but bad for privacy → unique identifier
- Batch issuance helps for unlinkability towards RPs (but is rather costly)
& does **not achieve** unlinkability/untraceability if IdP and RP collude

Anonymous Credentials & Zero-Knowledge Proofs



Anonymous Credentials & Zero-Knowledge Proofs

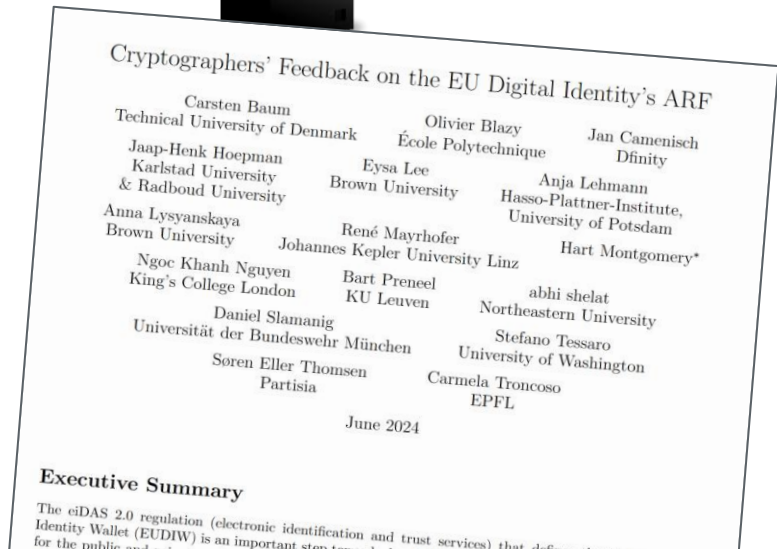


Zero-Knowledge Proof (ZKP)

Proof that reveals nothing beyond validity

→ multi-show unlinkability & no trust on IdP needed

Practical ZKP-based solutions exist since 25 years



Why does EUDI not use Anonymous Credentials (yet) ?



- Main focus so far: PID, i.e., national identity card

This comes with strong security requirements:

- 1) Must only use “Agreed Cryptographic Mechanisms”
 - 2) Security on Level of Assurance High
- ZKP not included in ACM or supported by secure hardware



KYC e.g.,
opening bank
account

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Adding lightweight ZKPs would still offer better privacy and security though

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Privacy

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Adding lightweight ZKPs would still offer better privacy and security though
- Current attention: age verification online (from age credential \neq PID)
High privacy needs, but lower security requirements than KYC
Right setting to deploy anonymous credentials



Towards Deploying Anonymous Credentials

1 | Standardisation of all required cryptographic protocols → interoperability

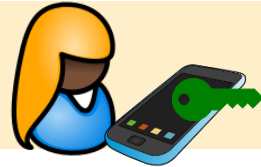
IETF: ongoing standards for all core crypto primitives (e.g., BBS)

ETSI: dedicated standard for ZKP-based EUDI Wallet under development



2 | Realize cloning-prevention for anonymous credentials

Usually done through device binding



→ Challenge: Secure Elements only support non-ZKP friendly signatures (ECDSA)

3 | Ensure that privacy is strong enough for real-world use case, e.g., for age proofs

→ Challenge: (Anonymous) Credentials usually reveal the issuer's identity

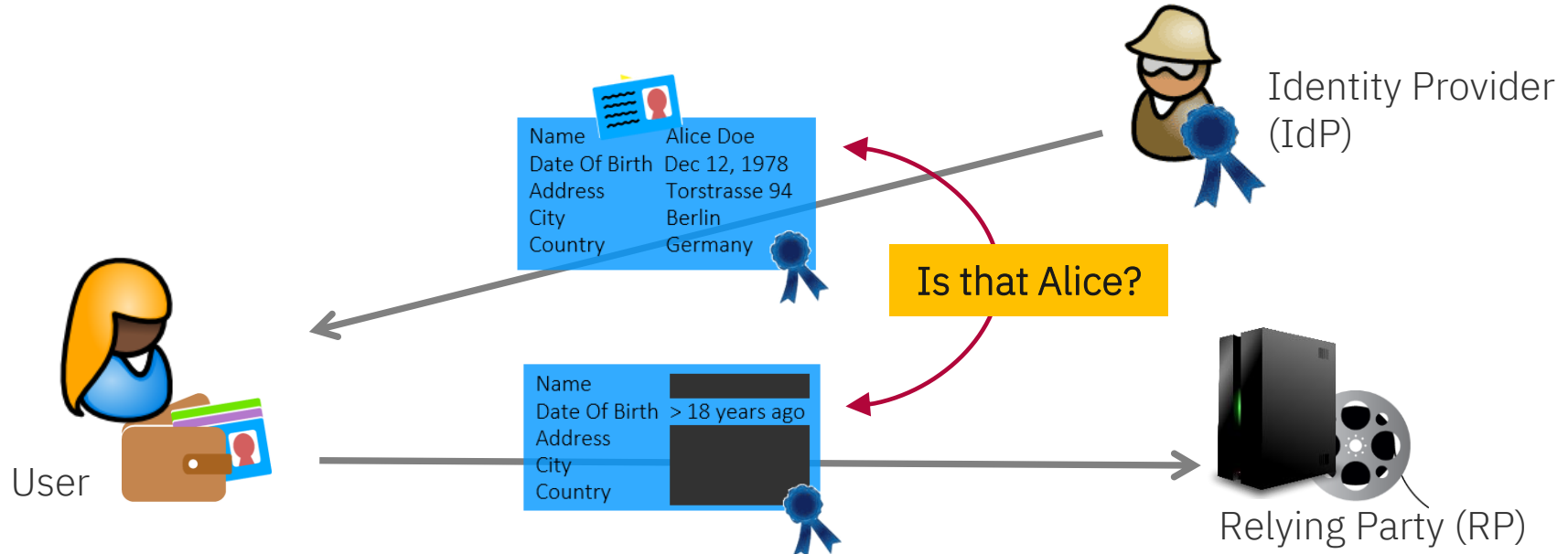
Device Binding for Anonymous Credentials

Alexandros Zacharakis

based on joint work with Sofia Celi, Anja Lehmann and Shai Levin

Recall Unlinkability

- Presentations cannot be traced back to a credential.
- Great for privacy! But...
- Makes misuse easier.



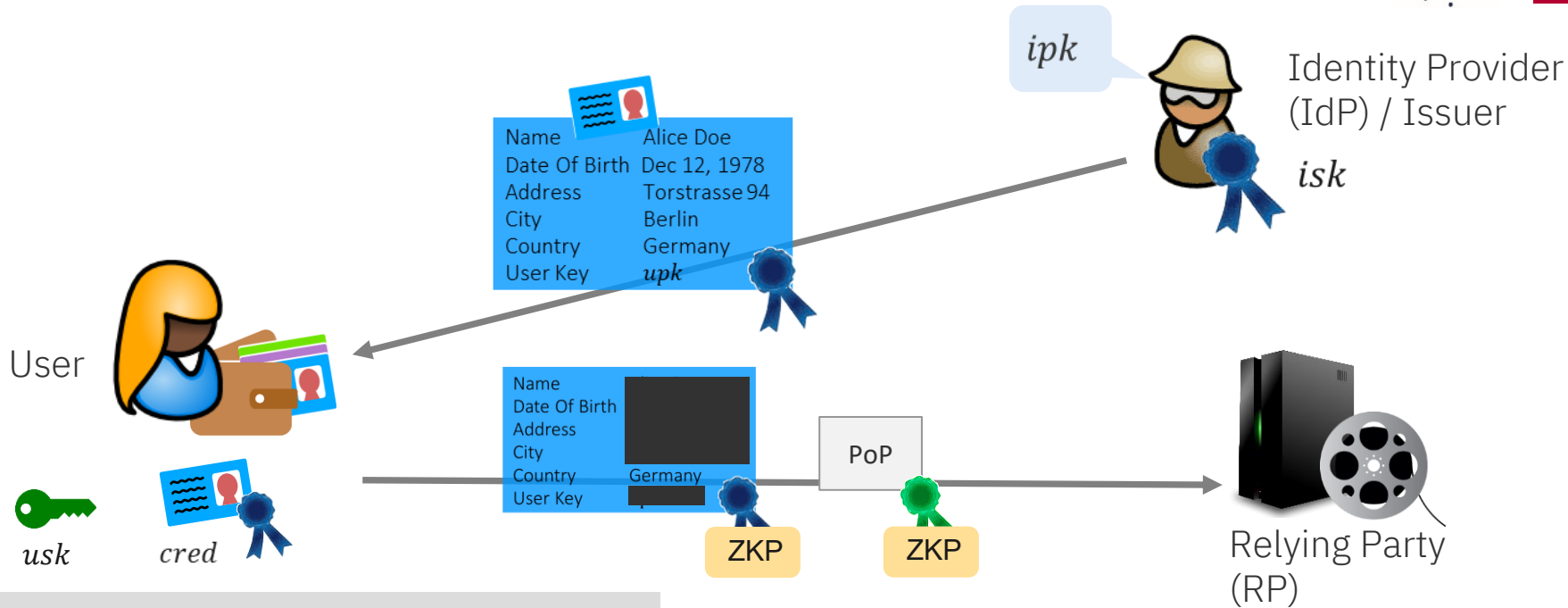
The core problem:

- Digital goods are easy to *clone*
- *Privacy vs Security*: misuse (e.g. sharing credentials) can happen *undetectably*
- Lack of security → reduced trust on the system: presentations are “vacant”

A (somewhat partial) solution

- Bind credentials to a device → harder to share “en masse”
- Valid presentation requires accessing the device
- Implemented via SE that stores an (inaccessible) secret signing key + *Proof-of-Possession*

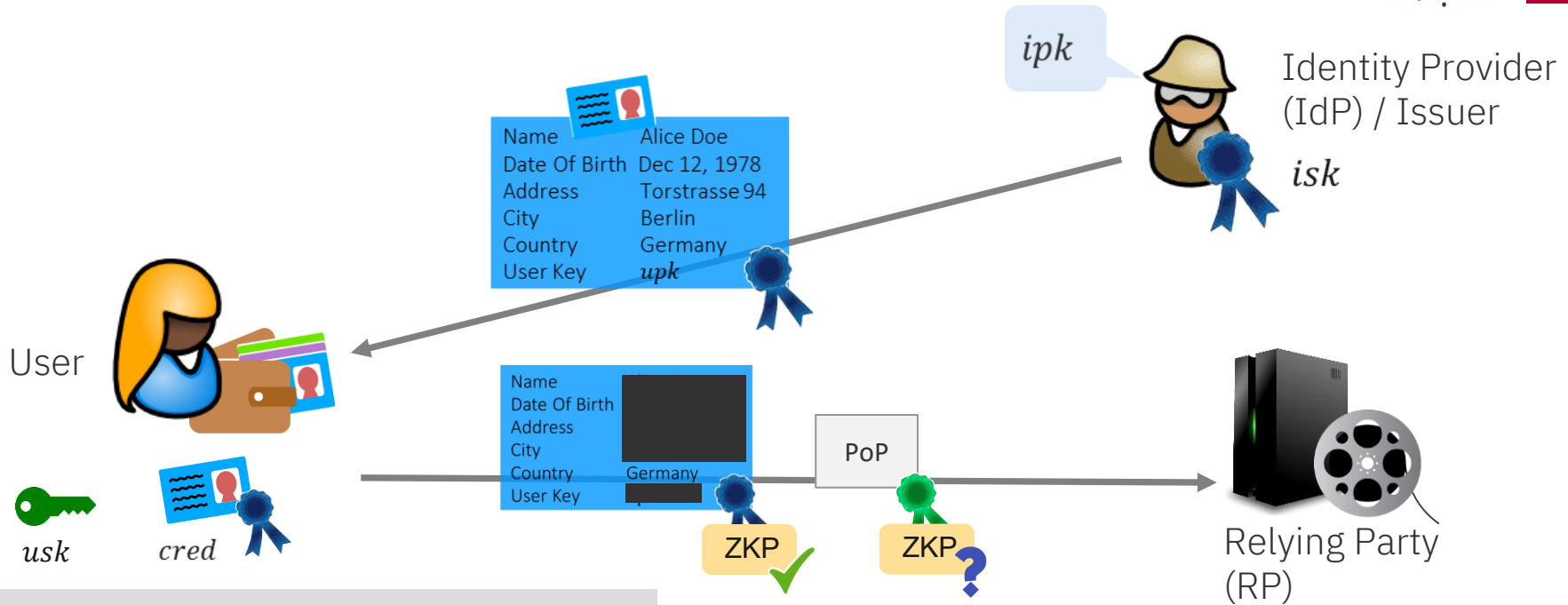
Device Bound Anonymous Credentials



$$cred := Sign(isk, (a_1, \dots, a_\ell, upk))$$

$$\sigma_{PoP} := Sign(usk, nonce)$$

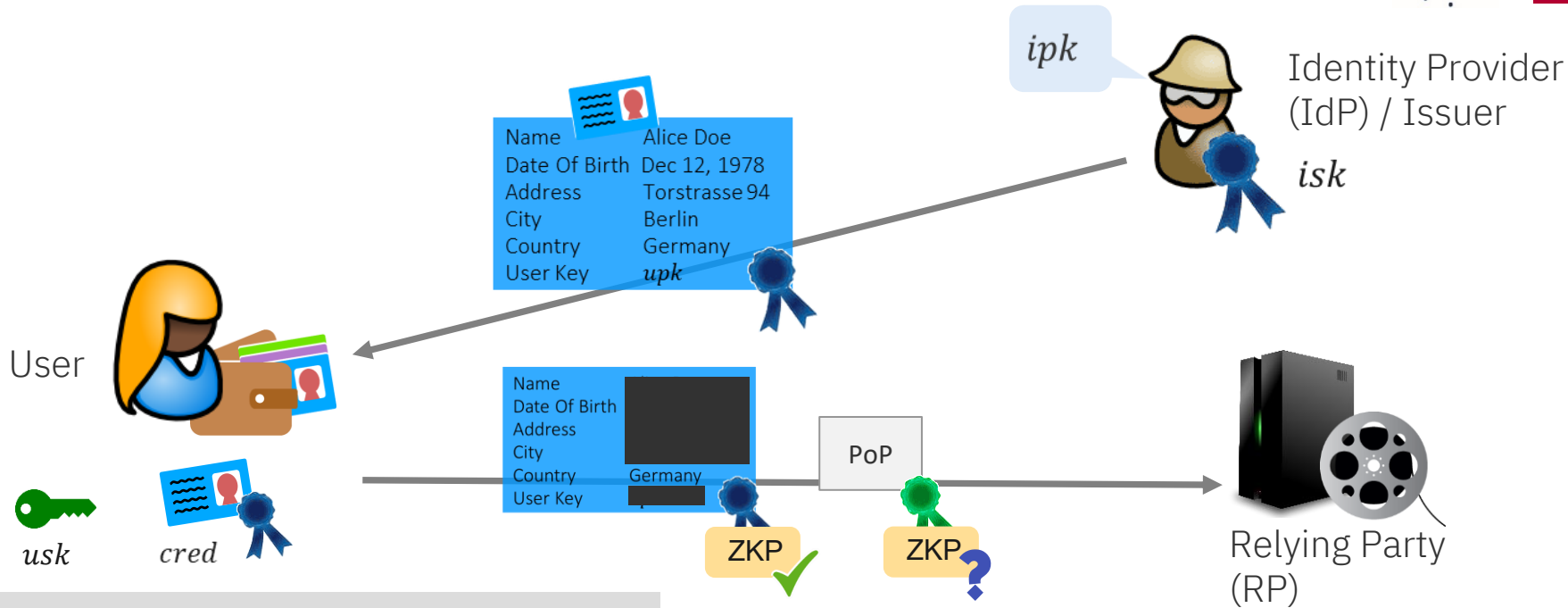
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- **Goal:** efficient ZKP for cred and PoP
- **Constraints:** ECDSA for PoP (at least for the short term)

- ECDSA was *not* designed to support ZKP
- We can use generic ZKP techniques → complex*



* This is an understatement.

How can we improve?

- PoP is an easier problem since the nonce is public
- Leverage the “structure” of ECDSA to get *specialized ZKPs* for PoP

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Device Binding for Anonymous Credentials on Legacy Phones

Sofia Celi*, Anja Lehmann†, Shai Levin‡, Alexandros Zacharakis†

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Email: {[anja.lehmann](mailto:anja.lehmann@hpi.de), [alexandros.zacharakis](mailto:alexandros.zacharakis@hpi.de)}@hpi.de

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Email: shai.levin@chalmers.se


The Current Landscape






Issuer	User	ZKP	Efficiency (rough guestimates)	Advantages

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BBS	ECDSA 	Schnorr (a lot for ECDSA)	~350ms, ~125kB	Relatively simple. No circuit Easy to extend & standardize
BBS	ECDSA 	Circuit (simple)	~350ms, ~1.5kB	Easy to extend & standardize Short proofs
ECDSA 	ECDSA 	Circuit (complex)	~400ms, ~300kB	No changes to issuer

Conclusion



Can we achieve efficient, privacy preserving device binding?

Can we do that now under the existing restrictions?

Are we done?

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- Various solutions for ECDSA with a *broad trade-off* spectrum

Are we done? **YES** and **NO**

- Efficient (and simple) solutions already exist!
- The remaining challenge: *standardization*

Towards Deploying Anonymous Credentials

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Anonymous Credentials with Issuer Hiding

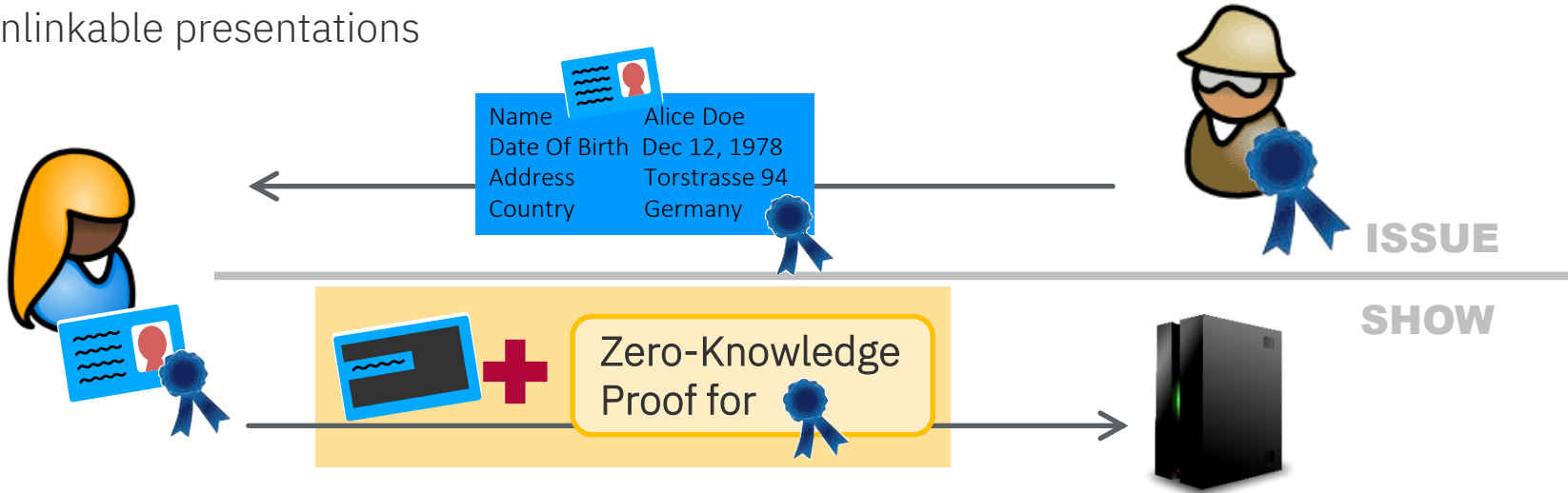
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Do anonymous credentials really enable minimal disclosure?

We have verifiable credentials supporting

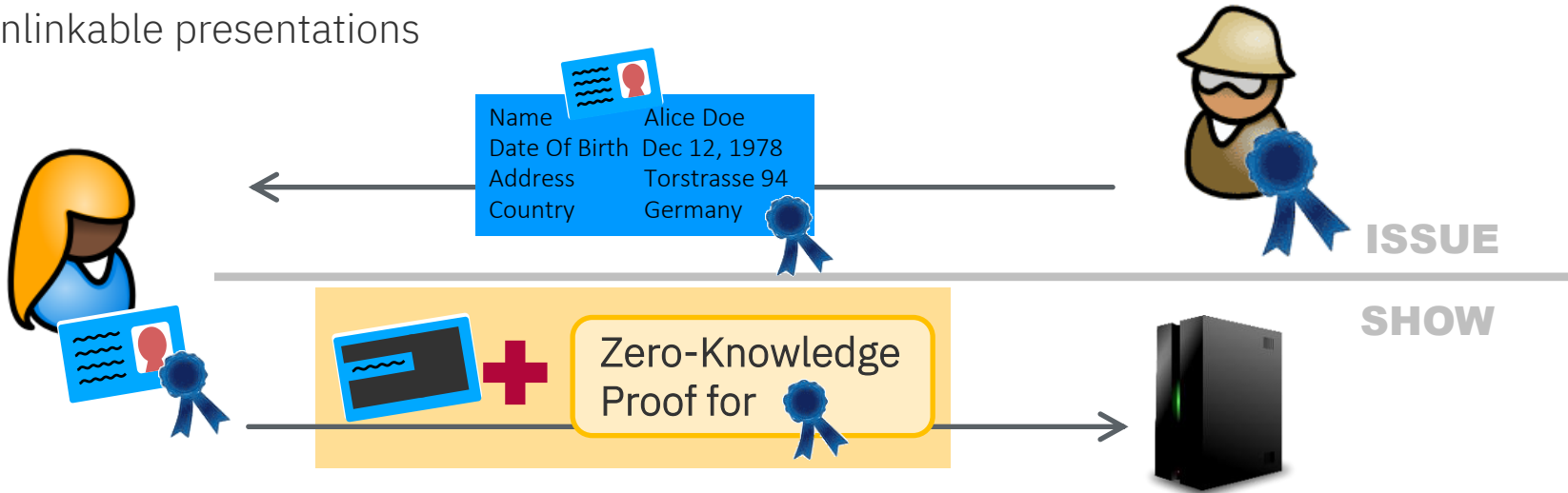
1. selective disclosure
2. unlinkable presentations



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Presentations **always reveal the issuer** to enable public verifiability

This can reveal your citizenship, your university, your device's manufacturer...

What we want to have

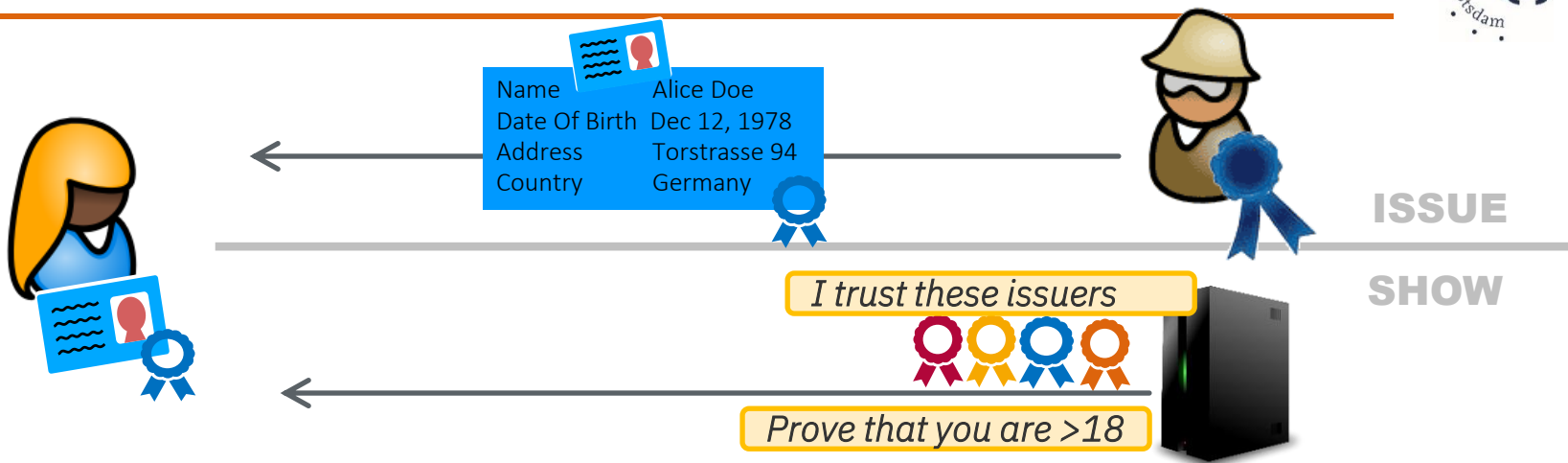


Name	Alice Doe
Date Of Birth	Dec 12, 1978
Address	Torstrasse 94
Country	Germany

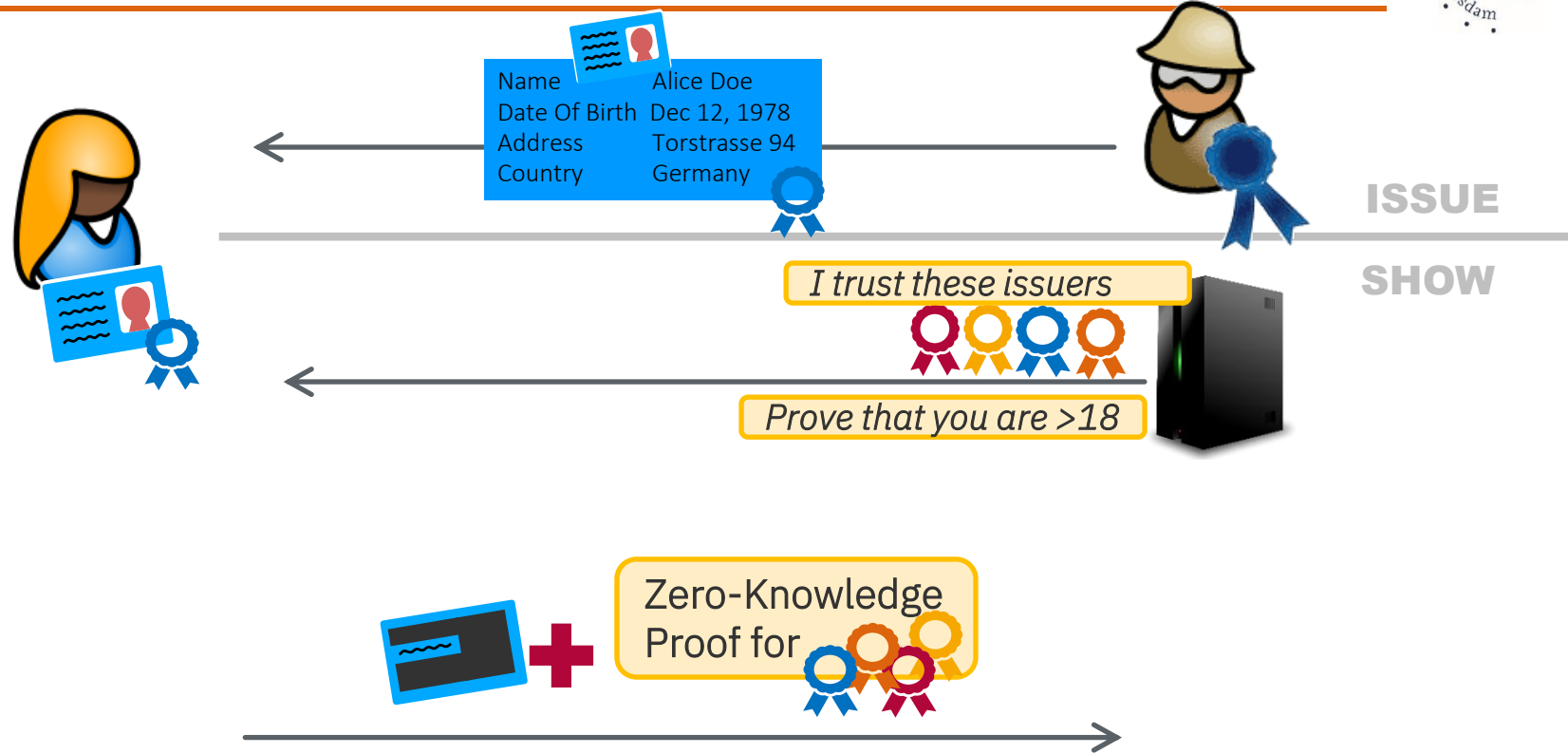


ISSUE

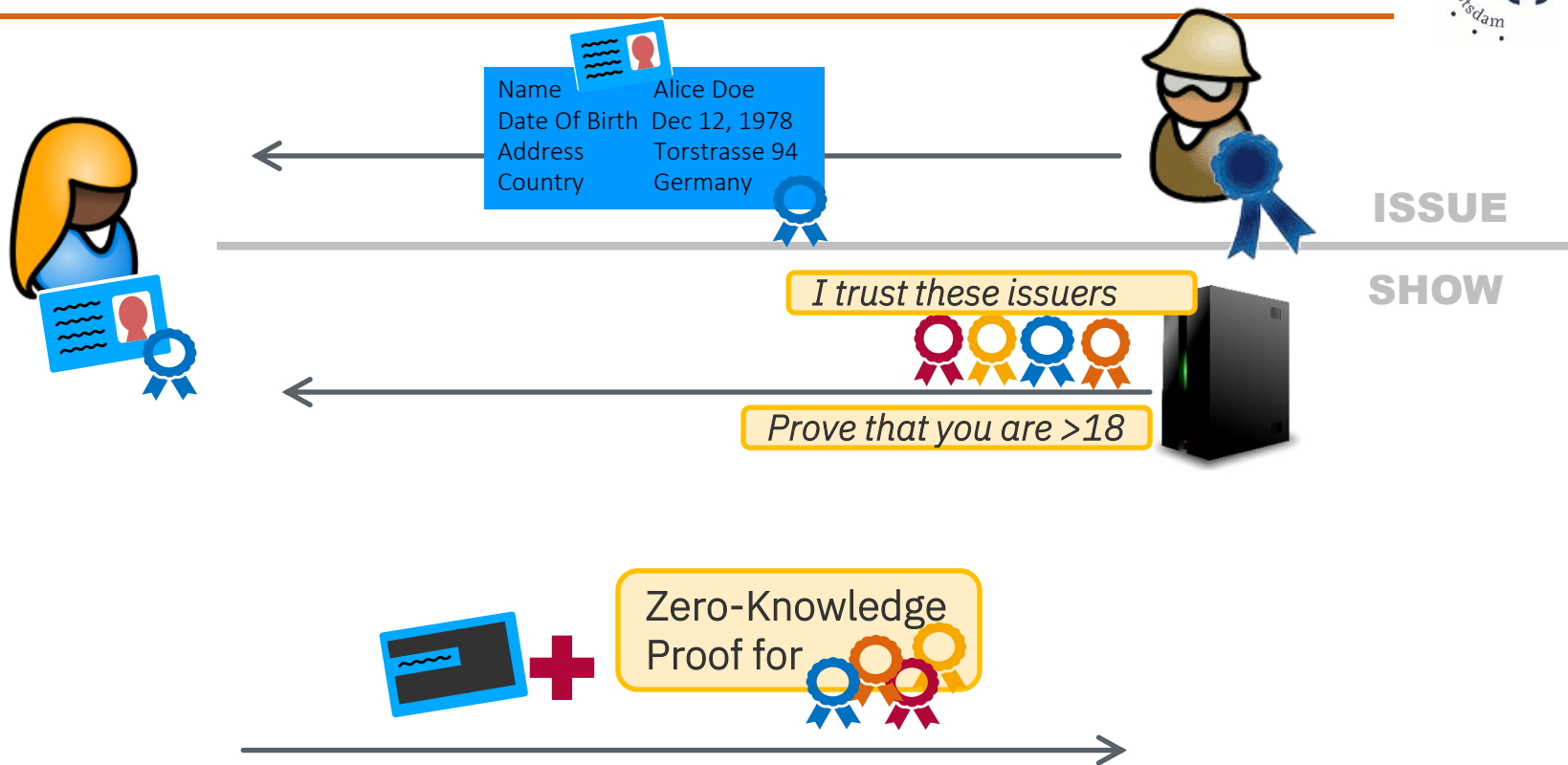
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**GOOD NEWS:
WE KNOW HOW TO DO IT!**

A useful property of BBS signatures



Given a BBS credential issued by



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We know how to transform it into a credential for the same attributes that verifies under a **random public key**



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This might look irrelevant, but it is actually **very useful!**

We have designed two solutions that match two use cases:

	Small sets of issuers	Large sets of issuers
Use case	Issuers are EU member states	Issuers are EU Universities
Trust	Trust only the issuers	Trust the issuers and a third party who certifies the public keys of the issuers in a set
Presentation size	overhead logarithmic in the number of issuers	Constant presentation size
Compatibility with BBS	✓	✓

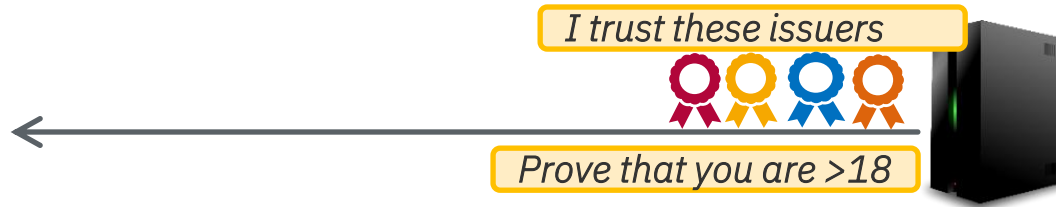
Simple construction for small policies



When the set of trusted issuers is small (e.g. <64) we have an efficient construction

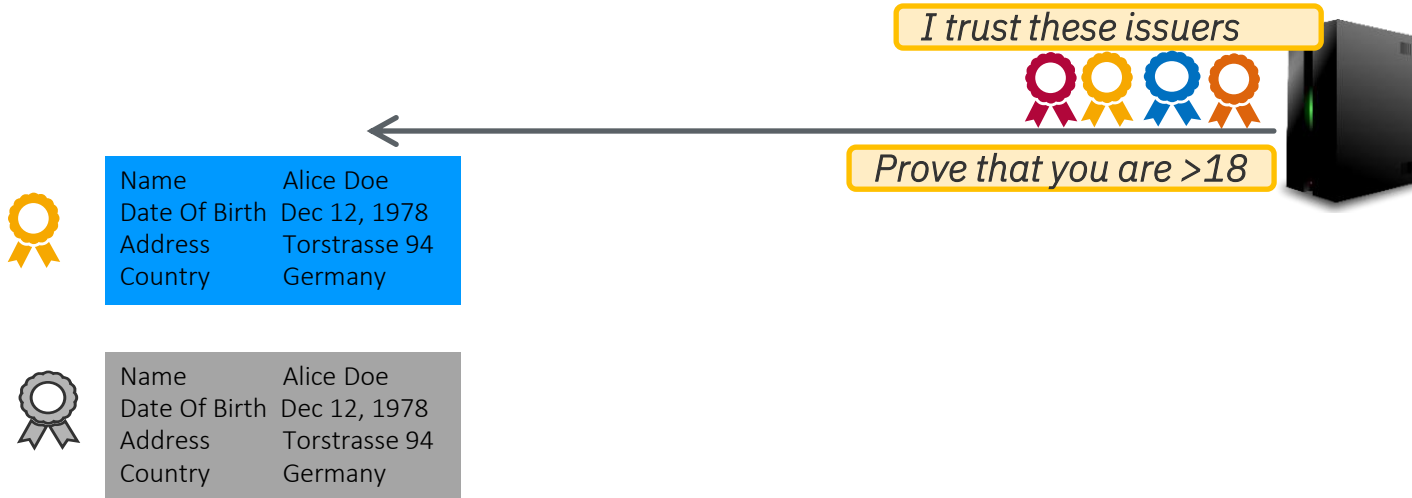
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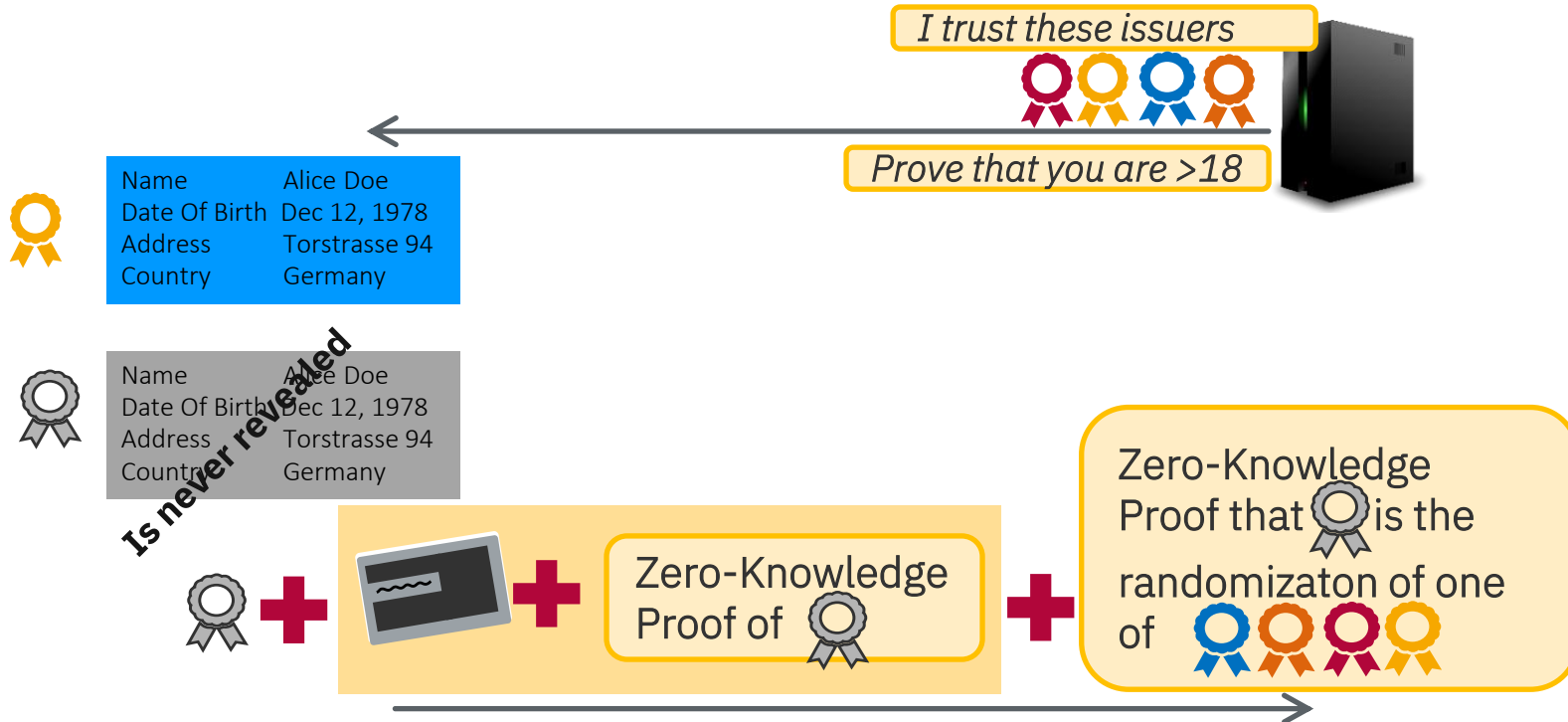
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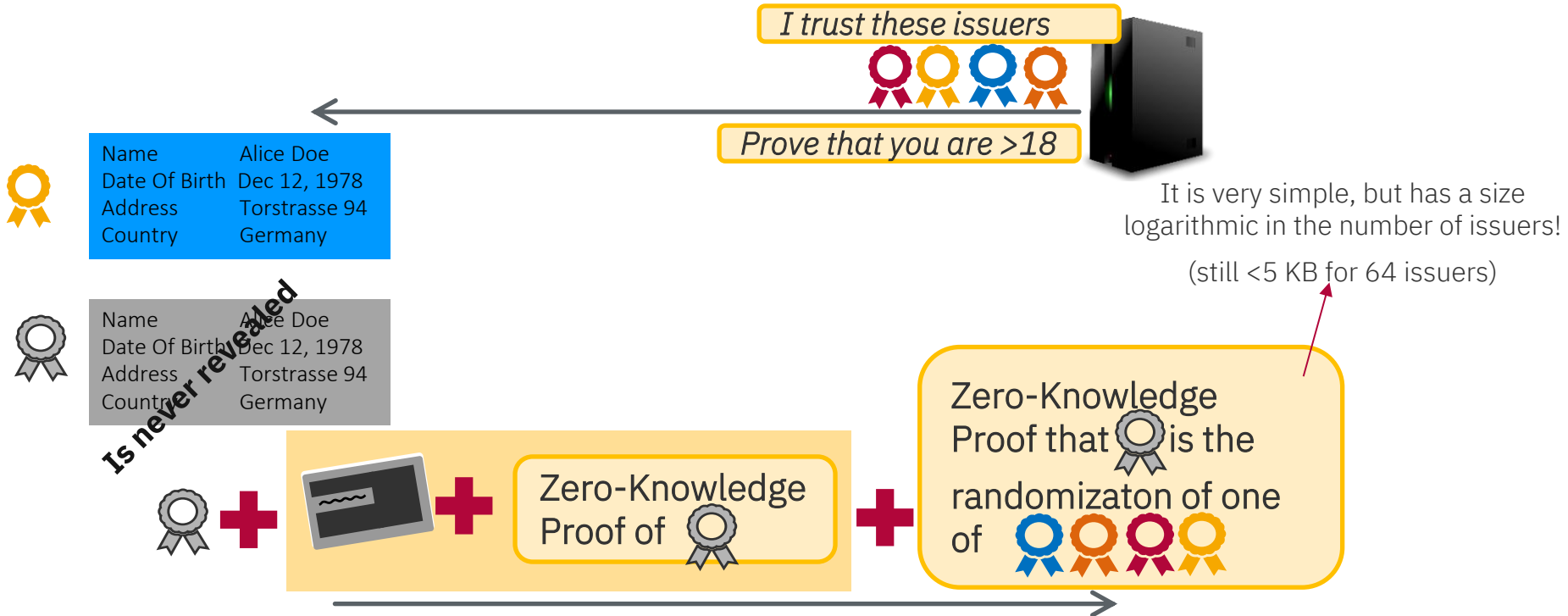
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Issuer-Hiding BBS-Based Anonymous Credentials without Policy Keys

Andrea Flamini¹, Karla Friedrichs¹, Jonathan Katz², Watson Ladd³,
Anja Lehmann¹, and Marek Sefranek⁴

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² Google, USA
jkatz2@gmail.com


³ Akamai, USA
watsonbladd@gmail.com

⁴ TU Wien, Austria
marek.sefranek@tuwien.ac.at

It is very simple, but has a size
arithmetic in the number of issuers!
(still <5 KB for 64 issuers)



Name	Alice
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Name	Alice
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Is never revealed



Proof of 



randomization of one
of 

Simple construction for larger policies



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Trusted party certifies

EU-Uni



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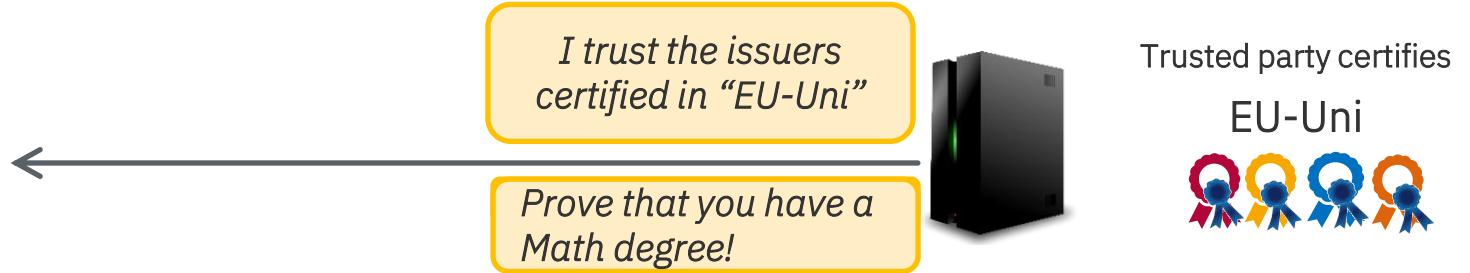
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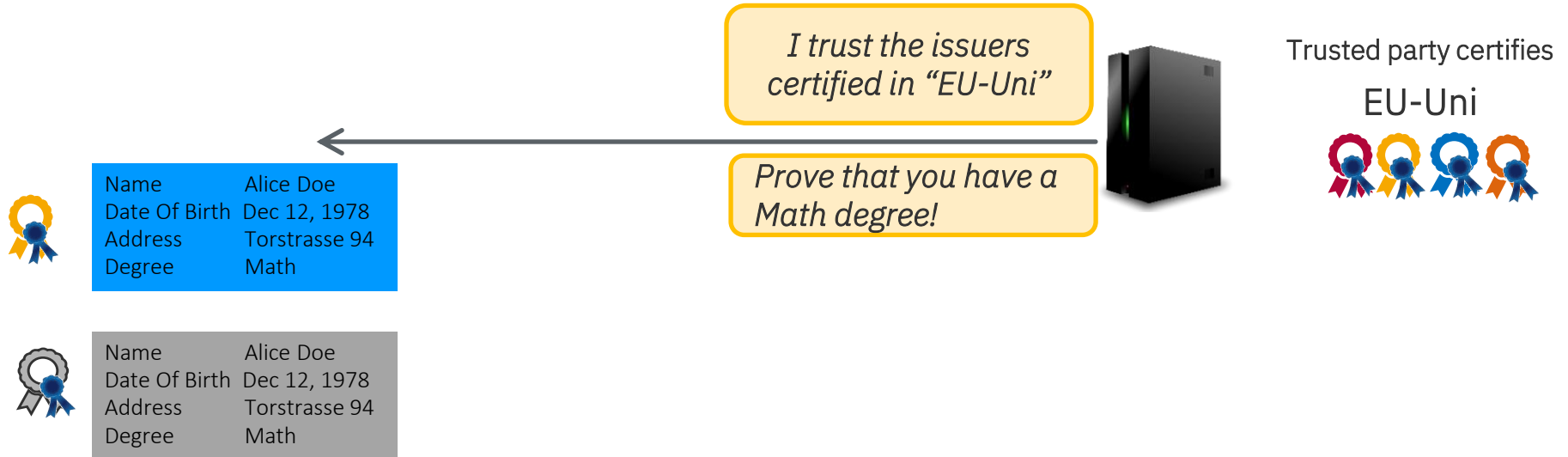


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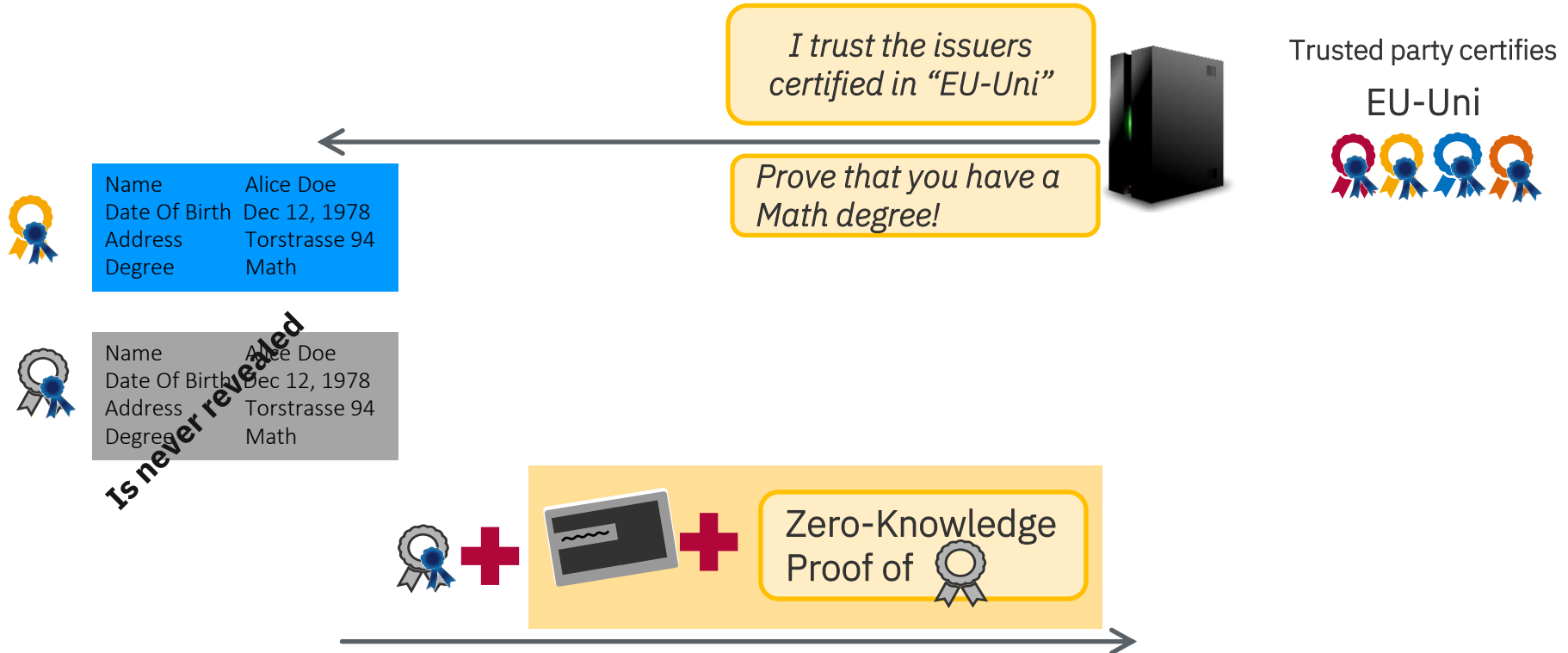
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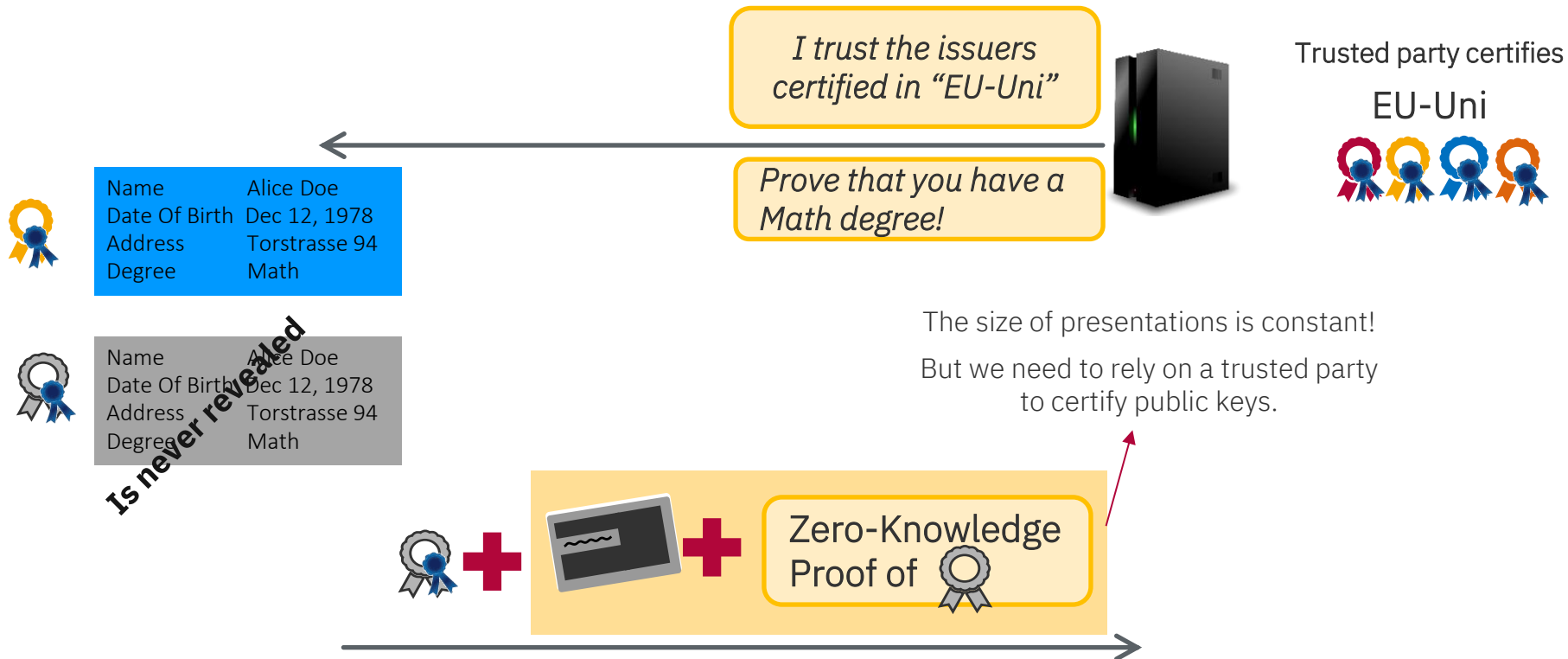
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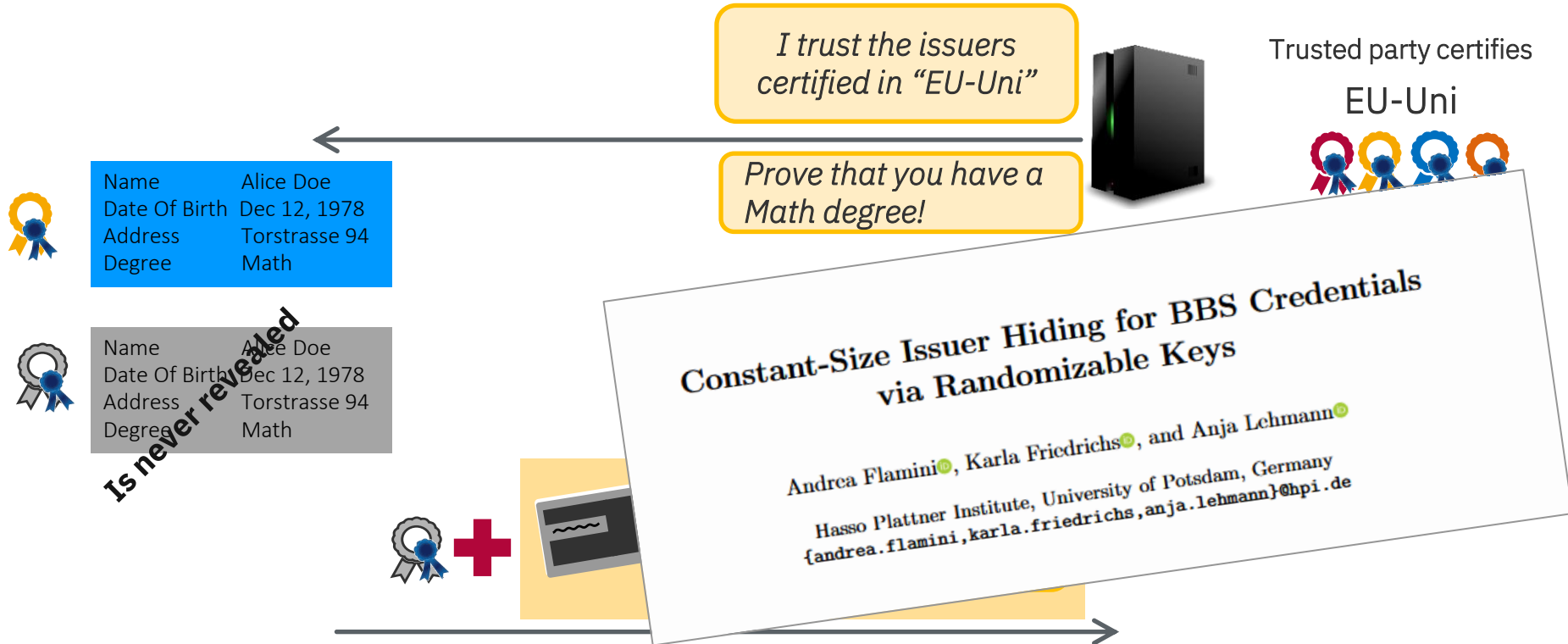
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The tricks we have shown do not work also for single-show credentials (due to salted hashes)

But we have shown that this is possible for BBS credentials!

- We know how to build Issuer-Hiding presentations efficiently.
- The protocols are **very simple**.
- The constructions **do not require updating the issuance** of BBS credentials.
- Issuer-Hiding presentations use **standard BBS presentations** as building blocks, and require only a few additional operations.

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→ Solved: Issuer-hiding protocols, e.g. for BBS

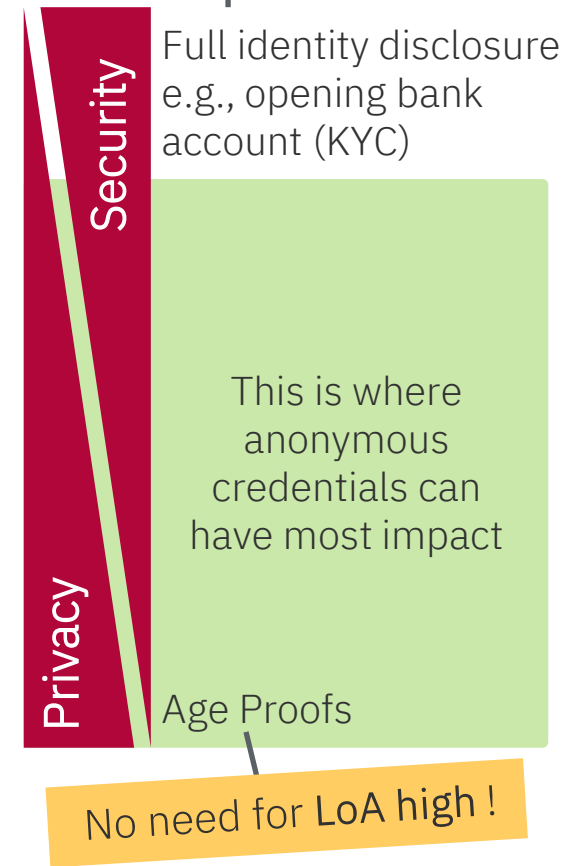


ZKP-EUDI Roadmap | Next Steps

No need for unlinkability!

Short/Mid-term: show feasibility and benefits

- Use ZKP for use cases with high privacy demands
- ZKP-compatible protocols (OID4VC) & data formats
- Many additional features: pseudonyms, blind issuance, ...
→ Build modular & crypto-agile protocols



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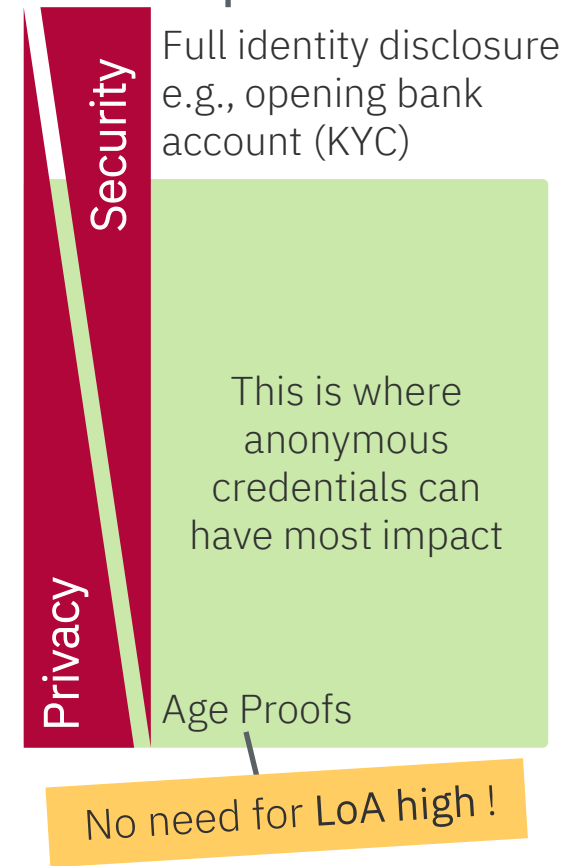
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Longterm: full post-quantum security

- Current solutions have PQC-privacy, but classic soundness
- PQC solutions need (a bit) more research & time to analyse
 - Current work provides concrete target for PQC research
Insights from pre-PQC serve as blueprint
 - Shape PQC base standards & hardware APIs now(ish)



Summary



- Anonymous credentials = only path to strong & multi-show unlinkability
Can even provide issuer hiding
- Start where privacy matters most, e.g., age proofs
(but whether we want age proofs is a different story ...)
- Main deployment blockers are solved: device binding & standards underway
- Cryptography for transitional phase is ready!

Summary

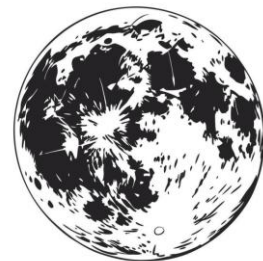
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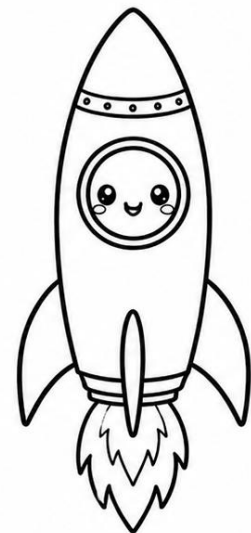
ARF: but you must use a car ...



ARF: ECDSA



eIDAS: privacy-preserving eID



*Ready to use & built for this:
anonymous credentials*

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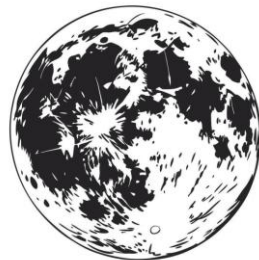
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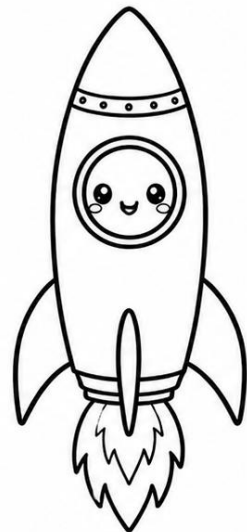
Interested in using ZKP? Please reach out:
<https://hpi.de/lehmann/eudi>



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