

IT Systems Engineering | Universität Potsdam

Memory-Based Cloud Architectures

(Or: Technical Challenges for

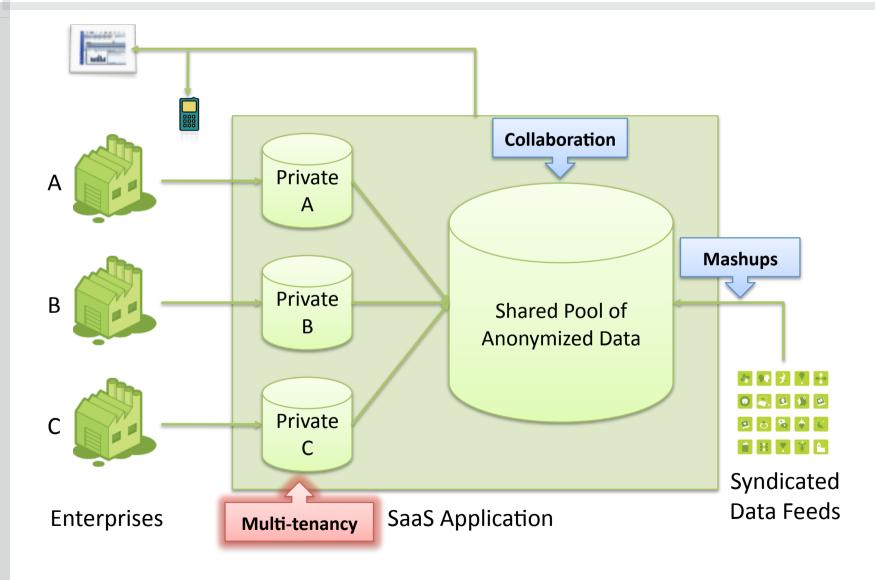
OnDemand Business Software)

Jan Schaffner

Enterprise Platform and Integration Concepts Group



Example: Enterprise Benchmarking



Definition



Cloud Computing

=

Data Center + API



What to take home from this talk?

Answers to four questions:

- Why are memory based architectures great for cloud computing?
- How predictable is the behavior of an in-memory column database?
- Does virtualization have a negative impact on in-memory databases?
- How do I assign tenants to servers in order to manage fault-tolerance and scalability?



Why are memory based architectures great for cloud computing?



Numbers everyone should know

■ L1 cache reference 0.5 ns

Branch mispredict5 ns

L2 cache reference7 ns

Mutex lock/unlock25 ns

Main memory reference100 ns (in 2008)

Compress 1K bytes with Zippy 3,000 ns

Send 2K bytes over 1 Gbps network 20,000 ns

Read 1 MB sequentially from memory 250,000 ns

Round trip within same datacenter 500,000 ns (in 2008)

■ Disk seek 10,000,000 ns

Read 1 MB sequentially from network 10,000,000 ns

Read 1 MB sequentially from disk
20,000,000 ns

■ Send packet CA \rightarrow Netherlands \rightarrow CA 150,000,000 ns

Source: Jeff Dean

Memory should be the system of record



- Typically disks have been the system of record
 - □ Slow → wrap them in complicated caching and distributed file systems to make them perform
 - Memory used as cache all over the place but it can be invalidated when something changes on disk
- Bandwidth:

□ Disk: 120 MB/s/controller

 \square DRAM (x86 + FSB): 10.4 GB/s/board

□ DRAM (Nehalem): 25.6 GB/s/socket

Latency:

Disk: 13 milliseconds (up to seconds when queuing)

□ InfiniBand: 1-2 microseconds

DRAM: 5 nanoseconds



High-end networks vs. disks

Maximum bandwidths:

| Hard Disk | 100-120 MB/s |
|-------------------------------|--|
| SSD | 250 MB/s |
| Serial ATA II | 600 MB/s |
| 10 GB Ethernet | 1204 MB/s |
| | |
| InfiniBand | 1250 MB/s (4 channels) |
| InfiniBand PCIe Flash Storage | 1250 MB/s (4 channels) 1400 MB/s |
| | |



Even more numbers...

| | Туре | Device / Medium | Latency | Throughput | |
|--------------|--------------|----------------------------------|---------------------------------------|------------------------------------|--|
| | Storage | L1 cache read (local) | 1.3 ns | 364.8 Gbps | |
| Same core | Storage | L2 cache read (local) | 3.4 ns | 248.8 Gbps | |
| | Storage | L3 cache read (local) | 13 ns | 209.6 Gbps | |
| | Storage | L1 cache read (remote, same die) | 13 - 28.3 ns | 75.2 - 154.4 Gbps | |
| Same die | Storage | L2 cache read (remote, same die) | 13 - 25.5 ns | 105.6 - 157.6 Gbps | |
| | Storage | L3 cache read (remote, same die) | 13 - 22.2 ns | 157.6 - 209.6 Gbps | |
| ard | Storage | L1 cache read (remote, via QPI) | 58 - 109 ns | 44.8 - 72 Gbps | |
| Same board | Storage | L2 cache read (remote, via QPI) | 58 - 109 ns | 44.8 - 73.6 Gbps | |
| San | Storage | L3 cache read (remote, via QPI) | 58 - 109 ns | 44.8 - 73.6 Gbps | |
| | Storage | DRAM (Nehalem) | 65 - 106 ns | 160 - 256 Gbps / socket | |
| ne | Interconnect | SATA 3.0 | at least 1 μs | 6 Gbps | |
| achi | Interconnect | Serial Attached SCSI | at least 1 μs | 6 Gbps | |
| Same machine | Interconnect | PCI Express | 3.8 - 5 μs | 4 Gbps x number of lanes | |
| San | Storage | Magnetical disk read / write | 3.2 - 13 ms | 0.96 - 1.12 Gbps | |
| | Storage | Solid State Disk read | 65 μs | 1.9 Gbps | |
| | Interconnect | RDMA over InfiniBand | 1 - 3 μs | 2.5 - 10 Gbps x number of channels | |
| ¥ | Interconnect | RDMA over iWARP | 6 μs | 10 Gbps / link | |
| Network | Interconnect | 10Gb Ethernet | 20 μs | 10 Gbps / link | |
| | Interconnect | Fibre channel | 3 - 10 μs (add 1 ms per 100 km) | 8 Gbps / channel | |



Designing a database for the cloud

- Disks are the limiting factor in contemporary database systems
 - Sharing a high performance disk on a machine/cluster/cloud is fine/troublesome/miserable
 - While one guy is fetching 100 MB/s, everyone else is waiting
- Claim: Two machines + network is better than one machine + disk
 - Log to disk on a single node:
 - $> 10,000 \,\mu s$ (not predictable)
 - Transactions only in memory but on two nodes:
 - < 600 µs (more predictable)
- Concept: Design to the strengths of cloud (redundancy) rather than their weaknesses (shared anything)



Design choices for a cloud database

- No disks (in-memory delta tables + async snapshots)
- Multi-master replication
 - Two copies of the data
 - Load balancing both reads and (monotonic) writes
 - □ (Eventual) consistency achieved via MVCC (+ Paxos, later)
- High-end hardware
 - Nehalem for high memory bandwidth
 - Fast interconnect
- Virtualization
 - Ease of deployment/administration
 - Consolidation/multi-tenancy



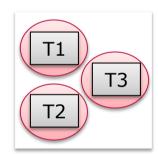
Why consolidation?

- In-memory column databases are ideal for mixed workload processing
- **But:** In a SaaS environment it seems costly to give everybody their private NewDB box
- How much consolidation is possible?
 - 3 years worth of sales records from our favorite
 Fortune 500 retail company
 - □ 360 million records
 - Less than 3 GB in compressed columns in memory
 - Next door is a machine with 2 TB of DRAM
 - (Beware of overhead)

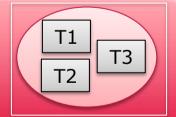
Multi-tenancy in the database – four different options



- No multi-tenancy one VM per tenant
 - Ex.: RightNow has 3000 tenants in 200 databases (2007):
 3000 vs. 200 Amazon VMs cost \$2,628,000 vs. \$175,200/year
 - Very strong isolation
- **Shared machine** one database process per tenant
 - Scheduler, session manager and transaction manager need
 live inside the individual DB processes: IPC for synchronization
 - Good for custom extensions, good isolation



- Shared instance one schema instance per tenant
 - Must support large numbers of tables
 - Must support online schema extension and evolution

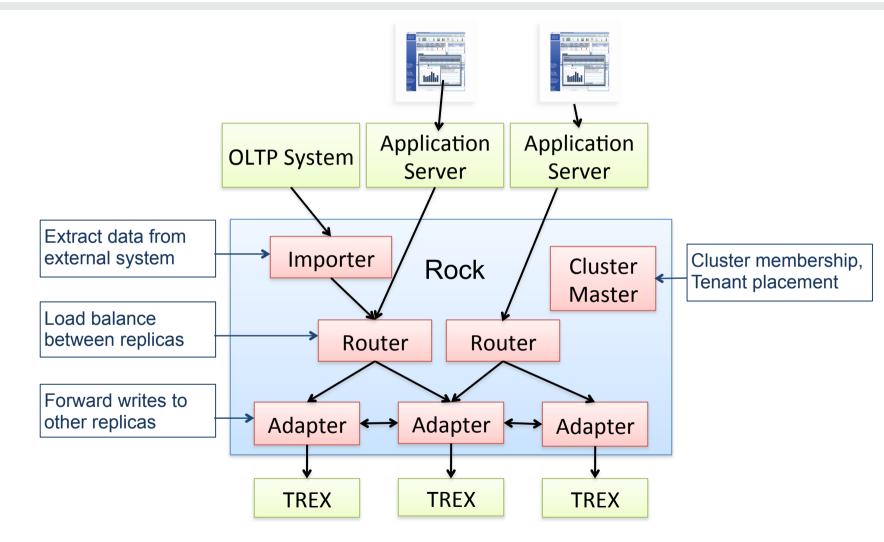


- Shared table use a tenant_id column and partitioning
 - Bad for custom extensions, bad isolation
 - Hard to backup/restore/migrate individual tenants



Putting it all together: Rock cluster architecture









How predictable is the behavior of an in-memory column database?



What does "predictable" mean?

- Traditionally, database people are concerned with the questions of type "how do I make a query faster?"
- In a SaaS environment, the question is "how do I get a fixed (low) response time as cheap as possible?"
 - Look at throughput
 - Look at quantiles (e.g. 99-th percentile)



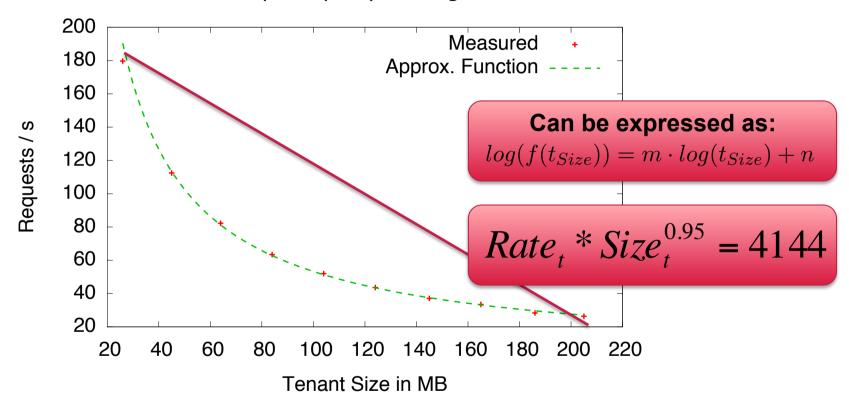


- Example formulation of desired performance:
 - Response time goal "1 second in the 99-th percentile"
 - Average response time around 200 ms
 - Less than 1% of all queries exceed 1,000 ms
 - Results in a maximum number of concurrent queries before response time goal is violated



System capacity

Fixed amount of data split equally among all tenants



- Capacity ≈ bytes scanned per second (there is a small overhead when processing more requests)
- In-memory databases behave very linearly!



Workload

- Tenants generally have different rates and sizes
- For a given set of T tenants (on one server) define

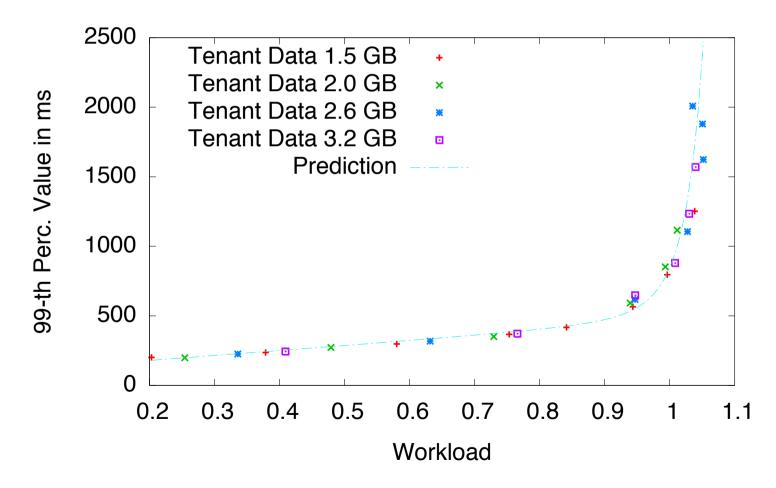
$$Workload = \sum_{t \in T} \frac{Rate_t * Size_t^{0.95}}{4144}$$

- When Workload = 1
 - System runs at it's maximum throughput level
 - Further increase of workload will result in violation of response time goal



Response time

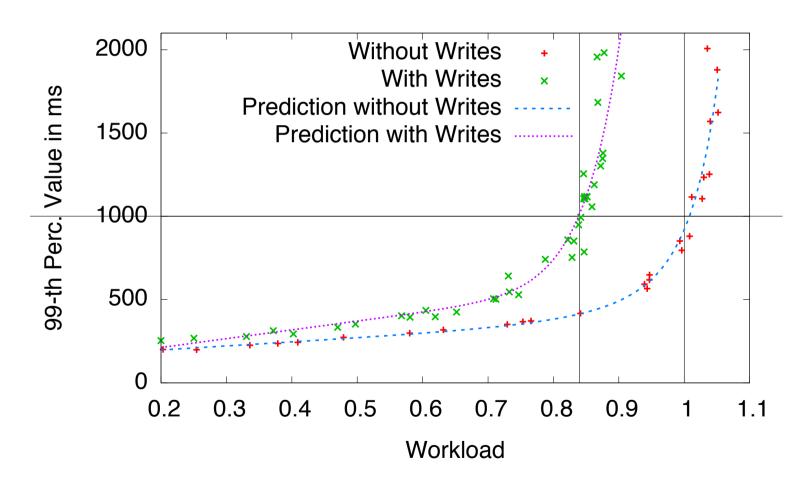
- Different amounts of data and different request rates ("assorted mix")
- Workload is varied by scaling the request rates





Impact of writes

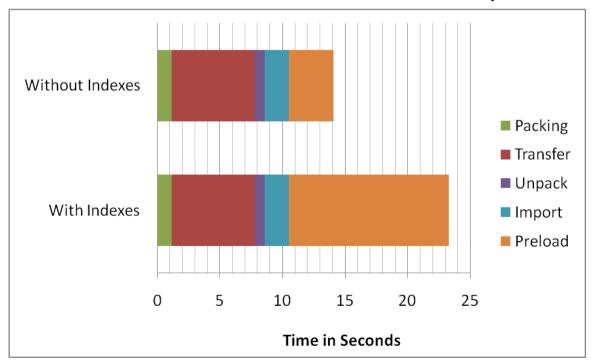
Added periodic batch writes (fact table grows by 0.5% every 5 minutes)





Why is predictability good?

- Ability to plan and perform resource intensive tasks during normal operations:
 - Upgrades
 - Merges
 - Migrations of tenants in the cluster (e.g. to dynamically re-balance the load situation in the cluster)



Cost breakdown for migration of tenants





Cloud Computing

Data Center + API



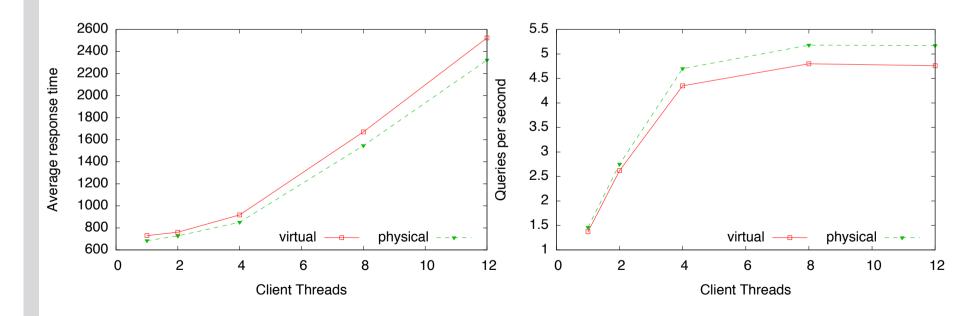


Does virtualization have a negative impact on in-memory databases?



Impact of virtualization

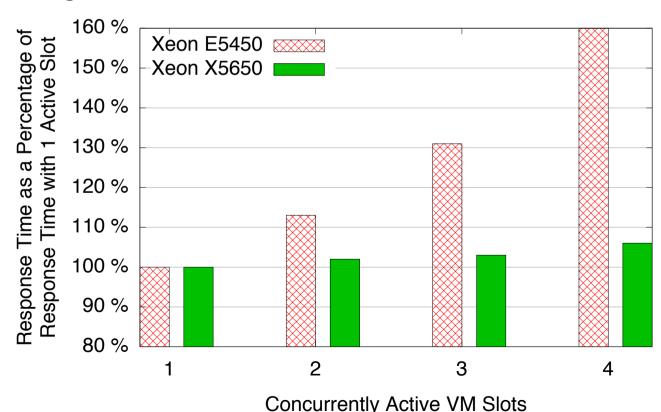
- Run multi-tenant OLAP benchmark on either:
 - one TREX instance directly on the physical host vs.
 - one TREX instance inside VM on the physical host
- Overhead is approximately 7% (both in response time and throughput)

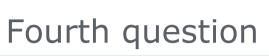




Impact of virtualization (contd.)

- Virtualization is often used to get "better" system utilization
 - What happens when a physical machine is split into multiple VMs?
 - □ Burning CPU cycles does not hurt → memory bandwidth is the limiting factor







How do I assign tenants to servers in order to manage fault-tolerance and scalability?

Why is it good to have multiple copies of the data?



- Scalability beyond a certain number of concurrently active users
- High availability during normal operations
- Alternating execution of resource-intensive operations (e.g. merge)
- Rolling upgrades without downtime
- Data migration without downtime
- Reminder: Two in-memory copies allow faster writes and are more predictable than one in-memory copy plus disk

Downsides:

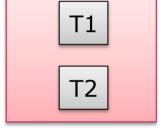
- Response time goal might be violated during reco Really?
- You need to plan for twice the capacity



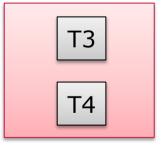


Conventional Mirrored Layout

T1

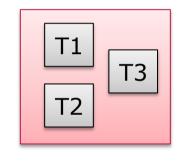


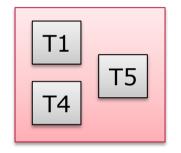


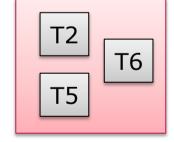


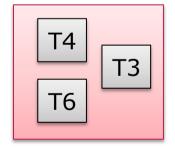
If a node fails, all work moves to one other node. The system must be **100% over-provisioned**.

Interleaved Layout









If a node fails, work moves to many other nodes. Allows **higher utilization** of nodes.



Handcrafted best case

- Perfect placement:
 - 100 tenants
 - 2 copies/tenant
 - All tenants have same size
 - 10 tenants/server



Mirrored



Interleaved

- Perfect balancing (same load on all tenants):
 - □ 6M rows (204 MB compressed) of data per tenant
 - The same (increasing) number of users per tenant

■ No writes

| L NO WITCES | Mirrored | Interleaved | Improvement |
|--------------------------|------------|-------------|-------------|
| No failures | 4218 users | 4506 users | 7% |
| Periodic single failures | 2265 users | 4250 users | 88% |

Throughput before violating response time goal



Requirements for placement algorithm

- An optimal placement algorithm needs to cope with multiple (conflicting) goals:
 - Balance load across servers
 - Achieve good interleaving
- Use migrations consciously for online layout improvements (no big bang cluster re-organization)
- Take usage patterns into account
 - Request rates double during last week before end of quarter
 - Time-zones, Christmas, etc.



Conclusion

Answers to four questions:



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How predictable is the behavior of an in-memory column database?



Does virtualization have a negative impact on in-memory databases?



How do I assign tenants to servers in order to manage fault-tolerance and scalability?

Questions?