

Advanced Data Profiling

Introduction to Order Dependencies

WS 2023/2024
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Create Knowledge.**

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Overview

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Intuition: Bidirectional Order Dependency (bOD)

salary ↗ tax

salary	tax
5k	1k
6k	1.5k
8k	2k
10k	3k

age ↗ year-of-birth ↘

age	yob
19	2001
25	1995
25	1995
31	1989
45	1975

[cases ↗, r0 ↘] ↗ [dt ↗]

cases	r0	doubling time
46	1.5	4 d
57	1.8	7 d
102	1.7	10 d
102	1.4	12 d
188	1.4	14 d

- Lexicographical ordering
- Think SQL:
 - ... ORDER BY cases asc, r0 desc ⇔ ... ORDER BY „doubling time“ asc
- Default order direction: asc = ↗

Notation

Consistent with [1,7]

R	relation
r	relational instance of R
t, s, t_0, t_1, \dots	tuples of r
A, B, C	attributes of R
X, Y, Z	list of attributes
Y = [A T]	A is the head element of list Y and T the tail list
X, Y, Z	sets of attributes

Definition box

Defined concept name is bold

An *example definition* is a list of attributes denoted as **X = [A | A ∈ R]**.

**Lists and sets with the same name reference the same data:
Set X contains all distinct elements of list X**

Running Example [7]

ID	Month	Day	OCode	OFips	OState	DepDelay	ArrDelay	ArrDelGrp
t_0	12	3	CLE	39	Ohio	80	72	4
t_1	12	14	ORD	17	Illinois	-4	-10	-1
t_2	12	14	JFK	36	New York	-6	0	0
t_3	12	16	ORD	17	Illinois	13	59	3
t_4	12	20	CLE	39	Ohio	⊥	⊥	⊥
t_5	12	24	ORD	17	Illinois	5	-2	-1
t_6	12	27	BWI	24	Maryland	-5	-23	-2
t_7	12	28	ORD	17	Illinois	57	82	5
t_8	12	30	SGF	29	Missouri	2	0	0
t_9	12	30	ORD	17	Illinois	-5	-14	-1

⊥ = null values

Order Specification [1,7,8]

- Compare to SQL ORDER BY
 - ORDER BY A asc, B desc

An **order specification** is a list of marked attributes denoted as $\bar{X} = [\bar{A} | A \in R]$.

A **marked attribute** \bar{A} can have two directions: ascending $A \uparrow$ or descending $A \downarrow$. The default order direction is ascending (\uparrow).

- Examples

[salary \uparrow]

[cases \uparrow , r0 \downarrow]

cases \uparrow , r0 \downarrow , doubling_time \uparrow

[OCode, Month, Day]

OFips, DepDelay, ArrDelay

Order Specification

- Example table ordered by [*OCode* ↑, *Day* ↓]:

ID	Month	Day	OCode	OFips	OState	DepDelay	ArrDelay	ArrDelGrp
t_6	12	27	BWI	24	Maryland	-5	-23	-2
t_4	12	20	CLE	39	Ohio	⊥	⊥	⊥
t_0	12	3	CLE	39	Ohio	80	72	4
t_2	12	14	JFK	36	New York	-6	0	0
t_9	12	30	ORD	17	Illinois	-5	-14	-1
t_7	12	28	ORD	17	Illinois	57	82	5
t_5	12	24	ORD	17	Illinois	5	-2	-1
t_3	12	16	ORD	17	Illinois	13	59	3
t_1	12	14	ORD	17	Illinois	-4	-10	-1
t_8	12	30	SGF	29	Missouri	2	0	0

Lexicographical Order Operator [1,7,8]

- Lexicographical ordering
 - Numbers ordered numerically
 - Strings ordered alphabetically
 - Timestamps ordered chronologically
- Examples

Let \bar{X} be an order specification, s and t be two tuples in relation instance r . The **lexicographical order operator** $\preceq_{\bar{X}}$ is defined as follows:

$$s \preceq_{\bar{X}} t = \begin{cases} \bar{X} = [] \\ \bar{X} = [A \uparrow | \bar{T}] \wedge s_A < t_A \\ \bar{X} = [A \downarrow | \bar{T}] \wedge s_A > t_A \\ \bar{X} = [\bar{A} | \bar{T}] \wedge s_A = t_A \wedge s \preceq_{\bar{T}} t \end{cases}$$

$t_5 \preceq_{[Day]} t_9$
 $t_5 \preceq_{[DepDelay\downarrow]} t_9$
 $t_9 \preceq_{[ArrDelGrp, ArrDelay]} t_5$
 $t_5 \preceq_{[ArrDelGrp, ArrDelay\downarrow]} t_9$
 $t_5 \preceq_{[OCode]} t_9 \wedge t_9 \preceq_{[OCode]} t_5$
 $\Rightarrow t_5 =_{[OCode]} t_9$

ID	Month	Day	OCode	OFips	OState	DepDelay	ArrDelay	ArrDelGrp
t_9	12	30	ORD	17	Illinois	-5	-14	-1
t_5	12	24	ORD	17	Illinois	5	-2	-1

Bidirectional Order Dependency [1,7,8]

- If a bOD $\bar{X} \mapsto \bar{Y}$ holds:
 - when the tuples are ordered by \bar{X} , then they are also ordered by \bar{Y}
 - but not necessarily vice versa!
- Simpler form: (unidirectional) order dependencies (ODs) $X \mapsto Y$
 - All attributes have the same order direction (either all \uparrow or all \downarrow)
 - Y 's values are monotonically non-decreasing with respect to the values in X

Let \bar{X} and \bar{Y} be order specifications over relation R . A **bidirectional order dependency (bOD)** $\bar{X} \mapsto \bar{Y}$ (read X orders Y) specifies that ordering a relation instance r by \bar{X} also orders r by \bar{Y} .

Table r over the relation R satisfies a bOD $\bar{X} \mapsto \bar{Y}$, where $X \subset R$ and $Y \subset R$, if

$$\forall s, t \in r: s \preceq_{\bar{X}} t \implies s \preceq_{\bar{Y}} t.$$

Order Equivalence (OE) ↔

[1,7,8]

Two order specifications \bar{X} and \bar{Y} are **order equivalent**, denoted as $\bar{X} \leftrightarrow \bar{Y}$, iff $\bar{X} \mapsto \bar{Y}$ and $\bar{Y} \mapsto \bar{X}$.

$[OFips] \leftrightarrow [OState]$

ID	Month	Day	OCode	OFips	OState	DepDelay	ArrDelay	ArrDelGrp
t_1	12	14	ORD	17	Illinois	-4	-10	-1
t_3	12	16	ORD	17	Illinois	13	59	3
t_5	12	24	ORD	17	Illinois	5	-2	-1
t_7	12	28	ORD	17	Illinois	57	82	5
t_9	12	30	ORD	17	Illinois	-5	-14	-1
t_6	12	27	BWI	24	Maryland	-5	-23	-2
t_8	12	30	SGF	29	Missouri	2	0	0
t_2	12	14	JFK	36	New York	-6	0	0
t_0	12	3	CLE	39	Ohio	80	72	4
t_4	12	20	CLE	39	Ohio	⊥	⊥	⊥

Order Compatibility (OC) ~ [1,7,8]

Two order specifications \bar{X} and \bar{Y} are **order compatible**, denoted as $\bar{X} \sim \bar{Y}$, iff $\bar{X}\bar{Y} \leftrightarrow \bar{Y}\bar{X}$.

$[ArrDelGrp] \sim [ArrDelay]$

ID	Month	Day	OCODE	OFips	OState	DepDelay	ArrDelay	ArrDelGrp
t_6	12	27	BWI	24	Maryland	-5	-23	-2
t_9	12	30	ORD	17	Illinois	-5	-14	-1
t_1	12	14	ORD	17	Illinois	-4	-10	-1
t_5	12	24	ORD	17	Illinois	5	-2	-1
t_2	12	14	JFK	36	New York	-6	0	0
t_8	12	30	SGF	29	Missouri	2	0	0
t_3	12	16	ORD	17	Illinois	13	59	3
t_0	12	3	CLE	39	Ohio	80	72	4
t_7	12	28	ORD	17	Illinois	57	82	5
t_4	12	20	CLE	39	Ohio	⊥	⊥	⊥

Violation of ODs (Splits and Swaps)

[1,7,8]

$\bar{X} \mapsto \bar{Y}$ is valid only iff $\bar{X} \mapsto \bar{XY}$ and $\bar{X} \sim \bar{Y}$ are valid.

**Functional
Dependency
(FD)**

**Order
Compatibility
(OC)**

- 2 ways for violation of an OD $\bar{X} \mapsto \bar{Y}$

1. Violation of the FD-part $\bar{X} \mapsto \bar{XY}$

A **split** w.r.t. a bOD $\bar{X} \mapsto \bar{XY}$ is a pair of tuples t and s in relational instance r , such that $t_X = s_X$ but $t_Y \neq s_Y$.

2. Violation of the OC-part $\bar{X} \sim \bar{Y}$

A **swap** w.r.t. a bOD $\bar{X} \sim \bar{Y}$ ($\bar{XY} \leftrightarrow \bar{YX}$) is a pair of tuples t and s in relational instance r , such that $t <_{\bar{X}} s$ but $s <_{\bar{Y}} t$.

Example:
[C1] \mapsto [C2]

ID	C1	C2	C3
t_0	1	4	3
t_1	1	4	3
t_2	2	2	5
t_3	2	3	4

[C1] \mapsto [C1, C2] (FD: $C1 \rightarrow C2$)
 \rightarrow Split in t_2 and t_3

[C1] \sim [C2] ($[C1, C2] \leftrightarrow [C2, C1]$)
 \rightarrow Swaps in t_0 and t_2 ,
 t_0 and t_3 ,
 t_1 and t_2 ,
 t_1 and t_3

Axioms (Inference of ODs)

[8]

- Sound and complete axioms:

OD1: Reflexivity

$$\mathbf{XY} \mapsto \mathbf{X}$$

OD2: Prefix

$$\frac{\mathbf{X} \mapsto \mathbf{Y}}{\mathbf{ZX} \mapsto \mathbf{ZY}}$$

OD3: Normalization

$$\mathbf{WXYXV} \leftrightarrow \mathbf{WXYV}$$

OD4: Transitivity

$$\frac{\begin{array}{l} \mathbf{X} \mapsto \mathbf{Y} \\ \mathbf{Y} \mapsto \mathbf{Z} \end{array}}{\mathbf{X} \mapsto \mathbf{Z}}$$

OD5: Suffix

$$\frac{\mathbf{X} \mapsto \mathbf{Y}}{\mathbf{X} \leftrightarrow \mathbf{YX}}$$

OD6: Chain

$$\frac{\begin{array}{l} \mathbf{X} \sim \mathbf{Y}_1 \\ \forall_{i \in [1, n-1]} \mathbf{Y}_i \sim \mathbf{Y}_{i+1} \\ \mathbf{Y}_n \sim \mathbf{Z} \\ \forall_{i \in [1, n]} \mathbf{Y}_i \mathbf{X} \sim \mathbf{Y}_i \mathbf{Z} \end{array}}{\mathbf{X} \sim \mathbf{Z}}$$

Axioms for bidirectional ODs (Inference of bODs) [1]

- Sound and complete axioms:

1. Reflexivity^L

$$\overline{XY} \mapsto \overline{X}$$

2. Prefix^L

$$\frac{\overline{X} \mapsto \overline{Y}}{\overline{ZX} \mapsto \overline{ZY}}$$

3. Transitivity^L

$$\frac{\overline{X} \mapsto \overline{Y} \quad \overline{Y} \mapsto \overline{Z}}{\overline{X} \mapsto \overline{Z}}$$

4. Normalization^L

$$\overline{WXYXV} \leftrightarrow \overline{WXYV}$$

5. Suffix^L

$$\frac{\overline{X} \mapsto \overline{Y}}{\overline{X} \leftrightarrow \overline{YX}}$$

6. Chain^L

$$\overline{X} \sim \overline{Y_1}$$

$$\forall_{i \in [1, n-1]} \overline{Y_i} \sim \overline{Y_{i+1}}$$

$$\overline{Y_n} \sim \overline{Z}$$

$$\frac{\forall_{i \in [1, n]} \overline{Y_i X} \sim \overline{Y_i Z}}{\overline{X} \sim \overline{Z}}$$

7. Reverse^L

$$\frac{\overline{X} \mapsto \overline{Y}}{\overline{X \downarrow} \mapsto \overline{Y \downarrow}}$$

8. Constant^L

$$\overline{XZ} \mapsto \overline{XY}$$

$$\overline{XZ} \mapsto \overline{XY \downarrow}$$

$$\frac{}{\overline{X} \mapsto \overline{XY}}$$

Set-based Canonical Form [1,7]

- Equivalence class

An **equivalence class** $\mathcal{E}(t_X)$ w.r.t. a given attribute set $X \in R$ groups tuples s and t together when their projection on X is equal: $\mathcal{E}(t_X) = \{s \in r \mid s_X = t_X\}$.

- Partition

A **partition** Π_X is a set of disjoint equivalence classes with the same set of attributes: $\Pi_X = \{\mathcal{E}(t_X) \mid t \in R\}$.

- Examples

$$\Pi_{\{Month\}} = \{\{t_0, \dots, t_9\}\}$$

$$\mathcal{E}(12) = \{t_0, \dots, t_9\}$$

$$\Pi_{\{OState\}} = \{\{t_0, t_4\}, \{t_1, t_3, t_5, t_7, t_9\}, \{t_2\}, \{t_6\}, \{t_8\}\}$$

$$\mathcal{E}(Ohio) = \{t_0, t_4\}$$

$$\mathcal{E}(New York) = \{t_2\}$$

$$\mathcal{E}(Missouri) = \{t_8\}$$

$$\mathcal{E}(Illinois) = \{t_1, t_3, t_5, t_7, t_9\}$$

$$\mathcal{E}(Maryland) = \{t_6\}$$

ID	Month	Day	OCode	OFips	OState
t_0	12	3	CLE	39	Ohio
t_1	12	14	ORD	17	Illinois
t_2	12	14	JFK	36	New York
t_3	12	16	ORD	17	Illinois
t_4	12	20	CLE	39	Ohio
t_5	12	24	ORD	17	Illinois
t_6	12	27	BWI	24	Maryland
t_7	12	28	ORD	17	Illinois
t_8	12	30	SGF	29	Missouri
t_9	12	30	ORD	17	Illinois

Set-based Canonical Form [1,7]

- Constant bOD from $X: [] \mapsto \bar{A}$

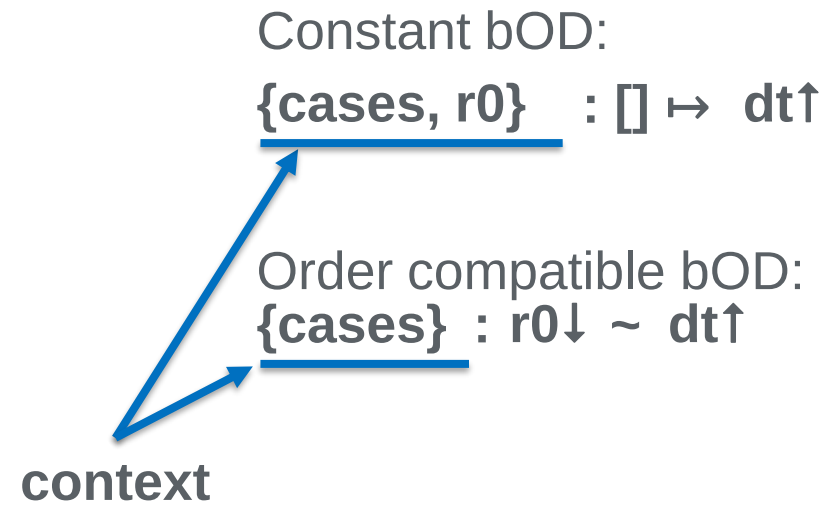
*A **constant bOD** $X: [] \mapsto \bar{A}$ is a marked attribute \bar{A} that is constant within each equivalence class w.r.t. the set of attributes in the context $X \in R$.*

- Order compatible bOD form: $X: \bar{A} \sim \bar{B}$

*An **order compatible bOD** $X: \bar{A} \sim \bar{B}$ are two marked attributes \bar{A} and \bar{B} that are order compatible within each equivalence class w.r.t. the set of attributes in the context $X \in R$.*

Set-based Canonical Form

	cases	r0	doubling time
t_1	46	1.5	4 d
t_2	57	1.8	7 d
t_3	102	1.7	10 d
t_4	102	1.4	12 d
t_5	188	1.4	14 d



Axioms (Inference for Set-based bODs) ^[1]

- Sound and complete axioms:

$$1. \text{ Reflexivity}^S \\ \mathcal{X}: [] \mapsto \overline{A}, \forall A \in \mathcal{X}$$

$$2. \text{ Identity}^S \\ \mathcal{X}: \overline{A} \sim \overline{A}$$

$$3. \text{ Commutativity}^S \\ \frac{\mathcal{X}: \overline{A} \sim \overline{B}}{\mathcal{X}: \overline{B} \sim \overline{A}}$$

$$4. \text{ Strengthen}^S \\ \frac{\mathcal{X}: [] \mapsto \overline{A} \\ \mathcal{X}A: [] \mapsto \overline{B}}{\mathcal{X}: [] \mapsto \overline{B}}$$

$$5. \text{ Propagate}^S \\ \frac{\mathcal{X}: [] \mapsto \overline{A}}{\mathcal{X}: \overline{A} \sim \overline{B}}$$

$$6. \text{ Augmentation-I}^S \\ \frac{\mathcal{X}: [] \mapsto \overline{A}}{\mathcal{Z}\mathcal{X}: [] \mapsto \overline{A}}$$

$$7. \text{ Augmentation-II}^S \\ \frac{\mathcal{X}: \overline{A} \sim \overline{B}}{\mathcal{Z}\mathcal{X}: \overline{A} \sim \overline{B}}$$

$$8. \text{ Chain}^S \\ \mathcal{X}: \overline{A} \sim \overline{B}_1 \\ \forall_{i \in [1, n-1]}, \mathcal{X}: \overline{B}_i \sim \overline{B}_{i+1} \\ \mathcal{X}: \overline{B}_n \sim \overline{C} \\ \frac{\forall_{i \in [1, n]}, \mathcal{X}B_i : \overline{A} \sim \overline{C}}{\mathcal{X}: \overline{A} \sim \overline{C}}$$

$$9. \text{ Reverse-I}^S \\ \frac{\mathcal{X}: [] \mapsto \overline{A}}{\mathcal{X}: [] \mapsto \overline{A}\uparrow}$$

$$10. \text{ Reverse-II}^S \\ \frac{\mathcal{X}: \overline{A} \sim \overline{B}}{\mathcal{X}: \overline{A}\uparrow \sim \overline{B}\uparrow}$$

$$11. \text{ Constant}^S \\ \mathcal{X}A: [] \mapsto \overline{B} \\ \mathcal{X}: \overline{A} \sim \overline{B} \\ \frac{\mathcal{X}: \overline{A} \sim \overline{B}\uparrow}{\mathcal{X}: [] \mapsto \overline{B}}$$

Axioms (Inference for Set-based bODs) ^[1]

- Examples

ID	Region	Name
1	AFRICA	Algeria
2	AFRICA	Angola
3	AFRICA	Benin
...
54	AFRICA	Zimbabwe
55	AMERICA	Argentina
...
193	OCEANIA	Vanuatu

3. Commutativity^S

$$\frac{\mathcal{X}: \overline{A} \sim \overline{B}}{\mathcal{X}: \overline{B} \sim \overline{A}}$$

{name}: [] ↦ region
 {name}: [] ↦ id
 {id}: [] ↦ name
 {id}: [] ↦ region

{}: id↑ ~ region↑
 {region}: id↑ ~ name↑



Axioms (Inference for Set-based bODs) ^[1]

- Examples

ID	Region	Name
1	AFRICA	Algeria
2	AFRICA	Angola
3	AFRICA	Benin
...
54	AFRICA	Zimbabwe
55	AMERICA	Argentina
...
193	OCEANIA	Vanuatu

$$5. \text{Propagate}^S \frac{\mathcal{X}: [] \mapsto \overline{A}}{\mathcal{X}: \overline{A} \sim B}$$

{name}: [] ↦ region
 {name}: [] ↦ id
{id}: [] ↦ name
 {id}: [] ↦ region

 {}: id↑ ~ region↑
 {region}: id↑ ~ name↑



{id}: name↑ ~ region↑
{id}: name↑ ~ region↓
{id}: name↑ ~ id↑
 ...

Axioms (Inference for Set-based bODs) ^[1]

- Examples

ID	Region	Name
1	AFRICA	Algeria
2	AFRICA	Angola
3	AFRICA	Benin
...
54	AFRICA	Zimbabwe
55	AMERICA	Argentina
...
193	OCEANIA	Vanuatu

$$6. \text{Augmentation-}I^S$$

$$\frac{\mathcal{X}: [] \mapsto \bar{A}}{\mathcal{Z}\mathcal{X}: [] \mapsto \bar{A}}$$

{name}: [] ↦ region
 {name}: [] ↦ id
{id}: [] ↦ name
 {id}: [] ↦ region

 {}: id↑ ~ region↑
 {region}: id↑ ~ name↑



{id,region}: [] ↦ name
 {id,name}: [] ↦ name
 {id,region,name}: [] ↦ name

Axioms (Inference for Set-based bODs) ^[1]

- Examples

ID	Region	Name
1	AFRICA	Algeria
2	AFRICA	Angola
3	AFRICA	Benin
...
54	AFRICA	Zimbabwe
55	AMERICA	Argentina
...
193	OCEANIA	Vanuatu

$$7. \text{Augmentation-II}^S$$

$$\frac{\mathcal{X}: \overline{A} \sim \overline{B}}{\mathcal{Z}\mathcal{X}: \overline{A} \sim \overline{B}}$$

{name}: [] ↦ region
 {name}: [] ↦ id
 {id}: [] ↦ name
 {id}: [] ↦ region

{}: id↑ ~ region↑
 {region}: id↑ ~ name↑



{name}: id↑ ~ region↑
 {name,id}: id↑ ~ region↑
 ...

Axioms (Inference for Set-based bODs) ^[1]

- Examples

ID	Region	Name
1	AFRICA	Algeria
2	AFRICA	Angola
3	AFRICA	Benin
...
54	AFRICA	Zimbabwe
55	AMERICA	Argentina
...
193	OCEANIA	Vanuatu

$$9. \text{Reverse-}I^S \frac{\mathcal{X}: [] \mapsto \bar{A}}{\mathcal{X}: [] \mapsto \bar{A}\updownarrow}$$

{name}: [] ↦ region
{name}: [] ↦ **id**
 {id}: [] ↦ name
 {id}: [] ↦ region


{name}: [] ↦ **id**↓

{}: id↑ ~ region↑
 {region}: id↑ ~ name↑

Axioms (Inference for Set-based bODs) ^[1]

- Examples

ID	Region	Name
1	AFRICA	Algeria
2	AFRICA	Angola
3	AFRICA	Benin
...
54	AFRICA	Zimbabwe
55	AMERICA	Argentina
...
193	OCEANIA	Vanuatu

$$10. \text{Reverse-II}^S$$

$$\frac{\mathcal{X}: \overline{A} \sim \overline{B}}{\mathcal{X}: \overline{A\uparrow} \sim \overline{B\uparrow}}$$

{name}: [] ↦ region
 {name}: [] ↦ id
 {id}: [] ↦ name
 {id}: [] ↦ region

{}: id↑ ~ region↑
 {region}: id↑ ~ name↑



Axioms (Inference for Set-based bODs) ^[1]

- Examples

ID	Region	Name
1	AFRICA	Algeria
2	AFRICA	Angola
3	AFRICA	Benin
...
54	AFRICA	Zimbabwe
55	AMERICA	Argentina
...
193	OCEANIA	Vanuatu

{name}: [] ↦ region
 {name}: [] ↦ id
 {id}: [] ↦ name
{id}: [] ↦ region

{}: id↑ ~ region↑
 {region}: id↑ ~ name↑

$$\begin{array}{l}
 11. \text{Constant}^S \\
 \mathcal{X}A: [] \mapsto \overline{B} \\
 \mathcal{X}: \overline{A} \sim \overline{B} \\
 \mathcal{X}: \overline{A} \sim \overline{B} \updownarrow \\
 \hline
 \mathcal{X}: [] \mapsto \overline{B}
 \end{array}$$

$X = \{\}$
 $A = \text{id}$
 $B = \text{region}$

{}: id↑ ~ region ↓ **missing**

Mapping between Set-based and List-based ODs

- List-based bODs and set-based bODs are equivalent

→ there exists a mapping between the representations

$$\bar{X} \mapsto \bar{Y} \quad \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{\quad} \end{array} \quad \begin{array}{l} X: [] \mapsto \bar{A} \\ X: \bar{A} \sim \bar{B} \end{array}$$

- But NO 1:1-mapping

Mapping from List-based to Set-based ODs [1,7]



$\bar{X} \mapsto \bar{Y}$ is valid iff $\bar{X} \mapsto \bar{X}\bar{Y}$ and $\bar{X} \sim \bar{Y}$ are valid.

No splits

No swaps

$\forall \bar{A} \in \bar{Y}$:
 $X: [] \mapsto \bar{A}$ must be valid

+

$\forall i \in \{1, \dots, |\bar{X}|\}$ and $\forall j \in \{1, \dots, |\bar{Y}|\}$:
 $\{X_1, \dots, X_{i-1}, Y_1, \dots, Y_{j-1}\}: \bar{X}_i \sim \bar{Y}_j$ must be valid

• Examples $[A \uparrow, B \downarrow] \mapsto [C \uparrow, D \uparrow]$

$[B \downarrow, A \uparrow] \mapsto [C \uparrow]$

- $\{A, B\}: [] \mapsto C \uparrow$
- $\{A, B\}: [] \mapsto D \uparrow$
- $\{ \}: A \uparrow \sim C \uparrow$
- $\{A\}: B \downarrow \sim C \uparrow$
- $\{C\}: A \uparrow \sim D \uparrow$
- $\{A, C\}: B \downarrow \sim D \uparrow$

- $\{A, B\}: [] \mapsto C \uparrow$
- $\{ \}: B \uparrow \sim C \uparrow$
- $\{B\}: A \downarrow \sim C \uparrow$

List-based ODs can share set-based ODs in their mapping

- Polynomial complexity $(|\bar{X}| \times |\bar{Y}|)$
- # set-based ODs is not necessarily larger than # list-based ODs

Mapping from Set-based to List-based ODs [1,7]



- Constant bODs:

$$X: [] \mapsto \bar{A}$$



\forall permutations \bar{X}' of \bar{X} :
 $\bar{X}' \mapsto \bar{X}'A$ must be valid

- Example

$$\{A, B\}: [] \mapsto C \uparrow$$

$$\begin{array}{ll} [A \uparrow, B \uparrow] \mapsto [A \uparrow, B \uparrow, C \uparrow] & [B \uparrow, A \uparrow] \mapsto [B \uparrow, A \uparrow, C \uparrow] \\ [A \uparrow, B \downarrow] \mapsto [A \uparrow, B \downarrow, C \uparrow] & [B \uparrow, A \downarrow] \mapsto [B \uparrow, A \downarrow, C \uparrow] \\ [A \downarrow, B \uparrow] \mapsto [A \downarrow, B \uparrow, C \uparrow] & [B \downarrow, A \uparrow] \mapsto [B \downarrow, A \uparrow, C \uparrow] \\ [A \downarrow, B \downarrow] \mapsto [A \downarrow, B \downarrow, C \uparrow] & [B \downarrow, A \downarrow] \mapsto [B \downarrow, A \downarrow, C \uparrow] \end{array}$$

- Factorial complexity ($|X|!$)
- Set-based is more concise representation

- Order compatible bODs:

$$X: \bar{A} \sim \bar{B}$$



\forall permutations \bar{X}' of \bar{X} :
 $\bar{X}'A \sim \bar{X}'B$ must be valid

- Example

$$\{A\}: B \uparrow \sim C \downarrow$$

$$\begin{array}{l} [A \uparrow, B \uparrow] \sim [A \uparrow, C \downarrow] \\ [A \downarrow, B \uparrow] \sim [A \downarrow, C \downarrow] \end{array}$$

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